

# Working with PLLs in PrimeTime – avoiding the “phase locked *oops*”

Paul Zimmer

Zimmer Design Services  
1375 Sun Tree Drive  
Roseville, CA 95661

[paulzimmer@zimmerdesignservices.com](mailto:paulzimmer@zimmerdesignservices.com)

[website: www.zimmerdesignservices.com](http://www.zimmerdesignservices.com)

## ABSTRACT

PLLs play an important role in modern high-speed designs, especially when configured for clock tree insertion delay cancellation (IDC). Modeling the behavior of such PLLs accurately in PrimeTime can be a challenge. This paper discusses basic modeling techniques for both standard and multiplier IDC PLLs, duty cycle modeling, jitter and skew, and on-chip-variation effects. The classic OCV/PLL excess pessimism problem will be explained and examined, and a couple of workarounds will be discussed, including a novel new technique developed by the author.

## Table of contents

1	Introduction.....	3
2	The basic insertion delay cancellation (idc) PLL.....	4
2.1	The problem – too much delay .....	4
2.2	The solution – the insertion delay cancellation PLL.....	5
2.3	Timing the basic idc PLL in PrimeTime .....	6
2.4	The IDC multiplier PLL .....	15
2.5	The PLL model itself .....	<a href="#">2324</a>
2.6	Performance considerations .....	<a href="#">2425</a>
3	Duty Cycle .....	<a href="#">2526</a>
3.1	Internal clocks (other than PLLs).....	<a href="#">2526</a>
3.2	Primary input clocks and PLLs .....	<a href="#">2728</a>
3.3	Applying this to our multiplier pll circuit .....	<a href="#">343533</a>
3.4	When to use these techniques .....	<a href="#">404139</a>
4	Jitter .....	<a href="#">414240</a>
4.1	Jitter, skew, and uncertainty.....	<a href="#">414240</a>
4.2	My definition of jitter .....	<a href="#">414240</a>
4.3	Sources of jitter.....	<a href="#">424341</a>
4.4	Effects of jitter on different sorts of paths .....	<a href="#">424341</a>
4.5	Modeling jitter with set_clock_uncertainty.....	<a href="#">464745</a>
4.6	Applying jitter specs to the example circuit – simple case .....	<a href="#">464745</a>
4.7	Generated clocks.....	<a href="#">535452</a>
4.8	What about falling edges?.....	<a href="#">555654</a>
4.9	Applying jitter specs to the example circuit – complex case .....	<a href="#">565755</a>
5	On-chip Variation .....	<a href="#">636462</a>
5.1	The classic OCV case .....	<a href="#">636462</a>
5.2	Enter CRPR .....	<a href="#">666765</a>
5.3	OCV and PLLs .....	<a href="#">686967</a>
5.4	The OCV/PLL excess pessimism problem .....	<a href="#">747573</a>
6	OCV/PLL excess pessimism workarounds.....	<a href="#">787977</a>
6.1	Forcing OCV off on the fb path.....	<a href="#">787977</a>
6.2	Referencing the i/os to the feedback clock .....	<a href="#">959694</a>
6.3	The shell game.....	<a href="#">103105103</a>
7	Conclusion .....	<a href="#">110111109</a>
8	Acknowledgements.....	<a href="#">111112110</a>
9	References.....	<a href="#">112113111</a>
10	Appendix.....	<a href="#">113114112</a>
10.1	The PLL model itself .....	<a href="#">113114112</a>
10.2	Why I do it my way.....	<a href="#">128130128</a>
10.3	Modeling duty cycle using set_clock_latency early/late.....	<a href="#">129130128</a>
10.4	Modeling duty cycle using set_clock_latency min/max.....	<a href="#">133135133</a>

## 1 Introduction

PLLs are an absolutely essential tool in high-speed design. Their ability to nearly zero out the delay of a large clock tree allows for much higher speed inter-chip communication.

However, modeling PLLs in static timing analysis is tricky. Many of the details have been glossed over or ignored in the past, but as circuit feature sizes decrease and speeds increase, it is no longer acceptable to depend on over-constrained budgets or just plain dumb luck to be sure the PLL will work in the intended application. Certain effects, such as on-chip variation and signal integrity analysis, make correct operation of the circuit absolutely dependent on a complete, accurate static timing analysis – and this includes the PLL.

## 2 The basic insertion delay cancellation (idc) PLL

PLLs are common in modern high-speed designs. The inner workings of the pll are more the domain of analog sorcerers, but the basic idea is this: “the pll makes its output do whatever is necessary to make its ref and fb inputs match in phase and frequency”.

There are many uses for PLLs. The one I’ll be discussing here is what I call the “insertion delay cancellation” pll.

### 2.1 The problem – too much delay

Here is a very basic i/o circuit. Data and clock are sent in and data comes out. The outgoing data is clocked externally by the same clock that is sent to the circuit.

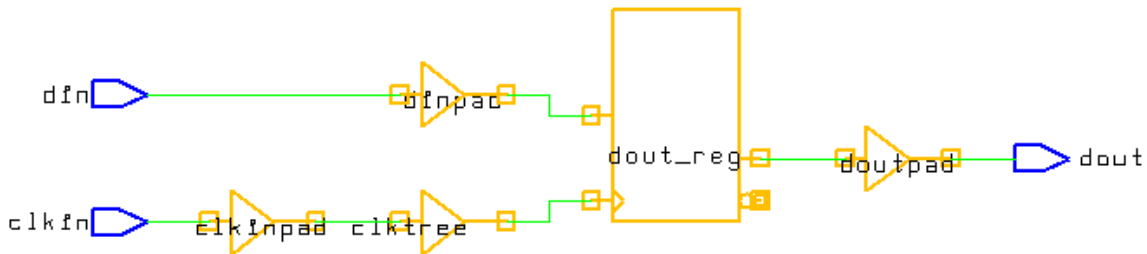


Figure 2-1

And here’s what the timing waveform looks like:

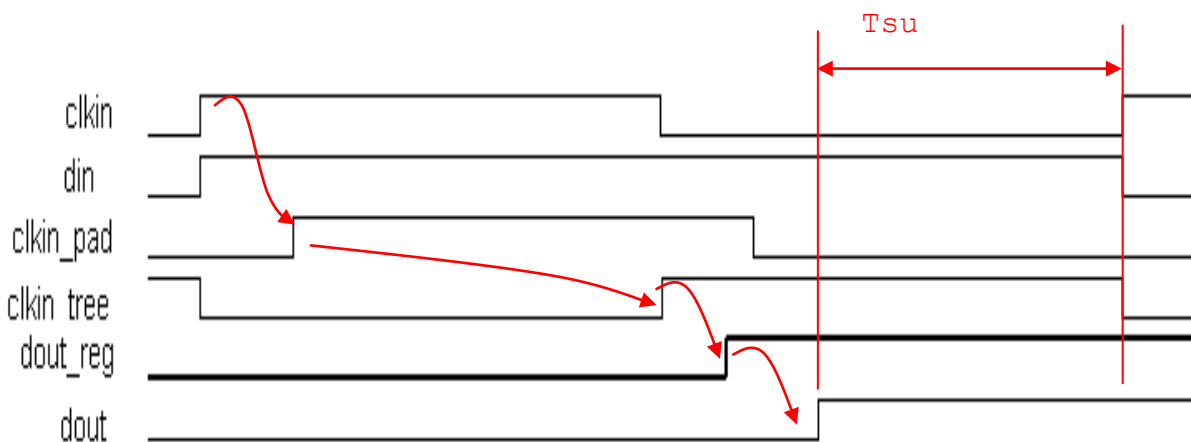


Figure 2-2

It is clear that all of the delays in generating the data (clock pad, clock tree insertion delay, etc)

will reduce the available setup time of dout and therefore limit the frequency of operation.

## 2.2 The solution – the insertion delay cancellation PLL

To get around this problem, designers often insert a PLL into the clock path. The configuration looks like this:

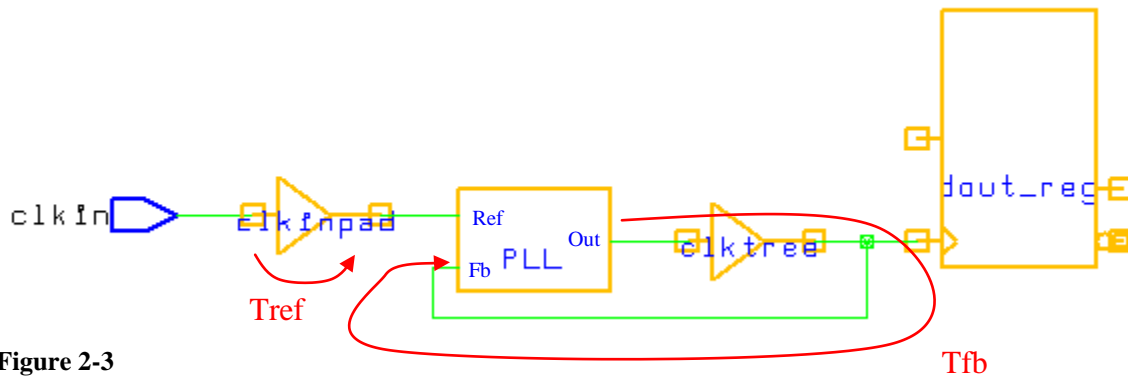


Figure 2-3

The Fb input of the pll is connected to the *end* of the clock tree.

Remember that the pll must drive its output such that its ref and fb inputs match in phase and frequency. The frequency part is easy – it just matches clk\_in's frequency, which reappears on the Fb pin. What about the phase?

Ignore Tref for now. If the clock at Ref arrived at time zero, the pll would have to drive its output *back in time* by Tfb to make the signal arrive at the Fb pin at time zero. So, it would launch the clock at time  $-Tfb$ . But, clk\_in appears at the pll Ref pin at time Tref. So, the pll doesn't have to drive its output back by the full Tfb, but only by  $Tfb - Tref$ . In other words, it launches its clock at time  $Tref - Tfb$ . After passing through the feedback loop, the clock arrives at the Fb input at time Tref, which is what we want.

The net effect of this is that the flops on the end of the clock tree get clocked at *almost* time zero. They get clocked at time Tref. If we add a little extra delay to the feedback path to match Tref, we can get the flops clocked at time zero relative to clkIn. This allows us to “cancel” the insertion delay of the clock tree. So, a more realistic idc pll might look like this:

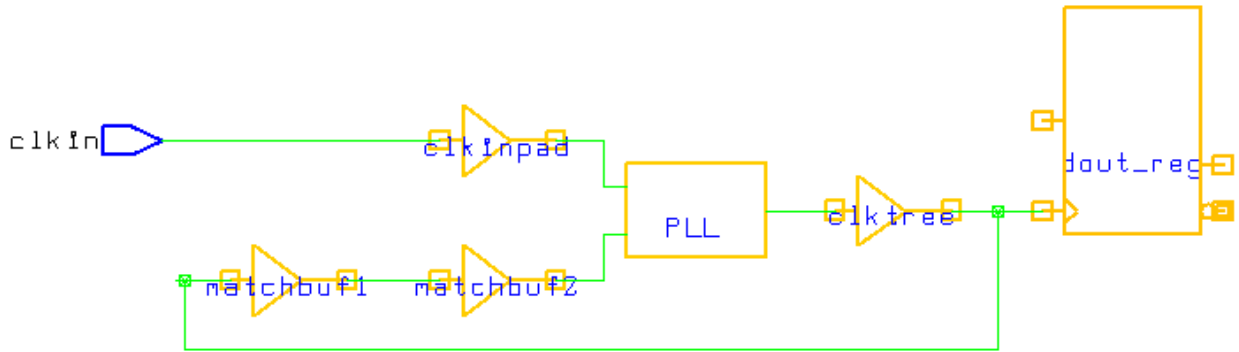


Figure 2-4

## 2.3 Timing the basic idc PLL in PrimeTime

So, how do we go about modeling this in Primetime? The most straightforward approach (although not the only approach, as we shall see later) is to model what happens in real life – create a clock on the pll output pin with a frequency to match clkIn, and then “launch” this pll output clock at time Tref-Tfb. We then allow the two clocks to time against one another. The “launch” is accomplished with the set\_clock\_latency command, so the commands would look something like this:

```
create_clock -period 10.0 -name clkIn [get_ports clkIn]
set_propagated_clock clkIn

create_clock -period 10.0 -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout

set_clock_latency -source \
  [expr $_ref_delay - $_fb_delay] \
  [get_clocks pllout]
```

The problem is, how do we get the values of \$\_ref\_delay and \$\_fb\_delay?

Reference [4] gives one approach. I use a different approach. I use get\_timing\_paths to get the values. I think mine is simpler to code. It also takes advantage of the fact that, by creating all the clocks first, it is possible to extract the feedback delay as an arrival attribute on the fb pin. Also, reference [4] puts the values into files that are then sourced by multiple Primetime runs. I’m not that trusting, so I force PT to calculate the value every time the script runs.

My approach is like this:

1. Create all the clocks, including the pll output clock (but don't set the source latency yet).
2. Get Tref and Tfb using get\_timing\_paths.
3. Calculate the source latency (Tref - Tfb) and apply it to the pll output clock. The source latency value will usually be negative.
4. Use set\_input\_delay and set\_output\_delay with the reference clock to specify the i/o timing.
5. Allow the clocks to time against each other (don't do set\_false\_path between them)

Here's an example circuit. I have added a flip-flop on the end of the clock tree. The flop clocks data in through the din pad, and sends data out through the dout pad. Notice also that this flop is on a slightly different branch of the clock tree than the one the pll feedback is hooked up to. They share clktree\_root, but not clktree1 (to the flop) or clktree2 (to the pll feedback).

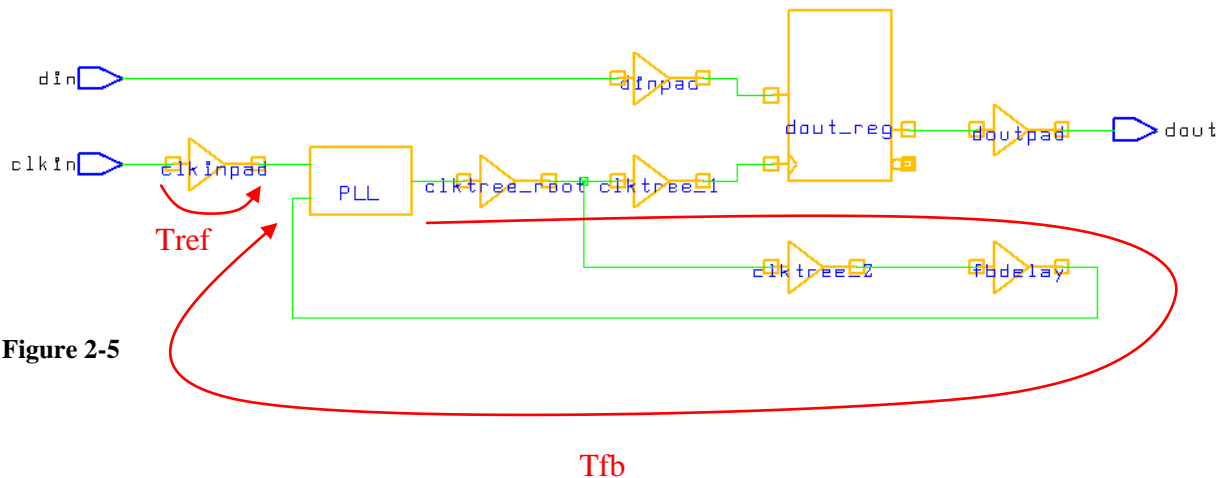


Figure 2-5

First, we'll create the clocks. Doing this first makes it easier to get the delay values we want:

```
create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout
```

Because the clocks exists, we can get the timing values we want by just getting the “arrival” attribute on the path:

```
set _path [get_timing_paths -delay max_rise \  
  -from [get_ports clk_in] \  
  -to [get_pins PLL/CKREF] \  
]  
  
set _ref_delay [get_attribute $_path arrival]  
  
set _path [get_timing_paths -delay max_rise \  
  -from [get_pins PLL/OUT] \  
  -to [get_pins PLL/FB] \  
]  
  
set _fb_delay [get_attribute $_path arrival]
```

Note the use of “-delay max\_rise”. We want a rise delay because the pll runs on rising edges (most do, anyway). Without on-chip-variation (OCV), the min\_rise and max\_rise values should be the same. OCV effects will be discussed later.

Also note that \_ref\_delay is set by getting the arrival attribute on the path returned by get\_timing\_paths, which is a collection. I expect this collection to have only one path in it. If I have done something wrong and the collection has more than one path, the “set \_ref\_delay” command will cause an error. I call this poor-man’s error checking...

When I run this script, I get:

```
pt_shell> echo $_ref_delay  
1.000000  
pt_shell> echo $_fb_delay  
3.300000  
pt_shell>
```



Let's verify this using report\_timing. We'll need to set the variable timing\_report\_unconstrained\_paths to "true" to see a result.

```
pt_shell> set timing_report_unconstrained_paths true
true
pt_shell> report_timing -input_pins -delay max_rise -from [get_ports clkkin] -
to [get_pins PLL/CKREF]
*****
Report : timing
        -path full
        -delay max_rise
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****
```

```
Startpoint: clkkin (clock source 'clkkin')
Endpoint: PLL/CKREF (internal pin)
Path Group: (none)
Path Type: max
```

Point	Incr	Path
clock source latency	0.00	0.00
clkkin (in)	0.00	0.00 r
clkkinpad/I (bufbd1)	0.00	0.00 r
clkkinpad/Z (bufbd1)	1.00 *	1.00 r
PLL/CKREF (DUMMYPLL)	0.00	1.00 r
data arrival time		1.00

(Path is unconstrained)

```

pt_shell> report_timing -input_pins -delay max_rise -from [get_pins PLL/OUT] -
to [get_pins PLL/FB]
*****
Report : timing
        -path full
        -delay max_rise
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

```

```

Startpoint: PLL/OUT (clock source 'pllout')
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_2/I (bufbd1)	0.00	2.20 r
clktree_2/Z (bufbd1)	0.10 *	2.30 r
fbdelay/I (bufbd1)	0.00	2.30 r
fbdelay/Z (bufbd1)	1.00 *	3.30 r
PLL/FB (DUMMYPLL)	0.00	3.30 r
data arrival time		3.30

(Path is unconstrained)

Looks correct.  $\$_{ref\_delay}$  is 1.0 and  $\$_{fb\_delay}$  is 3.3.

Now we'll apply this source latency:

```

set_clock_latency -source \
  [expr $_ref_delay - $_fb_delay] \
  [get_clocks pllout]

```

When we do this, we'll get the following warning:

```
Warning: Negative clock latency specified: -2.3 (UITE-150)
```

This is harmless. We normally expect a negative source delay on the pll output clock (because the fb path is usually longer than the refclk path).

To make sure we got what we wanted, do `report_clock -skew`:

```
pt_shell> report_clock -skew
*****
Report : clock_skew
Design : idc_pll_example
Version: V-2004.06
*****
```

Object	Min Condition Source Latency				Max Condition Source Latency				Rel_clk
	Early_r	Early_f	Late_r	Late_f	Early_r	Early_f	Late_r	Late_f	
pllout	-2.30	-2.30	-2.30	-2.30	-2.30	-2.30	-2.30	-2.30	--

Just to make absolutely sure this all worked, here's some code that verifies that the arrival time at the CKREF pin matches the arrival time at the FB pin within some small rounding error:

```
# Verify
# Get the new fb path delay
set_path [get_timing_paths -delay max_rise \
  -to [get_pins PLL/FB] \
]

set_new_fb_delay [get_attribute $_path arrival]

# Get the ref clock delay
set_path [get_timing_paths -delay max_rise \
  -from [get_ports clkkin] \
  -to [get_pins PLL/CKREF] \
]

set_new_ref_delay [get_attribute $_path arrival]

set_diff [expr $_new_ref_delay - $_new_fb_delay]

if { ($_diff > 0.01) || ($_diff < -0.01) } {
  echo "Error: Difference between FB and REF pins out of range!"
  echo "  Difference is $_diff"
  echo "  _ref_delay is $_ref_delay"
  echo "  _fb_delay is $_fb_delay"
} else {
  echo "PLL timing verified!"
}
```

When we run this, we get:

```
PLL timing verified!
```

Now we need to constrain the data paths. The din/dout paths are part of an interface that is referenced to clkin, so we'll set the input and output delays accordingly.

```
set_input_delay -max 8.0 -clock clkin [get_ports din]
set_input_delay -min 0.5 -clock clkin [get_ports din]
set_output_delay -max 2.0 -clock clkin [get_ports dout]
set_output_delay -min [expr -1.0 * 0.5] -clock clkin [get_ports dout]
```

Now let's look at the timing. Here's the schematic again, with the delays of each buffer shown:

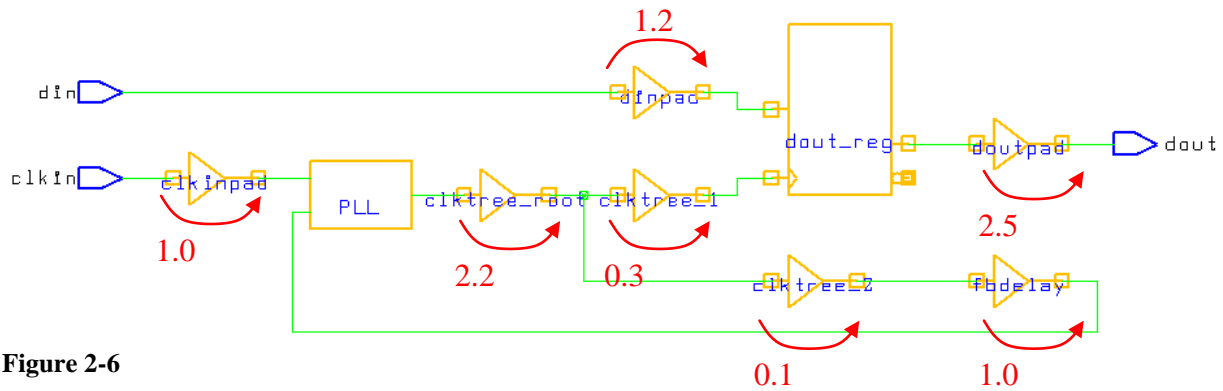


Figure 2-6

A couple of things to notice about this. I have made the fbdelay match the clkInpad delay exactly, but the clktree\_1 and clktree\_2 buffers don't match exactly. The clock tree has 0.2ns of *real* skew.

First we'll look at the din timing:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****
```

```
Startpoint: din (input port clocked by clkin)
Endpoint: dout_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type: max
```

Point	Incr	Path
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
dout_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.30	7.70
PLL/OUT (DUMMYPLL)	0.00	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
dout_reg/CP (dfnrb1)	0.00	10.20 r
library setup time	-0.08	10.12
data required time		10.12
-----		
data required time		10.12
data arrival time		-9.20
-----		
slack (MET)		0.92

If the insertion delay cancellation were perfect, we would expect the clock to arrive at dout\_reg/CP at time 10.0 (one full period of the clock). Instead, it arrives at 10.2. Why? Because the clock tree had skew. The clktree\_2 buffer was 0.2ns faster than the clktree\_1 buffer. This means the feedback delay from the pll was slightly less than it should have been, resulting in a late arriving clock.

Now let's look at the dout timing:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -to dout
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****
```

```
Startpoint: dout_reg (rising edge-triggered flip-flop clocked by pllout)
Endpoint:   dout (output port clocked by clkin)
Path Group: clkin
Path Type:  max
```

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
dout_reg/CP (dfnrb1)	0.00	0.20 r
dout_reg/Q (dfnrb1)	0.32	0.52 f
doutpad/I (bufbd1)	0.00	0.52 f
doutpad/Z (bufbd1)	2.50 *	3.02 f
dout (out)	0.00	3.02 f
data arrival time		3.02
clock clkin (rise edge)	10.00	10.00
clock network delay (propagated)	0.00	10.00
output external delay	-2.00	8.00
data required time		8.00
-----		
data required time		8.00
data arrival time		-3.02
-----		
slack (MET)		4.98

In this case, we're launching data from the pll domain and capturing it in the clkin domain. So, the capture clock is at 10.0. But the launch clock at dout\_reg/CP is 0.20 – once again it is 0.2ns late (for the same reason explained above).

Both of these path reports are correct – we have modeled the pll behavior correctly.

## 2.4 The IDC multiplier PLL

Now lets try a slightly more complicated example. Consider this circuit:

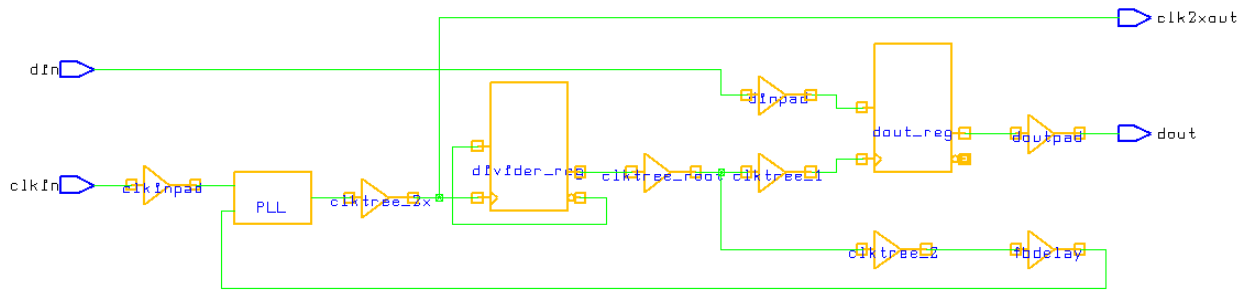


Figure 2-7

Remember that the pll must drive its output such that its ref and fb inputs match in phase and frequency. Since there is a divide-by 2 in the fb path, the pll will have to drive its output at twice the frequency of clkIn to make the frequency match at the fb pin. This is a frequency multiplier pll.

There are several reasons why you might want to do this. You might need the 2x clock internally for other functions, for example. In that case, the i/os would be connected to the divide-by output. Or you might be doing something really wierd, like running the data at 2x speeds anchored by the 1x reference clock.

How do we model this in PrimeTime? Well, it turns out that the procedure outlined above still works. Here's the circuit again with all the delays shown:

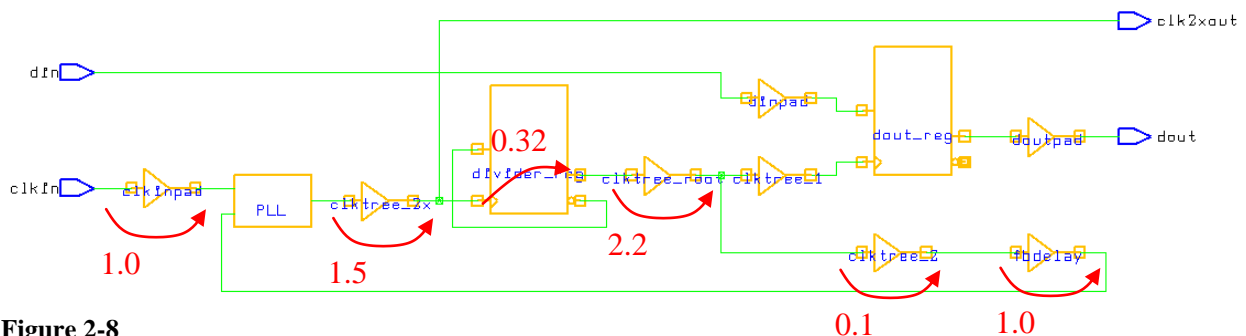


Figure 2-8

The first step was, “create all the clocks”. The first two clocks are the same as before, except that the period of the pll clock is now half that of clkIn:

```
create_clock -period 10.0 -name clkIn [get_ports clkIn]
set_propagated_clock clkIn
```

```
create_clock -period [expr 10.0 / 2] -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout
```

But there's another clock here. The divide-by 2 flop output is also a clock. It is a generated clock from pllout, with a divide-by of 2:

```
create_generated_clock \  
  -source [get_pins PLL/OUT] \  
  -name divclk \  
  -divide_by 2 \  
  [get_pins divider_reg/Q]  
set_propagated_clock divclk
```

Now we need to fetch the Tref and Tfb values.

```
set_path [get_timing_paths -delay max_rise \  
  -from [get_ports clkIn] \  
  -to [get_pins PLL/CKREF] \  
]  
set_ref_delay [get_attribute $_path arrival]  
set_path [get_timing_paths -delay max_rise \  
  -from [get_pins PLL/OUT] \  
  -to [get_pins PLL/FB] \  
]  
set_fb_delay [get_attribute $_path arrival]  
echo $_fb_delay
```

But if we look at the result, there's a problem:

```
pt_shell> echo $_ref_delay  
1.000000  
pt_shell> echo $_fb_delay  
  
pt_shell>
```

Let's look at the timing reports for the fb path.

```
pt_shell> set timing_report_unconstrained_paths true  
true  
pt_shell> report_timing -delay max_rise -from [get_pins PLL/OUT] -to  
[get_pins PLL/FB]  
*****  
Report : timing  
  -path full  
  -delay max_rise  
  -max_paths 1  
Design : idcm_pll_example  
Version: V-2004.06  
*****
```

No Paths.

The path no longer exists. That's because the divider flop is breaking the path.



Try again with just the pll FB pin endpoint:

```
pt_shell> report_timing -delay max_rise -to [get_pins PLL/FB]
*****
Report : timing
        -path full
        -delay max_rise
        -max_paths 1
Design : idcm_pll_example
Version: V-2004.06
*****
```

```
Startpoint: divider_reg/Q
             (clock_source 'divclk')
Endpoint:   PLL/FB (internal pin)
Path Group: (none)
Path Type:  max
```

Point	Incr	Path
clock source latency	1.82	1.82
divider_reg/Q (dfnrb1)	0.00	1.82 r
clktree_root/Z (bufbd1)	2.20 H	4.02 r
clktree_2/Z (bufbd1)	0.10 H	4.12 r
fbdelay/Z (bufbd1)	1.00 H	5.12 r
PLL/FB (DUMMYPLL)	0.00	5.12 r
data arrival time		5.12

(Path is unconstrained)

But is this correct? Let's see where the "clock source latency" comes from:

```
pt_shell> report_timing -delay max_rise -from [get_pins divider_reg/CP] -
to [get_pins divider_reg/Q]
*****
Report : timing
        -path full
        -delay max_rise
        -max_paths 1
Design : idcm_pll_example
Version: V-2004.06
*****
```

```
Startpoint: divider_reg
             (rising edge-triggered flip-flop clocked by pllout)
Endpoint:   divider_reg/Q
             (internal pin)
Path Group: (none)
Path Type:  max
```

Point	Incr	Path
clock network delay (propagated)	1.50	1.50
divider_reg/CP (dfnrb1)	0.00	1.50 r
divider_reg/Q (dfnrb1)	0.32	1.82 r
data arrival time		1.82

(Path is unconstrained)

```

pt_shell> report_timing -delay max_rise -from [get_pins PLL/OUT] -to [get_pins
divider_reg/CP] -input_pins
*****
Report : timing
        -path full
        -delay max_rise
        -input_pins
        -max_paths 1
Design : idcm_pll_example
Version: V-2004.06
*****

```

```

Startpoint: PLL/OUT (clock source 'pllout')
Endpoint: divider_reg/CP
          (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_2x/I (bufbd1)	0.00	0.00 r
clktree_2x/Z (bufbd1)	1.50 *	1.50 r
divider_reg/CP (dfnrb1)	0.00	1.50 r
data arrival time		1.50

(Path is unconstrained)

So, the source latency of 1.82 is the prop delay from the PLL/OUT pin to the divider\_reg/CP pin (1.50) plus the CP->Q rise delay through the divider flop (0.32). That's correct.

We can see this more clearly if we use `-path_type full_clock_expanded` with just the endpoint:

```
pt_shell> report_timing -delay max_rise -to [get_pins PLL/FB] -path_type
full_clock_expanded
*****
Report : timing
        -path full_clock_expanded
        -delay max_rise
        -max_paths 1
Design : idcm_pll_example
Version: V-2004.06-SP1
*****

Startpoint: divider_reg/Q
             (clock source 'divclk')
Endpoint:   PLL/FB (internal pin)
Path Group: (none)
Path Type:  max

Point              Incr      Path
-----
clock pllout (source latency)      -4.12    -4.12
PLL/OUT (DUMMYPLL)                  0.00    -4.12 r
clktree_2x/Z (bufbd1)                1.50 *   -2.62 r
divider_reg/Q (dfnrb1) (gclock source) 0.32 H   -2.30 r
divider_reg/Q (dfnrb1)                0.00    -2.30 r
clktree_root/Z (bufbd1)              2.20 H   -0.10 r
clktree_2/Z (bufbd1)                 0.10 H    0.00 r
fbdelay/Z (bufbd1)                   1.00 H    1.00 r
PLL/FB (DUMMYPLL)                    0.00    1.00 r
data arrival time                    1.00
-----
(Path is unconstrained)
```

So, let's try the `_fb_delay` code without the `-from` switch:

```
set_path [get_timing_paths -delay max_rise \
  -to [get_pins PLL/FB] \
]

set _fb_delay [get_attribute $_path arrival]
```

And the value is now correct:

```
pt_shell> echo $_fb_delay
5.120000
pt_shell>
```

If a simple “-to” is too open-ended for your taste, it turns out that you also get the correct result when you do “-from [get\_pins divider\_reg/Q]”, like this:

```
set_path [get_timing_paths -delay max_rise \
  -from [get_pins divider_reg/Q] \
  -to [get_pins PLL/FB] \
  ]

set _fb_delay [get_attribute $_path arrival]

pt_shell> echo $_fb_delay
5.120000
pt_shell>
```

Now we can apply the source latency as before:

```
# Set the source latency
set_clock_latency -source \
  [expr $_ref_delay - $_fb_delay] \
  [get_clocks pllout]
```

The resulting latency should be -4.12 (1.0 – 5.12):

```
pt_shell> report_clock -skew
*****
Report : clock_skew
Design : idcm_pll_example
Version: V-2004.06
*****
```

Object	Min Condition Source Latency				Max Condition Source Latency				Rel_clk
	Early_r	Early_f	Late_r	Late_f	Early_r	Early_f	Late_r	Late_f	
pllout	-4.12	-4.12	-4.12	-4.12	-4.12	-4.12	-4.12	-4.12	--
divclk	1.82	1.82	1.82	1.82	1.82	1.82	1.82	1.82	--

Looks right.

The verify code will similarly have to change (remove the `-from` from the `fb` path or replace with the `-from` from the `divider_reg`):

```
# Verify
# Get the new fb path delay
set _path [get_timing_paths -delay max_rise \
  -to [get_pins PLL/FB] \
]

set _new_fb_delay [get_attribute $_path arrival]

# Get the ref clock delay
set _path [get_timing_paths -delay max_rise \
  -from [get_ports clk] \
  -to [get_pins PLL/CKREF] \
]

set _new_ref_delay [get_attribute $_path arrival]

set _diff [expr $_new_ref_delay - $_new_fb_delay]

if { ($_diff > 0.01) || ($_diff < -0.01)
} {
  echo "Error: Difference between FB and REF pins out of range!"
  echo "  Difference is $_diff"
  echo "  _new_ref_delay is $_new_ref_delay"
  echo "  _new_fb_delay is $_new_fb_delay"
} else {
  echo "PLL timing verified!"
}
```

And the verify works:

```
PLL timing verified!
```

Apply the i/o constraints as before:

```
set_input_delay -max 8.0 -clock clk [get_ports din]
set_input_delay -min 0.5 -clock clk [get_ports din]
set_output_delay -max 2.0 -clock clk [get_ports dout]
set_output_delay -min [expr -1.0 * 0.5] -clock clk [get_ports dout]
```

Now let's look at some i/o timing:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : idcm_pll_example
Version: V-2004.06
*****
```

Startpoint: din (input port clocked by clkin)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by divclk)  
 Path Group: divclk  
 Path Type: max

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
dout_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
clock divclk (rise edge)	10.00	10.00
clock pllout (source latency)	-4.12	5.88
PLL/OUT (DUMMYPLL)	0.00	5.88 r
clktree_2x/I (bufbd1)	0.00	5.88 r
clktree_2x/Z (bufbd1)	1.50 *	7.38 r
divider_reg/CP (dfnrb1)	0.00	7.38 r
divider_reg/Q (dfnrb1) (gclock source)	0.32 *	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
dout_reg/CP (dfnrb1)	0.00	10.20 r
library setup time	-0.08	10.12
data required time		10.12
-----		
data required time		10.12
data arrival time		-9.20
-----		
slack (MET)		0.92

Notice that the slack is exactly the same as it was with the earlier non-multiplying idc pll example. This is to be expected. I made the pll run at 2x, but then clocked dout\_reg with the divide-by 2 clock, so nothing changed relative to the i/o timing. Extra delay in the clock path (the divider) is just “insertion delay” and gets cancelled out by the pll.

The same is true of the dout path:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -to dout
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : idcm_pll_example
Version: V-2004.06
*****
```

Startpoint: dout\_reg (rising edge-triggered flip-flop clocked by divclk)  
 Endpoint: dout (output port clocked by clkin)  
 Path Group: clkin  
 Path Type: max

Point	Incr	Path
clock divclk (rise edge)	0.00	0.00
clock pllout (source latency)	-4.12	-4.12
PLL/OUT (DUMMYPLL)	0.00	-4.12 r
clktree_2x/I (bufbd1)	0.00	-4.12 r
clktree_2x/Z (bufbd1)	1.50 *	-2.62 r
divider_reg/CP (dfnrb1)	0.00	-2.62 r
divider_reg/Q (dfnrb1) (gclock source)	0.32 *	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
dout_reg/CP (dfnrb1)	0.00	0.20 r
dout_reg/Q (dfnrb1)	0.32 *	0.52 r
doutpad/I (bufbd1)	0.00	0.52 r
doutpad/Z (bufbd1)	2.50 *	3.02 r
dout (out)	0.00	3.02 r
data arrival time		3.02
clock clkin (rise edge)	10.00	10.00
clock network delay (propagated)	0.00	10.00
output external delay	-2.00	8.00
data required time		8.00
data required time		8.00
data arrival time		-3.02
slack (MET)		4.98

## 2.5 The PLL model itself

Up to this point, I haven't talked about the model of the PLL itself. In the examples above, the model was an "empty shell" like this:

```

module DUMMYPLL (
    OUT,
    FB,
    CKREF
);

output OUT;
input CKREF;
input FB;

endmodule

```

It could also have been allowed to default to a "black box" by Primetime. The results would be the same.

This works fine if the flow uses SDF, but with parasitics this isn't sufficient. When using parasitics, Primetime needs to know more about the driving and load characteristics of the PLL itself. This requires a model of the PLL itself, which can be a bit tricky. This is covered in appendix 10.1. If you're using parasitics, please read this appendix carefully.

## 2.6 Performance considerations

Recall that the basic flow outlined above was:

1. Create all the clocks, including the pll output clock (but don't set the source latency yet).
2. Get Tref and Tfb using `get_timing_paths`.
3. Calculate the source latency ( $T_{ref} - T_{fb}$ ) and apply it to the pll output clock. The source latency value will usually be negative.
4. Use `set_input_delay` and `set_output_delay` with the reference clock to specify the i/o timing.
5. Allow the clocks to time against each other (don't do `set_false_path` between them)

And we could add another:

6. Run the verify code to ensure that the arrival time at CKREF matches the arrival time at FB.

The drawback to this is performance. Steps 2 and 6 cause timing updates. Once you have confidence in the flow, you could always turn off step 6. But step 2 always causes a timing update.

That's not a big problem if you only have 1 pll. But if you have several, you'll want to structure the code such that step 1 covers all clocks of all PLLs, step 2 fetches all Tref and Tfb values for all pll, etc. The code is much harder to read, but the performance impact can be very noticeable.



### 3 Duty Cycle

The term duty cycle refers to the percent of time that the clock signal is high. It matters when signals are sampled on the falling edge of the clock. We will examine two cases: internal clocks and primary input clocks / plls.

#### 3.1 Internal clocks (other than PLLs)

It is important to note that PT will handle the duty cycle calculations without user intervention for most internally-generated clocks. For example, take the 2x multiplier pll circuit shown above and add another flop in the data path. The first flop (din\_reg) will be clocked by the pll divide-by 2 output as before. But the second flop (dout\_reg) will be clocked by an inverted version of this clock.

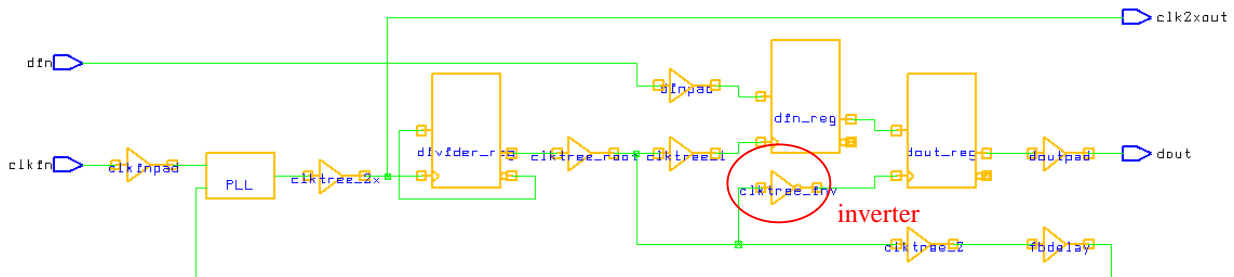


Figure 3-1

The script to handle this is exactly the same as before. To illustrate the effect, I will force the rise and fall delays of the divider\_reg to be different:

```
set_annotated_delay -cell -from divider_reg/CP -to divider_reg/Q -rise 0.32
set_annotated_delay -cell -from divider_reg/CP -to divider_reg/Q -fall 0.20
```

Since the signal falls more quickly than it rises, the high time will be reduced.

Now, report the timing between the 2 flops:

```
pt_shell> report_timing -from din_reg -to dout_reg -path_type
full_clock_expanded -input_pins
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_internal
Version: V-2003.12-SP1
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by divclk)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by divclk')  
 Path Group: divclk  
 Path Type: max

Point	Incr	Path
-----		
clock divclk (rise edge)	0.00	0.00
clock pllout (source latency)	-4.12	-4.12
PLL/OUT (DUMMYPLL)	0.00	-4.12 r
clktree_2x/I (bufbd1)	0.00	-4.12 r
clktree_2x/Z (bufbd1)	1.50 *	-2.62 r
divider_reg/CP (dfnrb1)	0.00	-2.62 r
divider_reg/Q (dfnrb1) (gclock source)	0.32 *	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/Q (dfnrb1) <-	0.32	0.52 r
dout_reg/D (dfnrb1)	0.00	0.52 r
data arrival time		0.52
clock divclk' (rise edge)	5.00	5.00
clock pllout (source latency)	-4.12	0.88
PLL/OUT (DUMMYPLL)	0.00	0.88 r
clktree_2x/I (bufbd1)	0.00	0.88 r
clktree_2x/Z (bufbd1)	1.50 *	2.38 r
divider_reg/CP (dfnrb1)	0.00	2.38 r
divider_reg/Q (dfnrb1) (gclock source)	0.20 *	2.58 f
clktree_root/I (bufbd1)	0.00	2.58 f
clktree_root/Z (bufbd1)	2.20 *	4.78 f
clktree_inv/I (inv0d2)	0.00	4.78 f
clktree_inv/ZN (inv0d2)	0.03	4.81 r
dout_reg/CP (dfnrb1)	0.00	4.81 r
library setup time	-0.08	4.73
data required time		4.73
-----		
data required time		4.73
data arrival time		-0.52
-----		
slack (MET)		4.21

A couple of things to notice about this trace. First, the capture edge is divclk` at time 5.0. The refclk period is 10.0. We're running the pll as a 2x multiplier, so its period is 5.0. These flops are on the div\_clk, which is divide-by-2 on the pll – which gives it a period of 10.0. So the fall edge of divclk is at 5.0 – we're timing a half-cycle path.

The second thing to notice is that the capture clock uses the *fall* edge through divider\_reg/CP->Q. Since this is set to 0.20 (instead of 0.32 for the rise edge), the capture clock will occur 0.12ns sooner. This reflects the non-ideal duty cycle of the generated clock. The effect will be to reduce the slack on this rise-to-fall path.

The important thing to notice, however, is that the duty cycle effect is handled automatically by the tool.

### 3.2 Primary input clocks and PLLs

Primary input clocks and pll output clocks are *not* handled automatically by the tool, since it can't deduce what the effect would be. You have to tell it via the `-waveform` option on `create_clock`. But there's a snag. The incoming clock (or the pll output clock) spec doesn't tell you that the duty cycle is, say, "55%". It usually says the duty cycle is, say, 50% +/- 5%. For internal clocks, the duty cycle is what it is, but for primary input clocks and pll output clocks, it's a range.

Let's take this simple example:

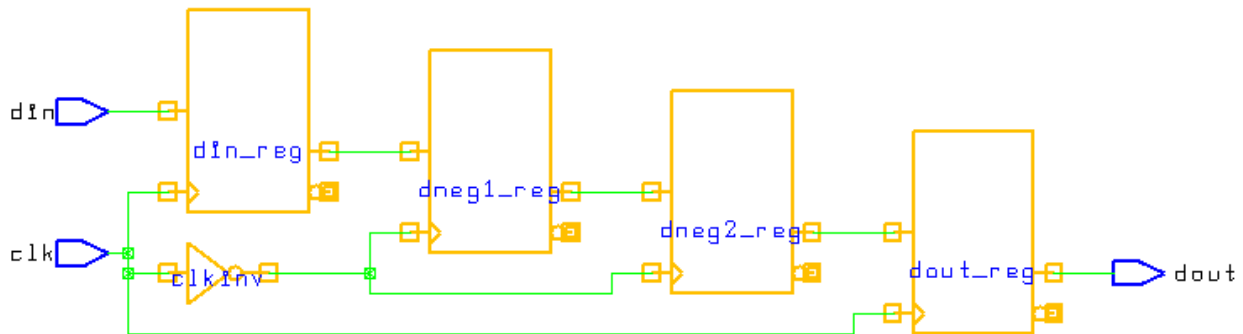


Figure 3-2

This circuit has paths from the rising clock to the falling clock, from the falling clock to the falling clock, and from the falling clock to the rising clock.

So, how do we model this in PrimeTime? Well, duty cycle seems like an uncertainty, so your first thought might be to use `set_clock_uncertainty`. Unfortunately, this command doesn't allow you to set uncertainty between the rising and falling edges *of the same clock*. It has the `-rise_to`, `-fall_from` options, but the man page says you can only use this for *inter-clock* uncertainty – that is, uncertainty between two clocks.

<added 3/17/2005>

The man page says:

-from from clock -to to clock

These two options specify the source and destination clocks for interclock

uncertainty. You must specify either

the pair of -from/-rise from/-fall from

and -to/-rise to/-fall to, or

object list; you cannot specify both.

...

-fall Indicates that uncertainty applies to only the falling edge of the destination clock. By default, the uncertainty applies to both rising and falling edges. This option is valid only for interclock uncertainty, and is now obsolete. Unless you need this option for backward-compatibility, use -fall to instead.

Sure enough, if you try to use edge switches with intraclock uncertainty, you get an error:

```
pt shell> set clock uncertainty -setup -rise 0.2 [get clocks clkin]
Error: Cannot specify '-rise or -fall' with 'clock_list'. (CMD-001)
0
```

But PT will accept this:

```
pt shell> set clock uncertainty -setup -rise from [get clocks clkin] -fall to [get clocks
clkin] 0.2
1
```

So, it appears that you can specify single-edge uncertainty intra-clock by using the inter-clock syntax. The man page doesn't say you can't make both clocks in the inter-clock syntax be the same.

It appears to work:

```
pt shell> report timing -from din reg -to dout reg
*****
Report : timing
-path full
-delay max
-max paths 1
Design : simple
Version: W-2004.12-SP1
Date : Thu Mar 17 18:26:54 2005
*****
```

Warning: There are 3 invalid start points. (UITE-416)
Warning: There are 2 invalid end points for constrained paths. (UITE-416)

```
Startpoint: din reg (rising edge-triggered flip-flop clocked by clkin)
Endpoint: dout reg (rising edge-triggered flip-flop clocked by clkin')
Path Group: clkin
Path Type: max
```

```
Point                               Incr                               Path
-----
```

Formatted  
Formatted

Formatted

clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
din reg/CP (dfnrb1)	0.00	0.00 r
din reg/Q (dfnrb1) <-	0.31	0.31 r
dout reg/D (dfnrb1)	0.00	0.31 r
data arrival time		0.31
-----		
clock clkin' (rise edge)	5.00	5.00
clock network delay (propagated)	0.02	5.02
inter-clock uncertainty	-0.20	4.82
dout reg/CP (dfnrb1)		4.82 r
library setup time	-0.09	4.73
data required time		4.73
-----		
data required time		4.73
data arrival time		-0.31
-----		
slack (MET)		4.42

Still, there are disadvantages to this approach. Rolling the duty cycle into uncertainty will make it difficult to separate duty cycle issues from other uses of set clock uncertainty described later.

On the other hand, the uncertainty method may have some advantages in runtime (fewer clocks) and possibly in SI analysis over the technique described below.

It is possible to model duty cycle using “set\_clock\_latency -fall”, but it doesn’t work out-of-the-box. You have to (ab)use some of the on-chip-variation features for something unrelated to on-chip-variation, and isn’t too clean (it is explained in Appendix [10.3/4]).

Until a few years ago, there seemed to be no really clean way to model this except to run the script twice – once for each of the extreme duty cycle waveforms (I call them min\_high and max\_high). With the introduction of multiclock propagation capabilities in PT, we can now do this in a single run.

Since we are going to use multiclock propagation, we first need to turn it on (it’s off by default):

```
set timing_enable_multiple_clocks_per_reg true
```

Now we can create the two clocks. I have defined a single variable “\_duty\_cycle\_min” to specify what the minimum high period can be. I’ll set “\_duty\_cycle\_max” to 1.0 minus this value (but you could set it to some other value if you wanted to):

```
set _period 10.0
set _duty_cycle_min 0.40
set _duty_cycle_max [expr 1.0 - $_duty_cycle_min]
create_clock -period $_period -name clk_minhigh \
  -waveform [list 0 [expr $_period * $_duty_cycle_min]] \
  [get_ports clk]
set_propagated_clock clk_minhigh

create_clock -period $_period -name clk_maxhigh \
  -waveform [list 0 [expr $_period * $_duty_cycle_max]] \
  -add \
  [get_ports clk]
set_propagated_clock clk_maxhigh
```

Notice the use of “-add” on the second create\_clock. This is necessary to have both clocks exist on the same pin.

Now, these clocks can never exist at the same time, so we don’t want them timing against one another. We could do set\_false\_path from each to the other, but instead let’s use the new set\_clock\_groups command. The two clocks can never coexist, so I’ll use the –exclusive option:

```
set_clock_groups -exclusive \
  -group {clk_minhigh} \
  -group {clk_maxhigh}
```

Now let’s look at the timing reports. We’ll time the path from the din\_reg to the dneg1\_reg first:

```
pt_shell> report_timing -input_pins -path_type full_clock -from din_reg -to
dneg1_reg
*****
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_piclck
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by clk\_maxhigh)  
 Endpoint: dnegl\_reg (rising edge-triggered flip-flop clocked by clk\_maxhigh')  
 Path Group: clk\_maxhigh  
 Path Type: max

Point	Incr	Path
-----		
clock clk_maxhigh (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clk (in)	0.00	0.00 r
din_reg/CP (dfnrb1)	0.00	0.00 r
din_reg/Q (dfnrb1) <-	0.31 *	0.31 r
dnegl_reg/D (dfnrb1)	0.00	0.31 r
data arrival time		0.31
clock clk_maxhigh' (rise edge)	6.00	6.00
clock source latency	0.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dnegl_reg/CP (dfnrb1)	0.00	7.00 r
library setup time	-0.09	6.91
data required time		6.91
-----		
data required time		6.91
data arrival time		-0.31
-----		
slack (MET)		6.60

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by clk\_minhigh)  
 Endpoint: dnegl\_reg (rising edge-triggered flip-flop clocked by clk\_minhigh')  
 Path Group: clk\_minhigh  
 Path Type: max

Point	Incr	Path
-----		
clock clk_minhigh (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clk (in)	0.00	0.00 r
din_reg/CP (dfnrb1)	0.00	0.00 r
din_reg/Q (dfnrb1) <-	0.31 *	0.31 r
dnegl_reg/D (dfnrb1)	0.00	0.31 r
data arrival time		0.31
clock clk_minhigh' (rise edge)	4.00	4.00
clock source latency	0.00	4.00
clk (in)	0.00	4.00 f
clkinv/I (inv0d2)	0.00	4.00 f
clkinv/ZN (inv0d2)	1.00 *	5.00 r
dnegl_reg/CP (dfnrb1)	0.00	5.00 r
library setup time	-0.09	4.91
data required time		4.91
-----		
data required time		4.91
data arrival time		-0.31
-----		
slack (MET)		4.60

We get two timing reports – one for clk\_minhigh and one for clk\_maxhigh. Since we are launching data from a rising edge (both clocks have the same rising edge waveform) and capturing with the falling edge, the worst case for setup is the shorter duty cycle clock – clk\_minhigh. As you can see, clk\_minhigh does indeed have less slack. Its fall edge (that weird “clock clk\_minhigh’ (rise edge)” is PT-speak for falling edge) is at time 4.0 – 40% of the 10ns period.

Next we’ll look at the falling-edge to falling-edge path.

```
pt_shell> report_timing -input_pins -path_type full_clock -from dneg1_reg -to
dneg2_reg
```

```
*****
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_piclk
Version: V-2004.06
*****
```

```
Startpoint: dneg1_reg (rising edge-triggered flip-flop clocked by
clk_maxhigh')
Endpoint: dneg2_reg (rising edge-triggered flip-flop clocked by
clk_maxhigh')
Path Group: clk_maxhigh
Path Type: max
```

Point	Incr	Path
clock clk_maxhigh' (rise edge)	6.00	6.00
clock source latency	0.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dneg1_reg/CP (dfnrb1)	0.00	7.00 r
dneg1_reg/Q (dfnrb1) <-	0.32 *	7.32 r
dneg2_reg/D (dfnrb1)	0.00	7.32 r
data arrival time		7.32
clock clk_maxhigh' (rise edge)	16.00	16.00
clock source latency	0.00	16.00
clk (in)	0.00	16.00 f
clkinv/I (inv0d2)	0.00	16.00 f
clkinv/ZN (inv0d2)	1.00 *	17.00 r
dneg2_reg/CP (dfnrb1)	0.00	17.00 r
library setup time	-0.09	16.91
data required time		16.91
data required time		16.91
data arrival time		-7.32
slack (MET)		9.59



Startpoint: dneg1\_reg (rising edge-triggered flip-flop clocked by clk\_minhigh')  
 Endpoint: dneg2\_reg (rising edge-triggered flip-flop clocked by clk\_minhigh')  
 Path Group: clk\_minhigh  
 Path Type: max

Point	Incr	Path
-----		
clock clk_minhigh' (rise edge)	4.00	4.00
clock source latency	0.00	4.00
clk (in)	0.00	4.00 f
clkinv/I (inv0d2)	0.00	4.00 f
clkinv/ZN (inv0d2)	1.00 *	5.00 r
dneg1_reg/CP (dfnrb1)	0.00	5.00 r
dneg1_reg/Q (dfnrb1) <-	0.32 *	5.32 r
dneg2_reg/D (dfnrb1)	0.00	5.32 r
data arrival time		5.32
clock clk_minhigh' (rise edge)	14.00	14.00
clock source latency	0.00	14.00
clk (in)	0.00	14.00 f
clkinv/I (inv0d2)	0.00	14.00 f
clkinv/ZN (inv0d2)	1.00 *	15.00 r
dneg2_reg/CP (dfnrb1)	0.00	15.00 r
library setup time	-0.09	14.91
data required time		14.91
-----		
data required time		14.91
data arrival time		-5.32
-----		
slack (MET)		9.59

Again, there are two traces, one for each clock. Although the launch and capture times are different, as you would expect, the slack values are the same.

Now let's look at the dneg2\_reg to dout\_reg path:

```
pt_shell> report_timing -input_pins -path_type full_clock -from dneg2_reg -to
dout_reg
*****
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_piclk
Version: V-2004.06
*****
```

Startpoint: dneg2\_reg (rising edge-triggered flip-flop clocked by clk\_maxhigh')  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by clk\_maxhigh)  
 Path Group: clk\_maxhigh  
 Path Type: max

Point	Incr	Path
clock clk_maxhigh' (rise edge)	6.00	6.00
clock source latency	0.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dneg2_reg/CP (dfnrb1)	0.00	7.00 r
dneg2_reg/Q (dfnrb1) <-	0.32 *	7.32 r
dout_reg/D (dfnrb1)	0.00	7.32 r
data arrival time		7.32
-----		
clock clk_maxhigh (rise edge)	10.00	10.00
clock source latency	0.00	10.00
clk (in)	0.00	10.00 r
dout_reg/CP (dfnrb1)	0.00	10.00 r
library setup time	-0.10	9.90
data required time		9.90
-----		
data required time		9.90
data arrival time		-7.32
-----		
slack (MET)		2.58

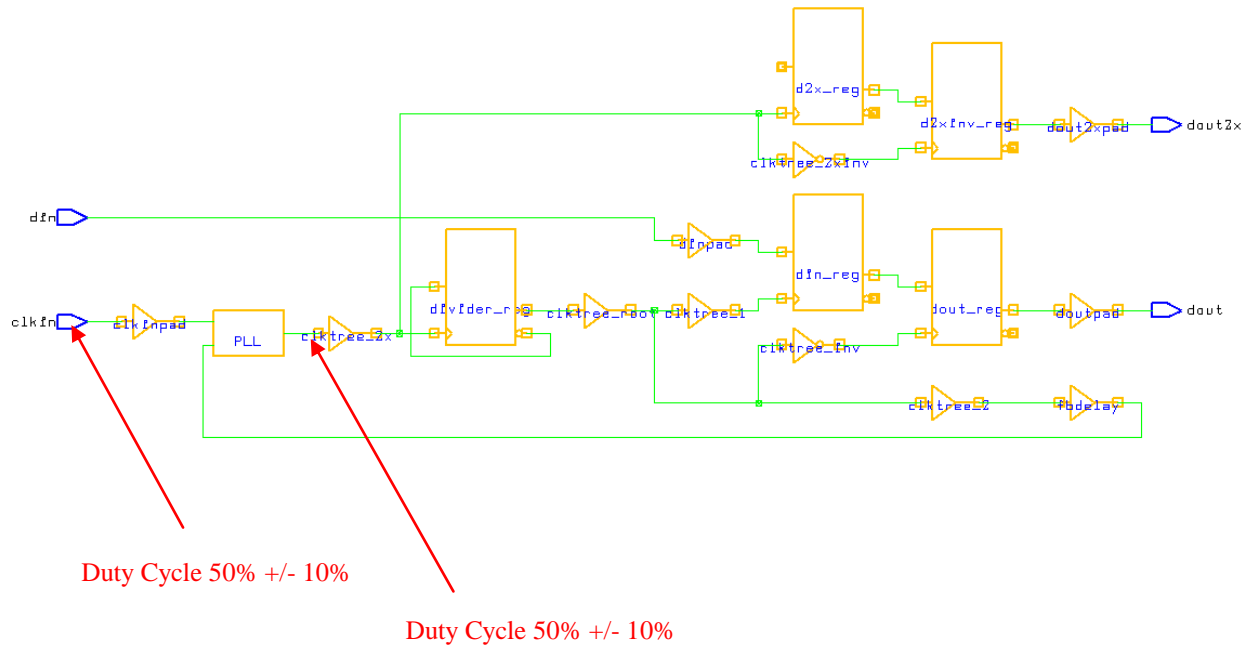
Since this is from the falling edge to the rising edge, the worst case path will be from the later falling edge, which is clk\_maxhigh. There is another trace for clk\_minhigh with 2 more ns of slack.

So, we can model duty cycle variation in a single pass using two “exclusive” clocks and multiclock propagation.

### 3.3 Applying this to our multiplier pll circuit

We can use this on our multiplier pll circuit in 2 places – the primary input clock clkin and the pll clock pllout.

We'll use the same circuit as before, except that I have added a pair of flops on the pllout clock, one rising edge and the other falling edge, to illustrate the duty cycle effect on pllout clock.



**Figure 3-3**

First we'll create the primary input clocks, one for each duty cycle case. We'll use a loop.

```

set _period 10.0
set _duty_cycle(min) 0.40
set _duty_cycle(max) [expr 1.0 - $_duty_cycle(min)]
foreach _dc {min max} {
    create_clock -period $_period -name clkin_${_dc}high \
        -waveform [list 0 [expr $_period * $_duty_cycle($_dc)]] \
        -add \
        [get_ports clkin]
    set_propagated_clock clkin_${_dc}high
}
set_clock_groups -exclusive \
    -group {clkin_minhigh} \
    -group {clkin_maxhigh}

```

Notice the “-add” switch again. Since the create\_clock is done in a loop, the -add will occur on both create\_clock commands, including the first one. This is harmless.

Creating the pll output clocks is similar:

```
set _period [expr 10.0 / 2]
set _duty_cycle(min) 0.40
set _duty_cycle(max) [expr 1.0 - $_duty_cycle(min)]
foreach _dc {min max} {
    create_clock -period $_period -name pllout_${_dc}high \
        -waveform [list 0 [expr $_period * $_duty_cycle($_dc)]] \
        -add \
        [get_pins PLL/OUT]
    set_propagated_clock pllout_${_dc}high
}
set_clock_groups -exclusive \
    -group {pllout_minhigh} \
    -group {pllout_maxhigh}
```

Now we need to create the generated divider clock. But now there are two clocks feeding the divider\_reg/CP pin. Which do we choose, and how do we tell PT about our choice?

Well, in this case, the choice is arbitrary. Since divider\_reg only runs on rising edges, it doesn't matter which pllout\_ clock we choose. If divider\_reg ran on falling edges, however, we would have to create two generated clocks – one for pllout\_minhigh and one for pllout\_maxhigh. We would then add these to their respective clock groups in set\_clock\_group. In fact, we could still do this even though divider\_reg runs on rising edges – the propagated clocks would be identical, but this is harmless. Some would argue that this is more consistent. It's a personal choice.

To keep things simple, we'll just create one generated clock. We'll use pllout\_maxhigh.

When multiple clocks feed into a point where we want to create a generated clock, it is necessary to tell PT which clock is the source. This is done using the “-master” switch. To use the “-master” switch, you must also use the “-add” switch, even though we're only creating one clock.

```
create_generated_clock \
    -source [get_pins PLL/OUT] \
    -name divclk \
    -divide_by 2 \
    -add \
    -master pllout_maxhigh \
    [get_pins divider_reg/Q]
set_propagated_clock divclk
```

The code to fetch the Tref and Tfb values is unchanged:

```
set _path [get_timing_paths -delay max_rise \  
  -from [get_ports clkIn] \  
  -to [get_pins PLL/CKREF] \  
]  
  
set _ref_delay [get_attribute $_path arrival]  
  
set _path [get_timing_paths -delay max_rise \  
  -from [get_pins divider_reg/Q] \  
  -to [get_pins PLL/FB] \  
]  
  
set _fb_delay [get_attribute $_path arrival]
```

Now to set the source latency. There are two pll output clocks now, and they need the same latency value. So, we do it in a loop again:

```
foreach _dc {minhigh maxhigh} {  
  set_clock_latency -source \  
    [expr $_ref_delay - $_fb_delay] \  
    [get_clocks pllout_${_dc}]  
}
```

Now we'll set the i/o timing. Since we now have two versions of clkIn, we'll have to create an input/output constraint for each. Also, I want to illustrate the duty cycle effects here. In this case, the duty cycle is on the external clock clkIn, so I'll do the i/o constraints relative to the falling edge.

```
foreach _dc {minhigh maxhigh} {  
  set_input_delay -max 4.0 -clock clkIn_${_dc} -clock_fall -add [get_ports  
din]  
  set_input_delay -min 0.5 -clock clkIn_${_dc} -clock_fall -add [get_ports  
din]  
  set_output_delay -max 1.0 -clock clkIn_${_dc} -clock_fall -add [get_ports  
dout]  
  set_output_delay -min [expr -1.0 * 0.5] -clock clkIn_${_dc} -clock_fall -add  
[get_ports dout]  
}
```

Now look at the timing from din:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din  
*****  
Report : timing  
  -path full_clock_expanded  
  -delay max  
  -input_pins  
  -max_paths 1  
Design : idcm_pll_dc  
Version: V-2004.06  
*****
```

Startpoint: din (input port clocked by clkin\_maxhigh)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by divclk)  
 Path Group: divclk  
 Path Type: max

Point	Incr	Path
clock clkin_maxhigh (fall edge)	6.00	6.00
clock network delay (propagated)	0.00	6.00
input external delay	4.00	10.00 r
din (in)	0.00	10.00 r
dinpad/I (bufbd1)	0.00	10.00 r
dinpad/Z (bufbd1)	1.20 *	11.20 r
dout_reg/D (dfnrb1)	0.00	11.20 r
data arrival time		11.20
-----		
clock divclk (rise edge)	10.00	10.00
clock pllout_maxhigh (source latency)		
	-4.12	5.88
PLL/OUT (DUMMYPLL)	0.00	5.88 r
clktree_2x/I (bufbd1)	0.00	5.88 r
clktree_2x/Z (bufbd1)	1.50 *	7.38 r
divider_reg/CP (dfnrb1)	0.00	7.38 r
divider_reg/Q (dfnrb1) (gclock source)		
	0.32 *	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
dout_reg/CP (dfnrb1)	0.00	10.20 r
library setup time	-0.08	10.12
data required time		10.12
-----		
data required time		10.12
data arrival time		-11.20
-----		
slack (VIOLATED)		-1.08

I have nworst defaulted to 1, so the report only shows the worst case – data launched by clkin\_maxhigh (latest possible falling edge of clkin) and captured by divclk. If I set nworst to 4, I'd see the clkin\_minhigh paths as well.

Now let's look at the pllout path:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
d2x_reg -to d2xinv_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : idcm_pll_dc
Version: V-2004.06
*****
```

Startpoint: d2x\_reg (rising edge-triggered flip-flop clocked by pllout\_maxhigh)

Endpoint: d2xinv\_reg (rising edge-triggered flip-flop clocked by pllout\_maxhigh')

Path Group: pllout\_maxhigh

Path Type: max

Point	Incr	Path
clock pllout_maxhigh (rise edge)	0.00	0.00
clock source latency	-4.12	-4.12
PLL/OUT (DUMMYPLL)	0.00	-4.12 r
clktree_2x/I (bufbd1)	0.00	-4.12 r
clktree_2x/Z (bufbd1)	1.50 *	-2.62 r
d2x_reg/CP (dfnrb1)	0.00	-2.62 r
d2x_reg/Q (dfnrb1) <-	0.34	-2.28 r
d2xinv_reg/D (dfnrb1)	0.00	-2.28 r
data arrival time		-2.28
-----		
clock pllout_maxhigh' (rise edge)	3.00	3.00
clock source latency	-4.12	-1.12
PLL/OUT (DUMMYPLL)	0.00	-1.12 f
clktree_2x/I (bufbd1)	0.00	-1.12 f
clktree_2x/Z (bufbd1)	1.50 *	0.38 f
clktree_2xinv/I (inv0d2)	0.00	0.38 f
clktree_2xinv/ZN (inv0d2)	0.04	0.42 r
d2xinv_reg/CP (dfnrb1)	0.00	0.42 r
library setup time	-0.08	0.33
data required time		0.33
-----		
data required time		0.33
data arrival time		2.28
-----		
slack (MET)		2.61

Startpoint: d2x\_reg (rising edge-triggered flip-flop clocked by pllout\_minhigh)

Endpoint: d2xinv\_reg (rising edge-triggered flip-flop clocked by pllout\_minhigh')

Path Group: pllout\_minhigh

Path Type: max

Point	Incr	Path
clock pllout_minhigh (rise edge)	0.00	0.00
clock source latency	-4.12	-4.12
PLL/OUT (DUMMYPLL)	0.00	-4.12 r
clktree_2x/I (bufbd1)	0.00	-4.12 r
clktree_2x/Z (bufbd1)	1.50 *	-2.62 r
d2x_reg/CP (dfnrb1)	0.00	-2.62 r
d2x_reg/Q (dfnrb1) <-	0.34	-2.28 r
d2xinv_reg/D (dfnrb1)	0.00	-2.28 r
data arrival time		-2.28
-----		
clock pllout_minhigh' (rise edge)	2.00	2.00
clock source latency	-4.12	-2.12
PLL/OUT (DUMMYPLL)	0.00	-2.12 f
clktree_2x/I (bufbd1)	0.00	-2.12 f
clktree_2x/Z (bufbd1)	1.50 *	-0.62 f
clktree_2xinv/I (inv0d2)	0.00	-0.62 f
clktree_2xinv/ZN (inv0d2)	0.04	-0.58 r
d2xinv_reg/CP (dfnrb1)	0.00	-0.58 r

library setup time	-0.08	-0.67
data required time		-0.67
-----		
data required time		-0.67
data arrival time		2.28
-----		
slack (MET)		1.61

It reports paths on both pll clocks (no need to change nworst here – there are two capture clocks instead of two launch clocks) – the worst being the capture with pllout\_minhigh because it is the nearest falling edge after launching data on the rising edge.

### 3.4 When to use these techniques

It is important to point out that you only need to create these pairs of clocks to represent duty cycle specs when you have opposite-edge clocking. If you're *sure* nothing happens on falling edges, you don't need to do all this.



## 4 Jitter

### 4.1 Jitter, skew, and uncertainty

I don't know if there are any "official" definitions of jitter and skew. I tend to think of jitter as a high-frequency, cycle-to-cycle phenomenon. I tend to think of skew as something static or at least very slow-changing. But there is one type of "skew" that I want to discuss up front – clock tree skew.

Clock tree skew is the difference in delays from the clock source to the various flop clock pins on the clock network. Because clock tree skew is often modeled in synthesis as clock uncertainty, it has come to be associated in many people's minds with all the other things that are modeled using clock uncertainty. However, for post-route analysis (which is what is being addressed in this paper), clock tree skew doesn't exist as a separate entity and does not need to be modeled. It is already there in the delay numbers or parasitics. For post-route analysis, any skew in the clock tree will automatically be handled by PT – it knows the exact delay to each element on the clock tree and will do slack calculations using these numbers. There is no need to budget for clock tree skew.

There is also a phenomenon that I would call "clock tree jitter". Since switching thresholds and delay can vary slightly according to switching activity in other gates, there may indeed be more cycle-to-cycle jitter at the end of the clock tree than there was at the beginning. I have seen this in an actual chip – a heavily loaded divide-by 2 clock was "modulating" the faster clock. Every time the slow clock switched from low to high, a lot of current was drawn, and the faster clock's edge rate would slow down, thus delaying the faster clock. This caused a fairly pronounced cycle-to-cycle jitter at the end of the fast clock tree. This is admittedly a somewhat extreme case, but we may someday have to account for "clock tree jitter" in our analysis.

### 4.2 My definition of jitter

There doesn't seem to be a common interpretation of what a jitter specification means. I'm going to define jitter somewhat loosely as follows:

"Jitter is the maximum/minimum variation in the length of a single clock *cycle*".

This means that a 10ns clock with jitter of +/-100ps can have a minimum cycle length of 9.9ns and a maximum cycle length of 10.1ns.

This is not the only possible definition of jitter. You can also define jitter as "the uncertainty in the location of a clock *edge* relative to its nominal location". By this definition, our 10ns clock with +/-100ps jitter can actually have a minimum period of 9.8 and a maximum period of 10.2 (jitter edge late followed by jitter edge early and jitter edge early followed by jitter edge late). Or you could define our original clock (period 9.9 to 10.1) as having *edge* jitter of +/-50ps.

These two different definitions stem from two different sources of jitter. The “cycle” definition represents “frequency” jitter – like a PLL adjusting its cycle to track a source. The “edge” definition represents “noise” jitter – switching thresholds and the like that cause uncertainty against a steady frequency.

Both types of jitter may be real. In most cases, you can convert “noise” jitter to an equivalent 2x value and add it to the “frequency” jitter to get *cycle* jitter (which is what I’m using here). Whenever you divide the clock in some way, this isn’t so clean. More on this later (see the section on jitter and generated clocks later in this chapter).

### 4.3 Sources of jitter

There are several sources of jitter related to the pll:. They are:

1. PLL cycle-to-cycle jitter. This is the change in period of the PLL on a cycle-to-cycle basis.
2. Refclk cycle-to-cycle jitter. This is the change in period of the Refclk on a cycle-to-cycle basis.
3. PLL phase error. The job of the pll is to make the Fb pin match the Ref pin in phase and frequency. But the pll is not perfect. Depending on process, temperature, etc there may be some small difference between the arrival time at the Ref pin and at the Fb pin. PLL phase error is the measure of this difference.

PLL phase error isn’t technically a form a jitter, but I have included it because it is modeled in the same way.

### 4.4 Effects of jitter on different sorts of paths

When modeling jitter effects in PT, it is important to consider the effect of different types of jitter on different paths.

First, consider the simple case of two flops on a single clock:

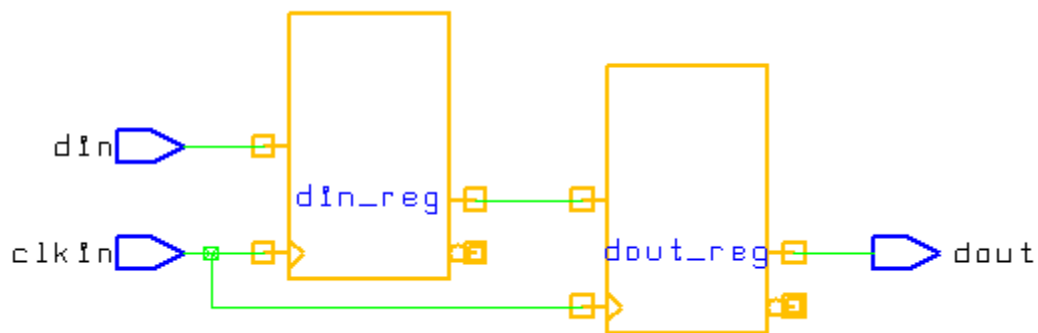
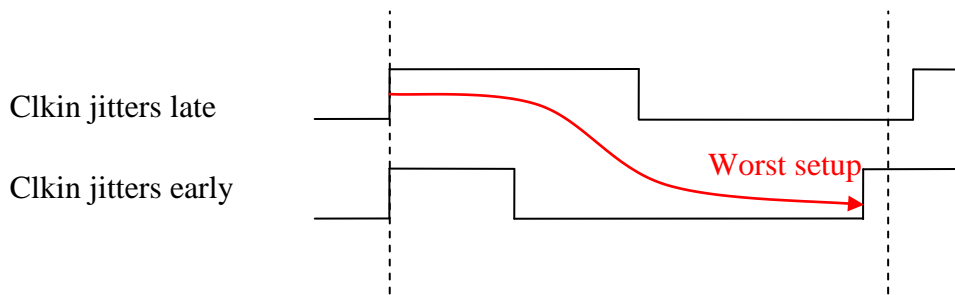


Figure 4-1

Once we select an arbitrary time zero for the first edge, here are two possibilities for the second clock edge – early and late.



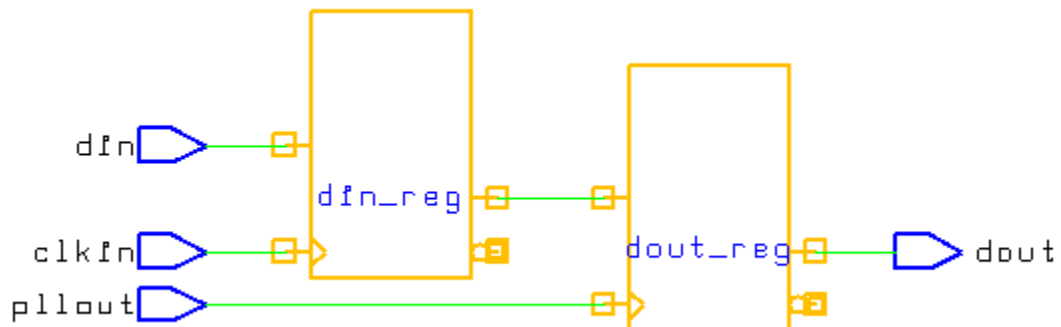
**Figure 4-2**

It's easy to see that the worst-case setup is affected by the jitter. If the clock jitters by +/- 100ps, that means that the capture clock could be 100ps earlier or 100ps later. Later won't matter (for setup), but earlier will reduce the available clock period and therefore reduce the path slack.

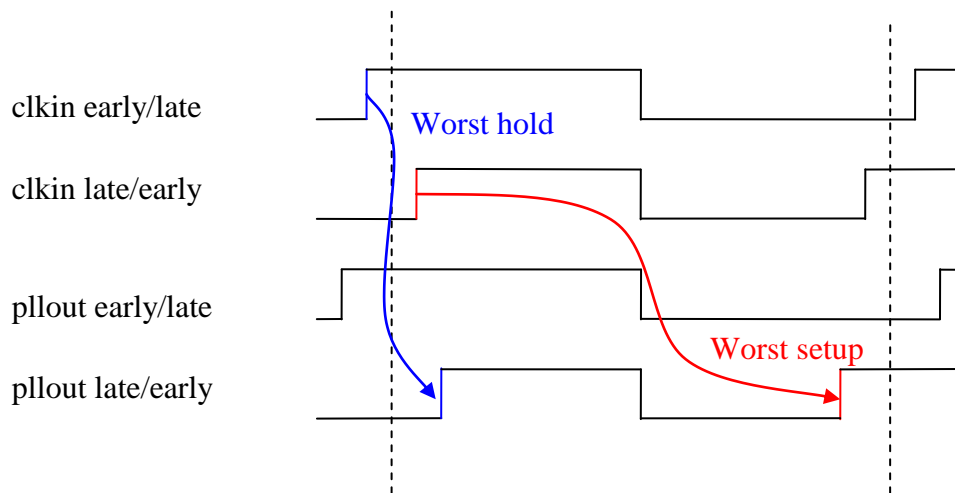
PLL cycle-to-cycle jitter will not, however, affect hold paths within the PLL's clock tree. This is because hold is a "single-edge" or "same-edge" phenomenon. The question is whether a source flop will change its data too soon to avoid being captured by the capture flop *on the same edge*. If the next edge jitters, it won't affect hold.

However, hold paths *between* the pll clock and another clock (like the refclk) *will be* affected by the PLL's cycle-to-cycle jitter.

Here's an example circuit and waveforms to illustrate this:



**Figure 4-3**

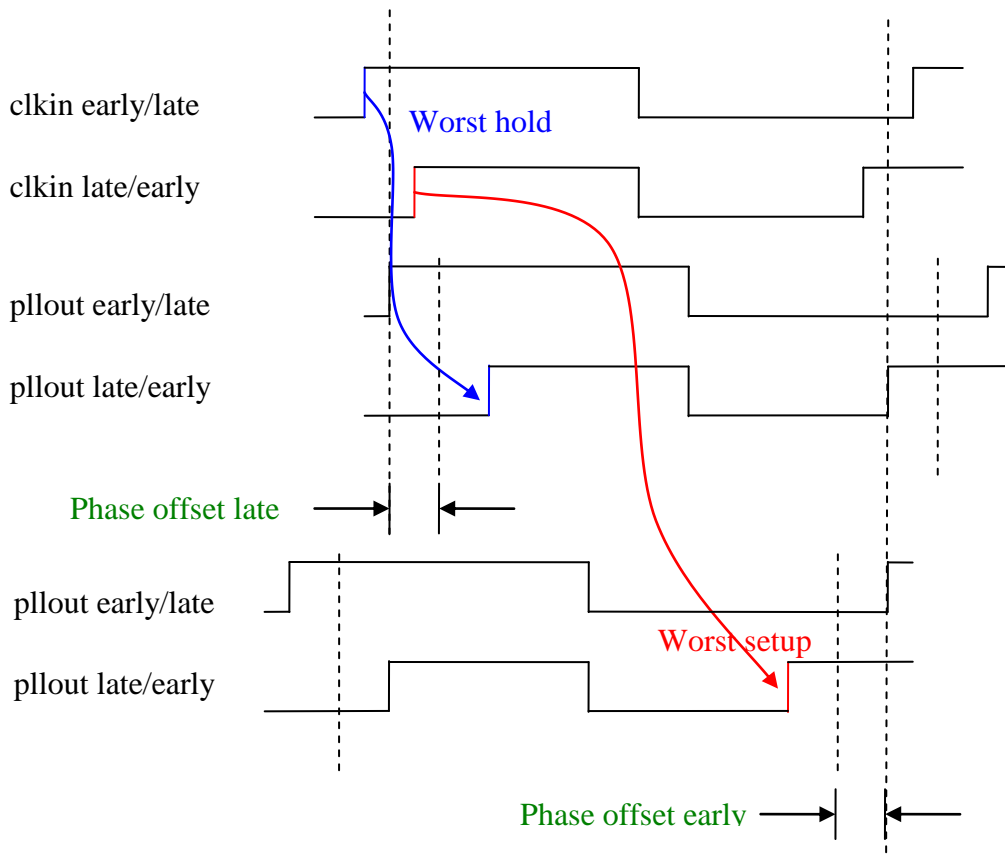


**Figure 4-4**

Although the path uses clocks that are nominally the “same edge” they are really not the *same* edge. They are really edges of two different clocks that are supposed to occur at the same time. But the “at the same time” is affected by jitter. A flop clocked by the pll could in fact launch its data earlier due to the PLL’s cycle-to-cycle jitter without there being any matching movement of the capture clock (refclk). It’s easy to see from the above waveform that both setup and hold margins will be reduced by the *sum* of the jitters of both clocks.

Now consider pll phase error. Unlike jitter, pll phase error doesn’t affect edge-to-edge timing at all. It is simply a time shift relative to some arbitrary external standard. Thus, it has absolutely no effect on internal paths within the PLL’s clock tree. If the pll clock is shifted from the refclk by 100ps or 100seconds, it won’t matter to flops along the PLL’s clock tree for either setup or hold.

But pll phase error will certainly affect paths between the pll clock and the ref clock, just as if it were pll cycle-to-cycle jitter:



Setup margin reduced by :  $clkin\_c2jitter(+)$  +  $phase\_offset(-)$  +  $pll\_c2jitter(-)$   
 Hold margin reduced by :  $clkin\_c2jitter(-)$  +  $phase\_offset(+)$  +  $pll\_c2jitter(+)$

**Figure 4-5**

Here is a table summarizing which type of jitter affects which paths:

Type of jitter	Same-clock setup	Same-clock hold	Refclk to/from pllclk setup	Refclk to/from pllclk hold
Pll cycle-to-cycle	Yes	No	Yes	Yes
Refclk cycle-to-cycle	Yes	No	Yes	Yes
Pll phase error	No	No	Yes	Yes

**Figure 4-6**

Note that I have used a very simplified definition of jitter in these examples. It might be worse than this. If you allow the refclock to have its maximum jitter for N cycles, and the pll clock to have its maximum jitter for M consecutive cycles, then you get setup/hold reductions of:  $N * \text{refclk\_max\_jitter} + M * \text{pllclk\_max\_jitter}$  (where “max\_jitter” means the max excursion from nominal – 100ps for a jitter of +/- 100ps). Ultimately this will have to be limited by the “long-term jitter” spec. If we define “long term jitter” as “the absolute maximum difference between the edges of the refclk and the pll clock under any and all circumstances”, then perhaps we should replace the “refclk\_c2c\_jitter + pllclk\_c2c\_jitter + phase\_error” with a single term – the long-term jitter. Or, since the long-term jitter spec was probably derived using a jitter-free reference clock, perhaps the correct answer is “refclk\_c2c\_jitter + pll\_long\_term\_jitter”.

Of course, *we’re* not free to define any of these terms. What matters is how the PLL designer defines them. Since there is, as far as I can tell, no clear industry-accepted definition of jitter specs, it is important to know what the person who created the spec meant.

The important point to remember is that the uncertainty *between* the ref clock and the pll clock is *different* from the uncertainty of the pll clock itself, and of the ref clock itself. You’ll have to work with the PLL designer to make sure you get the correct inter-clock value.

## 4.5 Modeling jitter with set\_clock\_uncertainty

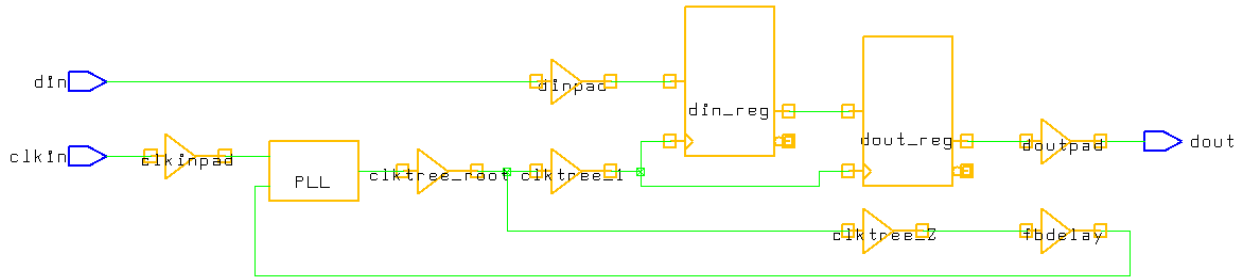
The PT command “set\_clock\_uncertainty” is really like two commands rolled into one. It can really be used only in one of two modes – setting uncertainty *within* a clock and setting uncertainty *between* clocks. Many of the options, in particular the edge-related options, only apply in inter-clock mode. Different values can be assigned for setup and hold uncertainty.

Looking at the tables above, it is pretty easy to figure out what values to use for the various uncertainties. For intra-clock (same clock) uncertainties, setup will use the cycle-to-cycle jitter spec for that clock (the “+/-“ value, or half of the total range). Hold uncertainty will use zero.

For inter-clock uncertainties between the refclk and the pll clock, setup will use the sum of all three jitter sources, as will hold.

## 4.6 Applying jitter specs to the example circuit – simple case

First we’ll apply this to the simple case – no divider in the feedback path and no duty cycle checks. This circuit is just like the one in section 2.2, except that I have added another flop in the data path to show internal path timing checks.



**Figure 4-7**

Here are the jitters specs we're going to use:

```
# Jitter specs
set _pll_c2cjitter 0.200 ;# +/- 0.200
set _phase_error 0.150 ;# +/- 0.150
set _refclk_c2cjitter 0.120 ;# +/- 0.120
```

After doing all the clock creation and setting the correct source latency as before, we can set the jitter. First we'll do the intra-clock uncertainty. Although the hold value would default to 0 anyway, I will set it here to make it clear that 0 is what I want:

```
# refclk internal
set_clock_uncertainty \
  -setup $_refclk_c2cjitter \
  [get_clocks clkIn]
set_clock_uncertainty \
  -hold 0 \
  [get_clocks clkIn]

# pllout internal
set_clock_uncertainty \
  -setup $_pll_c2cjitter \
  [get_clocks pllout]
set_clock_uncertainty \
  -hold 0 \
  [get_clocks pllout]
```

Now, we'll set the uncertainty between the two clocks:

```
# refclk to pllout
set_clock_uncertainty \
  -setup [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \
  -from [get_clocks clkkin] \
  -to [get_clocks pllout]
set_clock_uncertainty \
  -hold [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \
  -from [get_clocks clkkin] \
  -to [get_clocks pllout]

# pllout to refclk
set_clock_uncertainty \
  -setup [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \
  -from [get_clocks pllout] \
  -to [get_clocks clkkin]
set_clock_uncertainty \
  -hold [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \
  -from [get_clocks pllout] \
  -to [get_clocks clkkin]
```

Set the same i/o constraints as before:

```
set_input_delay -max 8.0 -clock clkkin [get_ports din]
set_input_delay -min 0.5 -clock clkkin [get_ports din]
set_output_delay -max 2.0 -clock clkkin [get_ports dout]
set_output_delay -min [expr -1.0 * 0.5] -clock clkkin [get_ports dout]
```

Now, let's run some reports.

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : jitter_simple
Version: V-2004.06
*****
```



Startpoint: din (input port clocked by clkin)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
din_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.30	7.70
PLL/OUT (DUMMYPLL)	0.00	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
din_reg/CP (dfnrb1)	0.00	10.20 r
inter-clock uncertainty	-0.47	9.73
library setup time	-0.08	9.65
data required time		9.65
-----		
data required time		9.65
data arrival time		-9.20
-----		
slack (MET)		0.45

Notice that this report is just the same as the report in section 2.2, except the slack is reduced by the inter-clock latency between clkin and the pll output clock. This latency value is 0.2 (the pll cycle-to-cycle jitter) plus 0.15 (the potential phase error between clkin and pllout) plus 0.120 (clkin's own cycle-to-cycle jitter).

The hold path also shows this slack reduction due to the inter-clock latency between the clocks:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din -
delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : jitter_simple
Version: V-2004.06-SP1
*****
```

Startpoint: din (input port clocked by clkin)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	0.50	0.50 f
din (in)	0.00	0.50 f
dinpad/I (bufbd1)	0.00	0.50 f
dinpad/Z (bufbd1)	1.20 *	1.70 f
din_reg/D (dfnrb1)	0.00	1.70 f
data arrival time		1.70
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
inter-clock uncertainty	0.47	0.67
library hold time	0.01	0.68
data required time		0.68
-----		
data required time		0.68
data arrival time		-1.70
-----		
slack (MET)		1.02

The dout path is also the same as in section 2.2, except for the inter-clock latency.

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -to dout
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : jitter_simple
Version: V-2004.06
*****
```

Startpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout (output port clocked by clkin)  
 Path Group: clkin  
 Path Type: max

Point	Incr	Path
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
dout_reg/CP (dfnrb1)	0.00	0.20 r
dout_reg/Q (dfnrb1)	0.33	0.53 f
doutpad/I (bufbd1)	0.00	0.53 f
doutpad/Z (bufbd1)	2.50 *	3.03 f
dout (out)	0.00	3.03 f
data arrival time		3.03
-----		
clock clkin (rise edge)	10.00	10.00
clock network delay (propagated)	0.00	10.00
inter-clock uncertainty	-0.47	9.53
output external delay	-2.00	7.53
data required time		7.53
-----		
data required time		7.53
data arrival time		-3.03
-----		
slack (MET)		4.50

Here's the path between the data flops:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : jitter_simple
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/Q (dfnrb1) <-	0.33	0.53 r
dout_reg/D (dfnrb1)	0.00	0.53 r
data arrival time		0.53
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.30	7.70
PLL/OUT (DUMMYPLL)	0.00	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
dout_reg/CP (dfnrb1)	0.00	10.20 r
clock uncertainty	-0.20	10.00
library setup time	-0.08	9.92
data required time		9.92
-----		
data required time		9.92
data arrival time		-0.53
-----		
slack (MET)		9.39

Notice the clock uncertainty of 0.20. This is the PLL's cycle-to-cycle jitter. The same path on hold has no uncertainty:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg -delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : jitter_simple
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/Q (dfnrb1) <-	0.33	0.53 f
dout_reg/D (dfnrb1)	0.00	0.53 f
data arrival time		0.53
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
dout_reg/CP (dfnrb1)	0.00	0.20 r
library hold time	0.01	0.21
data required time		0.21
-----		
data required time		0.21
data arrival time		-0.53
-----		
slack (MET)		0.32

As discussed earlier, the hold path within a clock network is not affected by jitter.

## 4.7 Generated clocks

What about generated clocks? For example, if the circuit contains a divide-by clock derived from the pll output clock, which types of jitter affect which paths on this clock?

Just as with the refclk/pllclk interface, we are faced with the question of what cycle-to-cycle jitter means on subsequent cycles.

Remember that my (cycle) definition of jitter was:

“Jitter is the maximum/minimum variation in the length of a single clock *cycle*”.

If you assume that subsequent cycles can also vary by this same amount, at least over the number of periods concerned, a divide-by N clock can end up with N times the original source clock jitter. Let's take the min example (10ns period, +/-100ps jitter):

1x clock	Div2 clock	Div4 clock
0	0	0
9.9		
19.8	19.8	
29.7		
39.6	39.6	39.6
49.5		
59.4	59.4	
69.3		
79.2	79.2	79.2

Jitter is 2x master clock (200ps)

Jitter is 4x master clock (400ps)

If, however, you were to take an edge-centric view of jitter, the uncertainty of any divided clocks would be the same as the uncertainty of the master clock.

Again, we get back to what the terms mean, and, again, you'll have to work this out with the PLL designer. Now let's look at the various inter- and intra-clock paths related to the generated clock.

We'll start with paths between the master clock (pll out clock in this case) and the generated clock. Assuming the path is single-cycle, pll cycle-to-cycle jitter will have no effect on hold for the same reason it has no effect on hold within the pll clock network – it's the same edge. Now consider the setup case (again, assuming a single-cycle path). If the generated clock is the capture clock, then the data will have been launched one master clock earlier. This launch time could be late by `pll_c2c_jitter`, reducing the setup margin by that amount. You cannot simultaneously pull the capture (generated) clock *in* because it is directly created by the master clock. Similarly, if the generated clock is the launch clock, then the next master clock is the capture clock, and it could be early by `pll_c2c_jitter`. So, the setup uncertainty for paths between the master and generated clocks is `pll_c2c_jitter`.

PLL phase error won't have any effect on these paths, nor will `refclk` cycle-to-cycle jitter.

Paths between the `refclk` and the pll generated clock will be affected in the same way as paths between the `refclk` and the pll output clock itself (affected by everything).

So, we can continue our table like this:

Type of jitter	Genclk internal setup	Genclk internal hold	Pllclk to/from pllgenclk setup	Pllclk to/from pllgenclk hold	Refclk to/from pllgenclk setup	Refclk to/from pllgenclk hold
Pll cycle-to-cycle	Yes (may be multiplied)	No	Yes	No	Yes	Yes
Refclk cycle-to-cycle	No	No	No	No	Yes	Yes
Pll phase error	No	No	No	No	Yes	Yes

Figure 4-8

## 4.8 What about falling edges?

What does a pll cycle-to-cycle jitter spec of “+/- 100ps” imply about the pll clock’s *falling* edge? The PLL’s clock is typically generated from a voltage-controlled oscillator converted to digital. If the cycle-to-cycle jitter reflects mostly this analog-to-digital conversion, then it would be fair to say that the entire cycle-to-cycle jitter would apply to the falling edge, too. This is the *edge* jitter definition. If however, the jitter is primarily due changes in the vco’s period as it attempts to track the refclk frequency, then it would seem that the falling edge should have about half the rated jitter. This is the *cycle* jitter definition.

If the answer is anything other than 100%, you cannot model it within a clock network with the set\_clock\_uncertainty command in PT. PT will *only* allow edge-specific uncertainty *between* clocks. Any intra-clock uncertainty will be applied in full force to both edges.

<added 3/17/2005>

As mentioned in the earlier update sidebar, it appears this may be possible using inter-clock syntax to describe an intra-clock uncertainty:

```
pt shell> set clock uncertainty -setup -rise from [get clocks clk_in] -fall to [get clocks clk_in] 0.2
```

1

So, you can probably use this to model asymmetrical jitter.

If you actually had this level of detail about the PLL’s behavior, and you needed to model the falling edge uncertainty as a different value from the rising edge, you would probably have to do

this by adjusting the duty cycle clocks discussed earlier. I've never had this level of detail to work with.

Opposite-edge clocking between the refclk and the pll clock is something I haven't run into. If an interface is slow enough to use opposite-edge clocking, it probably doesn't have an idc pll.

## 4.9 Applying jitter specs to the example circuit – complex case

*Author's note: If you're getting a little tired by this point, you might want to skip this section and go back to it later. It gets a little hairy...*

Ok, now let's try something a little harder. We'll go back to the 2x multiplier configuration, only this time I will include some cross-clock paths. And we'll handle the duty cycle as well.

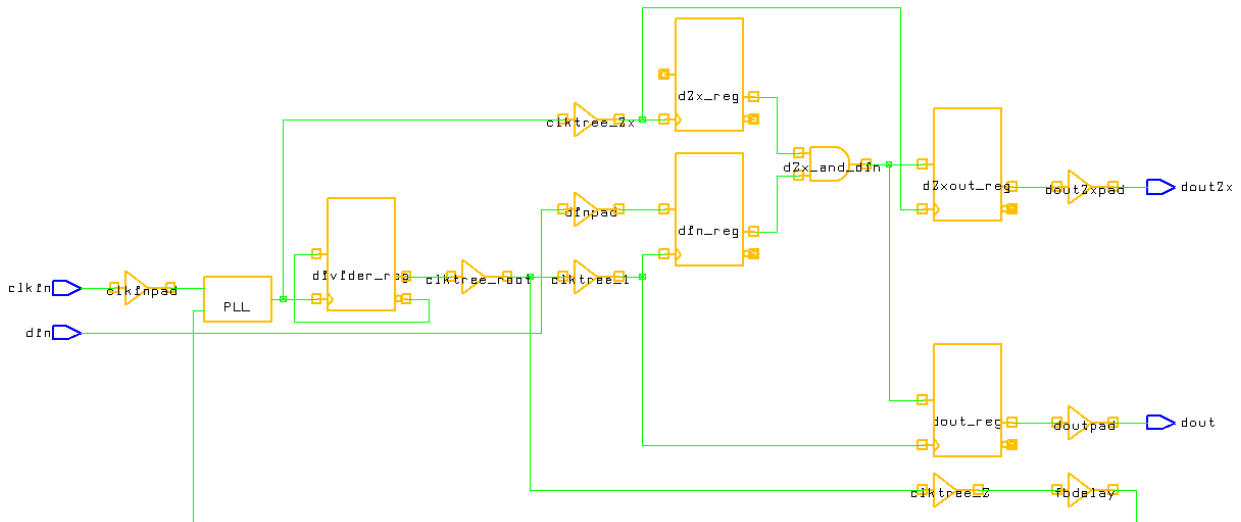


Figure 4-9

Note: I'm going to assume that the `divclk` uncertainty is 2 times the `pll_c2c_jitter` (see discussion above).

After setting up the clocks (2 primary clocks at `clkIn`, 2 primary clocks at PLL/OUT, and 1 generated clock at `divider_reg/Q`), and calculating and applying the pll source latency, we're ready to apply the uncertainties. We have a total of five clocks, but since the duty cycle clocks are exclusive to one another, we can treat them as three groups. We'll set up some collections to make this easier:

```
set_pll_clks [get_clocks pllout*high]
set_refclks [get_clocks clkIn*high]
```



First, we'll set the uncertainty for refclk (clkin) internal paths:

```
set_clock_uncertainty \  
  -setup $_refclk_c2cjitter \  
  $_refclks  
set_clock_uncertainty \  
  -hold 0 \  
  $_refclks
```

Now we'll do the paths from refclk to the pll clocks and divclk. They have the same value, so we can do it in one pair of statements.

```
set_clock_uncertainty \  
  -setup [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \  
  -from $_refclks \  
  -to [list $_pll_clks [get_clocks divclk] ]  
set_clock_uncertainty \  
  -hold [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \  
  -from $_refclks \  
  -to [list $_pll_clks [get_clocks divclk] ]
```

Now the pll output clk(s) internal paths:

```
set_clock_uncertainty \  
  -setup $_pll_c2cjitter \  
  $_pll_clks  
set_clock_uncertainty \  
  -hold 0 \  
  $_pll_clks
```

From pll output clock to divclk is special. It is similar to pll output clock's internal paths:

```
set_clock_uncertainty \  
  -setup $_pll_c2cjitter \  
  -from $_pll_clks \  
  -to [get_clocks divclk]  
set_clock_uncertainty \  
  -hold 0 \  
  -from $_pll_clks \  
  -to [get_clocks divclk]
```

Now pll output clocks to refclk:

```
set_clock_uncertainty \  
  -setup [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \  
  -from $_pll_clks \  
  -to $_refclks  
set_clock_uncertainty \  
  -hold [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \  
  -from $_pll_clks \  
  -to $_refclks
```

And finally we can do divclk itself.

First the internal paths. Here's where we apply the 2x multiplier.

```
set_clock_uncertainty \  
  -setup [expr 2 * $_pll_c2cjitter] \  
  [get_clocks divclk] \  
set_clock_uncertainty \  
  -hold 0 \  
  [get_clocks divclk]
```

Then the paths to pll output clock (again, similar to internal clocks):

```
set_clock_uncertainty \  
  -setup $_pll_c2cjitter \  
  -from [get_clocks divclk] \  
  -to $_pll_clks \  
set_clock_uncertainty \  
  -hold 0 \  
  -from [get_clocks divclk] \  
  -to $_pll_clks
```

And finally from divclk to refclk:

```
set_clock_uncertainty \  
  -setup [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \  
  -from [get_clocks divclk] \  
  -to $_refclks \  
set_clock_uncertainty \  
  -hold [expr $_refclk_c2cjitter + $_phase_error + $_pll_c2cjitter] \  
  -from [get_clocks divclk] \  
  -to $_refclks
```

Ok, now we set the i/o constraints as before:

```
foreach _dc {minhigh maxhigh} {  
  set_input_delay -max 4.0 -clock clk_in_${_dc} -add [get_ports din]  
  set_input_delay -min 0.5 -clock clk_in_${_dc} -add [get_ports din]  
  set_output_delay -max 1.0 -clock clk_in_${_dc} -add [get_ports dout]  
  set_output_delay -min [expr -1.0 * 0.5] -clock clk_in_${_dc} -add [get_ports  
dout]  
}
```

Here's the input report:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din  
*****  
Report : timing  
  -path full_clock_expanded  
  -delay max  
  -input_pins  
  -max_paths 1  
Design : jitter_example  
Version: V-2004.06  
*****
```

Startpoint: din (input port clocked by clkin\_minhigh)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by divclk)  
 Path Group: divclk  
 Path Type: max

Point	Incr	Path
-----		
clock clkin_minhigh (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
din_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
clock divclk (rise edge)	10.00	10.00
clock pllout_maxhigh (source latency)		
	-2.61	7.39
PLL/OUT (DUMMYPLL)	0.00	7.39 r
divider_reg/CP (dfnrb1)	0.00	7.39 r
divider_reg/Q (dfnrb1) (gclock source)		
	0.31	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
din_reg/CP (dfnrb1)	0.00	10.20 r
inter-clock uncertainty	-0.47	9.73
library setup time	-0.08	9.65
data required time		9.65
-----		
data required time		9.65
data arrival time		-9.20
-----		
slack (MET)		0.45

We expect the input and output path slacks to be the same as they were in the last section. Although now being captured by divclk, the pll will drive divclk with the same timing as the pllout clock was driven with in the previous example. The source latency and clock path delays will be different, but the net result will be the same. And it is.

The dout path also matches the previous example.

How about those cross-clock path?. We'll look at din\_reg to d2xout\_reg first.

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to d2xout_reg -delay max
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : jitter_example
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by divclk)  
 Endpoint: d2xout\_reg (rising edge-triggered flip-flop clocked by  
 pllout\_maxhigh)  
 Path Group: pllout\_maxhigh  
 Path Type: max

Point	Incr	Path
clock divclk (rise edge)	0.00	0.00
clock pllout_maxhigh (source latency)		
	-2.61	-2.61
PLL/OUT (DUMMYPLL)	0.00	-2.61 r
divider_reg/CP (dfnrb1)	0.00	-2.61 r
divider_reg/Q (dfnrb1) (gclock source)		
	0.31	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/Q (dfnrb1) <-	0.33	0.53 r
d2x_and_din/A2 (an02d2)	0.00	0.53 r
d2x_and_din/Z (an02d2)	0.13	0.66 r
d2xout_reg/D (dfnrb1)	0.00	0.66 r
data arrival time		0.66
clock pllout_maxhigh (rise edge)	5.00	5.00
clock source latency	-2.61	2.39
PLL/OUT (DUMMYPLL)	0.00	2.39 r
clktree_2x/I (bufbd1)	0.00	2.39 r
clktree_2x/Z (bufbd1)	1.50 *	3.89 r
d2xout_reg/CP (dfnrb1)	0.00	3.89 r
inter-clock uncertainty	-0.20	3.69
library setup time	-0.08	3.61
data required time		3.61
data required time		3.61
data arrival time		-0.66
slack (MET)		2.96

. . . (duplicate trace for pllout\_minhigh deleted)

Data is launched by divclk at time zero. The earliest that it can be captured by pllout is 5.0. Since the pllout clock has a cycle-to-cycle jitter of 0.200, the slack is reduced by this much via the inter-clock uncertainty value.

The same path for hold looks like this:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to d2xout_reg -delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : jitter_example
Version: V-2004.06
*****
```

```
Startpoint: din_reg (rising edge-triggered flip-flop clocked by divclk)
Endpoint: d2xout_reg (rising edge-triggered flip-flop clocked by
pllout_maxhigh)
Path Group: pllout_maxhigh
Path Type: min
```

Point	Incr	Path
clock divclk (rise edge)	0.00	0.00
clock pllout_maxhigh (source latency)		
	-2.61	-2.61
PLL/OUT (DUMMYPLL)	0.00	-2.61 r
divider_reg/CP (dfnrb1)	0.00	-2.61 r
divider_reg/Q (dfnrb1) (gclock source)		
	0.31	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/Q (dfnrb1) <-	0.33	0.53 f
d2x_and_din/A2 (an02d2)	0.00	0.53 f
d2x_and_din/Z (an02d2)	0.12	0.66 f
d2xout_reg/D (dfnrb1)	0.00	0.66 f
data arrival time		0.66
clock pllout_maxhigh (rise edge)	0.00	0.00
clock source latency	-2.61	-2.61
PLL/OUT (DUMMYPLL)	0.00	-2.61 r
clktree_2x/I (bufbd1)	0.00	-2.61 r
clktree_2x/Z (bufbd1)	1.50 *	-1.11 r
d2xout_reg/CP (dfnrb1)	0.00	-1.11 r
library hold time	0.01	-1.10
data required time		-1.10
data required time		-1.10
data arrival time		-0.66
slack (MET)		1.75

. . . (duplicate trace for pllout\_minhigh deleted)

Notice again the *lack* of an uncertainty value. If the cycle-to-cycle jitter caused the pllout clock edge to move, the divclk edge moved with it.

The other cross-clock paths are similar.

## 5 On-chip Variation

On-chip variation (OCV) refers to the fact that a physical device fabbed at a particular process point (best-case, worst-case, or somewhere in-between) will have some amount of variation in delays. Reference [5] has an excellent discussion of on-chip variation (OCV) and how it is handled by PT (and by Einstimer, if you're interested). I'm not going to go into all the details here, but I will present a simple example as background, and then show how this impacts pll timing.

### 5.1 The classic OCV case

Consider this simple circuit:

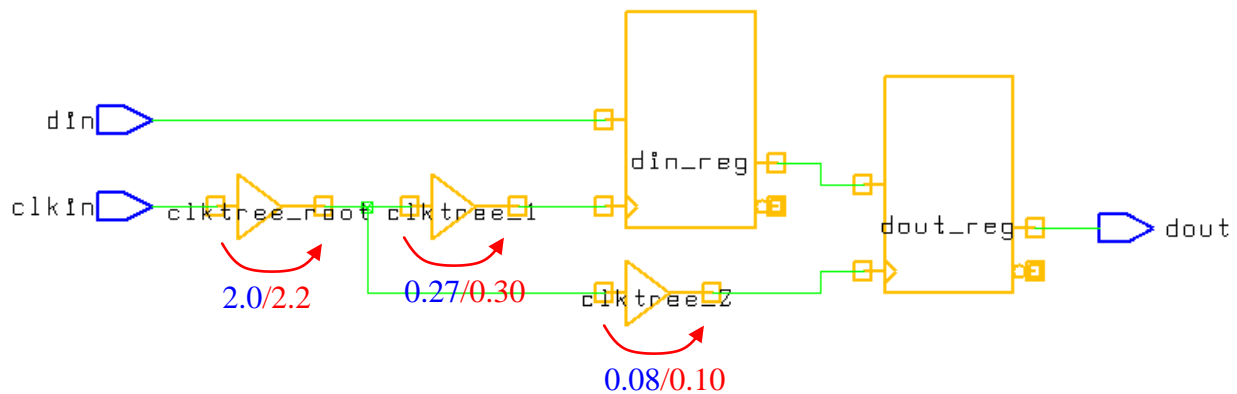


Figure 5-1

I have annotated the **min** and **max** delays onto the buffers (I left the CP->Q paths as constants to simplify the example).

Without OCV, the timing will use the worst-case (red) numbers and looks like this:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg
Information: Using automatic max wire load selection group 'predcaps'. (ENV-
003)
Information: Using automatic min wire load selection group 'predcaps'. (ENV-
003)
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_basic
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by clkin)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by clkin)  
 Path Group: clkin  
 Path Type: max

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clkin (in)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_1/I (bufbd1)	0.00	2.20 r
clktree_1/Z (bufbd1)	0.30 *	2.50 r
din_reg/CP (dfnrb1)	0.00	2.50 r
din_reg/Q (dfnrb1) <-	0.32 *	2.82 r
dout_reg/D (dfnrb1)	0.00	2.82 r
data arrival time		2.82
clock clkin (rise edge)	10.00	10.00
clock source latency	0.00	10.00
clkin (in)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.20 *	12.20 r
clktree_2/I (bufbd1)	0.00	12.20 r
clktree_2/Z (bufbd1)	0.10 *	12.30 r
dout_reg/CP (dfnrb1)	0.00	12.30 r
library setup time	-0.08	12.22
data required time		12.22
-----		
data required time		12.22
data arrival time		-2.82
-----		
slack (MET)		9.40

Let's look at that trace in a little more detail. The timing calculation is like this:

$$\begin{aligned}
 \text{slack} &= \text{period} + \text{capture\_clock} \\
 &\quad - \text{data\_arrival} \\
 \\
 \text{slack} &= \text{period} + T_{\text{clktree\_root}} + T_{\text{clktree\_2}} - T_{\text{su}} \\
 &\quad - ( T_{\text{clktree\_root}} + T_{\text{clktree\_1}} + T_{\text{cp2q}}) \\
 \\
 &= 10.0 \quad + 2.20 \quad \quad \quad + 0.10 \quad \quad - 0.08 \\
 &\quad - ( \quad \quad 2.20 \quad \quad \quad + 0.30 \quad \quad + 0.32) \\
 \\
 &= 9.40
 \end{aligned}$$

Now let's turn on OCV analysis. In this example, I'm using an SDF file with the worst case values in the 3<sup>rd</sup> part of the triplet, and the not-quite-worst-case values in the 1<sup>st</sup> part of the triplet. So, I turn on OCV with the following command:

```

read_sdf \
  -analysis_type on_chip_variation \
  -min_type sdf_min \
  -max_type sdf_max \
  ocv.sdf

```



If I were using parasitics, I might use something like this:

```
set_operating_conditions \  
-analysis_type on_chip_variation \  
-min MIN -max MAX
```

What happens when we turn on OCV? Well, PT will calculate the worst case. The worst case is where the terms in the first line are at their minimums and those on the second line are at their maximums:

$$\begin{aligned} &= 10.0 + 2.00 + 0.08 - 0.08 \\ &- ( 2.20 + 0.30 + 0.32) \\ &= 9.18 \end{aligned}$$

And here's the trace:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from  
din_reg -to dout_reg  
*****  
Report : timing  
-path full_clock_expanded  
-delay max  
-input_pins  
-max_paths 1  
Design : ocv_basic  
Version: V-2004.06  
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by clkin)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by clkin)  
 Path Group: clkin  
 Path Type: max

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clkin (in)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_1/I (bufbd1)	0.00	2.20 r
clktree_1/Z (bufbd1)	0.30 *	2.50 r
din_reg/CP (dfnrb1)	0.00	2.50 r
din_reg/Q (dfnrb1) <-	0.32 *	2.82 r
dout_reg/D (dfnrb1)	0.00	2.82 r
data arrival time		2.82
clock clkin (rise edge)	10.00	10.00
clock source latency	0.00	10.00
clkin (in)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.00 *	12.00 r
clktree_2/I (bufbd1)	0.00	12.00 r
clktree_2/Z (bufbd1)	0.08 *	12.08 r
dout_reg/CP (dfnrb1)	0.00	12.08 r
library setup time	-0.08	12.00
data required time		12.00
-----		
data required time		12.00
data arrival time		-2.82
-----		
slack (MET)		9.18

## 5.2 Enter CRPR

But is that really correct? The value for Tclktree\_root shows up twice in the calculation. We have used its min value in one place, and its max value in another. This is not correct. OCV is not a cycle-to-cycle variation. The clktree\_root buffer can be an max or min (or somewhere in between), but it can't be at different values on subsequent cycles. Whether we use the min or max value, the slack will end up being 9.38.

This phenomenon is known as “common path pessimism” or “clock reconvergence pessimism”. PT can correct this error by examining the clock paths and identifying the common elements. Instead of using the same value for common elements, it does the simplistic (incorrect) calculation above, and then adds in a correction factor to the slack. The correction factor (called “clock reconvergence pessimism”) is the sum of all the differences between min and max for the common elements.

PT can do this, but you have to turn the feature on:

```
set timing_remove_clock_reconvergence_pessimism true
```

Also, because of the values I have chosen for delays and OCV, it is necessary to change the default value for `timing_crpr_threshold_ps`. This variable sets the value (in ps) below which PT will not bother to adjust for OCV.

```
set timing_crpr_threshold_ps 1
```

Now if we run the report, we get the correction factor and the correct slack:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_basic
Version: V-2004.06
*****
```

```
Startpoint: din_reg (rising edge-triggered flip-flop clocked by clkin)
Endpoint:   dout_reg (rising edge-triggered flip-flop clocked by clkin)
Path Group: clkin
Path Type:  max
```

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clkin (in)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_1/I (bufbd1)	0.00	2.20 r
clktree_1/Z (bufbd1)	0.30 *	2.50 r
din_reg/CP (dfnrb1)	0.00	2.50 r
din_reg/Q (dfnrb1) <-	0.32 *	2.82 r
dout_reg/D (dfnrb1)	0.00	2.82 r
data arrival time		2.82
-----		
clock clkin (rise edge)	10.00	10.00
clock source latency	0.00	10.00
clkin (in)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.00 *	12.00 r
clktree_2/I (bufbd1)	0.00	12.00 r
clktree_2/Z (bufbd1)	0.08 *	12.08 r
dout_reg/CP (dfnrb1)	0.00	12.08 r
clock reconvergence pessimism	0.20	12.28
library setup time	-0.08	12.20
data required time		12.20
-----		
data required time		12.20
data arrival time		-2.82
-----		
slack (MET)		9.38

The “clock reconvergence pessimism” reflects the difference between min and max values for the shared clock element `clktree_root` ( $2.2 - 2.0 = 0.2$ ).

### 5.3 OCV and PLLs

Now let’s drive this little circuit from a pll:

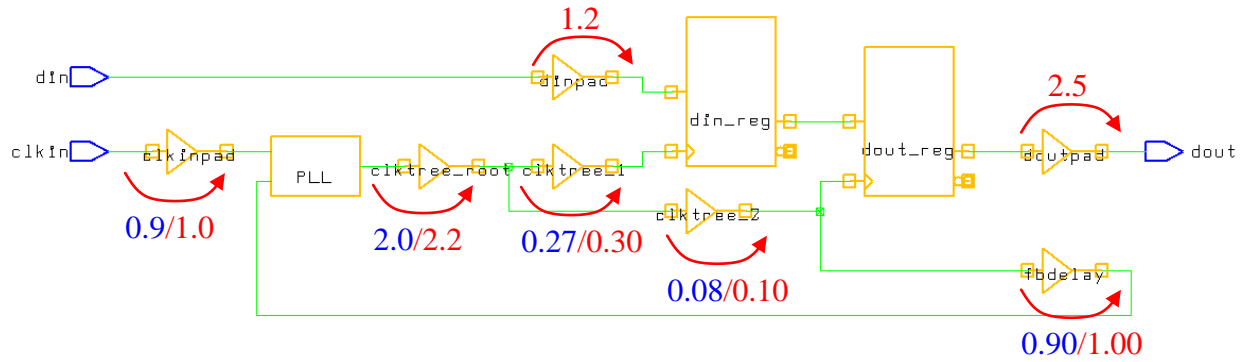


Figure 5-2

Once again, I have only applied OCV to the clock elements to simplify the example.

Looking at this, we are immediately faced with an issue: how do we assign the source latency to the pll clock? The delays through the elements in the feedback path have variable delays.

The most obvious approach is to calculate a min path and a max path, and apply min and max source latencies. The min value will be  $T_{ref\_min} - T_{fb\_max}$ ; the max value will be  $T_{ref\_max} - T_{fb\_min}$ .

```
set timing_remove_clock_reconvergence_pessimism true

create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout

set _path [get_timing_paths -delay max_rise \
  -from [get_ports clkin] \
  -to [get_pins PLL/CKREF] \
]
set _ref_delay_max [get_attribute $_path arrival]

set _path [get_timing_paths -delay min_rise \
  -from [get_ports clkin] \
  -to [get_pins PLL/CKREF] \
]
set _ref_delay_min [get_attribute $_path arrival]

set _path [get_timing_paths -delay max_rise \
  -from [get_pins PLL/OUT] \
  -to [get_pins PLL/FB] \
]
set _fb_delay_max [get_attribute $_path arrival]

set _path [get_timing_paths -delay min_rise \
  -from [get_pins PLL/OUT] \
  -to [get_pins PLL/FB] \
]
set _fb_delay_min [get_attribute $_path arrival]

set_clock_latency -early -source \
  [expr $_ref_delay_min - $_fb_delay_max] \
  [get_clocks pllout]

set_clock_latency -late -source \
  [expr $_ref_delay_max - $_fb_delay_min] \
  [get_clocks pllout]
```

Now look at the clock skew report:

```
pt_shell> report_clock -skew
*****
Report : clock_skew
Design : ocv_pll
Version: V-2004.06
*****
```

Object	Min Condition Source Latency				Max Condition Source Latency				Rel_clk
	Early_r	Early_f	Late_r	Late_f	Early_r	Early_f	Late_r	Late_f	
pllout	-2.40	-2.40	-1.98	-1.98	-2.40	-2.40	-1.98	-1.98	--

Note: set\_clock\_latency also has -min/-max options. If you use those instead, you get the following skew report:

```
pt_shell> report_clock -skew
*****
Report : clock_skew
Design : ocv_pll
Version: V-2004.06
*****
```

Object	Min Condition Source Latency				Max Condition Source Latency				Rel_clk
	Early_r	Early_f	Late_r	Late_f	Early_r	Early_f	Late_r	Late_f	
pllout	-2.40	-2.40	-2.40	-2.40	-1.98	-1.98	-1.98	-1.98	--

Since all the timing reports seem to use early/min against late/max, both sets of options have the same effect (in the reports that I have run). In all the timing traces shown below, it makes no difference whether you use early/late or min/max. I think that early/late is “more correct” in this case. I’d be interested to know if there are any timing reports where the distinction matters.

Anyway, the report has two values now.

But does this approach get the correct slack values?

Let's start with the flop-to-flop path. To avoid the chance of using different delays for the same gate in a single calculation, let's look at the two extreme cases (Tref@min/Tfb@max and Tref@max/Tfb@min) separately:

	Case 1: Trefclk@min, Tfb@max	Case2: Trefclk@max, Tfb@min
Tclkinpad	0.9	1.0
Tclktree_root	2.2	2.0
Tclktree_2	0.10	0.08
Tfbdelay	1.0	0.9
Source latency	-2.40	-1.98

Figure 5-3

### Flop-to-flop setup – Case 1:

$$\text{slack} = \text{period} + \text{capture\_clock} - \text{data\_arrival}$$

$$\begin{aligned} \text{slack} &= \text{period} + \text{Tsrclat} + \text{Tclktree\_root} + \text{Tclktree\_2} - \text{Tsu} \\ &- ( \text{Tsrclat} + \text{Tclktree\_root} + \text{Tclktree\_1(max)} + \text{Tcp2q} ) \\ &= 10.0 + -2.40 + 2.2 + 0.10 - 0.08 \\ &- ( + -2.40 + 2.2 + 0.30 + 0.32 ) \\ &= 9.4 \end{aligned}$$

### Flop-to-flop setup – Case 2:

$$\begin{aligned} \text{slack} &= \text{period} + \text{Tsrclat} + \text{Tclktree\_root} + \text{Tclktree\_2} - \text{Tsu} \\ &- ( \text{Tsrclat} + \text{Tclktree\_root} + \text{Tclktree\_1(max)} + \text{Tcp2q} ) \\ &= 10.0 + -1.98 + 2.0 + 0.08 - 0.08 \\ &- ( + -1.98 + 2.0 + 0.30 + 0.32 ) \\ &= 9.38 \end{aligned}$$

So, the minimum slack should be 9.38. What did PT get?

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-1.98	-1.98
PLL/OUT (DUMMYPLL)	0.00	-1.98 r
clktree_root/I (bufbd1)	0.00	-1.98 r
clktree_root/Z (bufbd1)	2.20 *	0.22 r
clktree_1/I (bufbd1)	0.00	0.22 r
clktree_1/Z (bufbd1)	0.30 *	0.52 r
din_reg/CP (dfnrb1)	0.00	0.52 r
din_reg/Q (dfnrb1) <-	0.32 *	0.84 r
dout_reg/D (dfnrb1)	0.00	0.84 r
data arrival time		0.84
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.40	7.60
PLL/OUT (DUMMYPLL)	0.00	7.60 r
clktree_root/I (bufbd1)	0.00	7.60 r
clktree_root/Z (bufbd1)	2.00 *	9.60 r
clktree_2/I (bufbd1)	0.00	9.60 r
clktree_2/Z (bufbd1)	0.08 *	9.68 r
dout_reg/CP (dfnrb1)	0.00	9.68 r
clock reconvergence pessimism	0.62	10.30
library setup time	-0.08	10.22
data required time		10.22
-----		
data required time		10.22
data arrival time		-0.84
-----		
slack (MET)		9.38

PT got the correct answer, but the trace looks a little strange. In the explicit calculations above, it is easy to see that the pll source latency will cancel. But PT used the max source latency value for the data launch, and the min value for the data capture (don't forget that  $-1.98 > -2.40$  !). Yet it still got the correct slack. How?

Take a look at the “clock reconvergence pessimism” value. The only common element in the clock paths of the two flops is clktree\_root. The min/max difference on this gate is 0.2. So why is the clock reconvergence pessimism 0.62? Because PT cleverly recognized that the pll source latency is *also* a common path. So it gave back an additional  $-1.98 - (-2.40) = 0.42$  of credit! It works!



The same happens on the hold calculation.

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg -delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

```
Startpoint: din_reg (rising edge-triggered flip-flop clocked by pllout)
Endpoint:   dout_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type:  min
```

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.40	-2.40
PLL/OUT (DUMMYPLL)	0.00	-2.40 r
clktree_root/I (bufbd1)	0.00	-2.40 r
clktree_root/Z (bufbd1)	2.00 *	-0.40 r
clktree_1/I (bufbd1)	0.00	-0.40 r
clktree_1/Z (bufbd1)	0.27 *	-0.13 r
din_reg/CP (dfnrb1)	0.00	-0.13 r
din_reg/Q (dfnrb1) <-	0.32 *	0.19 f
dout_reg/D (dfnrb1)	0.00	0.19 f
data arrival time		0.19
clock pllout (rise edge)	0.00	0.00
clock source latency	-1.98	-1.98
PLL/OUT (DUMMYPLL)	0.00	-1.98 r
clktree_root/I (bufbd1)	0.00	-1.98 r
clktree_root/Z (bufbd1)	2.20 *	0.22 r
clktree_2/I (bufbd1)	0.00	0.22 r
clktree_2/Z (bufbd1)	0.10 *	0.32 r
dout_reg/CP (dfnrb1)	0.00	0.32 r
clock reconvergence pessimism	-0.62	-0.30
library hold time	0.01	-0.29
data required time		-0.29
-----		
data required time		-0.29
data arrival time		-0.19
-----		
slack (MET)		0.48

## 5.4 The OCV/PLL excess pessimism problem

So much for the good news. Now let's look at the i/o timing. We'll use the same i/o constraints as before:

```
set_input_delay -max 8.0 -clock clkin [get_ports din]
set_input_delay -min 0.5 -clock clkin [get_ports din]
set_output_delay -max 2.0 -clock clkin [get_ports dout]
set_output_delay -min [expr -1.0 * 0.5] -clock clkin [get_ports dout]
```

Let's start with the setup check from input din:

Din setup – Case 1:

```
slack = period + capture_clock - Tsu
        - data_arrival

slack = period + Tsrclat + Tclktree_root + Tclktree_1(min) - Tsu
        - ( Tin    + Tinpad
          = 10.0   + -2.40   + 2.2           + 0.27           - 0.08
          - ( 8.0   + 1.2
          = 0.79
```

Din setup – Case 2:

```
slack = period + Tsrclat + Tclktree_root + Tclktree_1(min) - Tsu
        - ( Tin    + Tinpad
          = 10.0   + -1.98   + 2.0           + 0.27           - 0.08
          - ( 8.0   + 1.2
          = 1.01
```

So, the minimum slack should be 0.79.

What did PT get?

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din (input port clocked by clkin)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
din_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.40	7.60
PLL/OUT (DUMMYPLL)	0.00	7.60 r
clktree_root/I (bufbd1)	0.00	7.60 r
clktree_root/Z (bufbd1)	2.00 *	9.60 r
clktree_1/I (bufbd1)	0.00	9.60 r
clktree_1/Z (bufbd1)	0.27 *	9.87 r
din_reg/CP (dfnrb1)	0.00	9.87 r
clock reconvergence pessimism	0.00	9.87
library setup time	-0.08	9.79
data required time		9.79
-----		
data required time		9.79
data arrival time		-9.20
-----		
slack (MET)		0.59

Oops. What went wrong?

The problem here is an undetected clock reconvergence pessimism. PT used the min source latency (-2.40), which is based on the max Tfb, which uses Tclktree\_root=2.2. Then the path went through Tclktree\_root directly, and PT used 2.0. This is clock reconvergence pessimism, but PT doesn't know about it, so the clock reconvergence pessimism value is 0.00. This isn't really PT's fault. PT doesn't know about the pll, so it can't see the common path.

The hold is wrong, too:

Din hold – Case 1:

$$\begin{aligned}
 \text{slack} &= \text{early data} \\
 &\quad - (\text{late clock} + T_h) \\
 \\
 \text{slack} &= T_{\text{sid\_min}} + T_{\text{indpad}} \\
 &\quad - (T_{\text{src\bar{lat}}} + T_{\text{clktree\_root}} + T_{\text{clktree\_1(max)}} + T_h) \\
 &= 0.5 \quad + 1.2 \\
 &\quad - (-2.40 \quad + 2.2 \quad + 0.30 \quad + 0.01) \\
 &= 1.59
 \end{aligned}$$

## Din hold – Case 2:

$$\begin{aligned}
 \text{slack} &= \text{Tsid\_min} + \text{Tindpad} \\
 &- (\text{Tsrclat} + \text{Tclktree\_root} + \text{Tclktree\_1(max)} + \text{Th}) \\
 &= 0.5 + 1.2 \\
 &- (-1.98 + 2.0 + 0.30 + 0.01) \\
 &= 1.37
 \end{aligned}$$

So, the minimum slack is 1.37. But PT got:

```

pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din -
delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****

```

```

Startpoint: din (input port clocked by clkin)
Endpoint:   din_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type:  min

```

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	0.50	0.50 f
din (in)	0.00	0.50 f
dinpad/I (bufbd1)	0.00	0.50 f
dinpad/Z (bufbd1)	1.20 *	1.70 f
din_reg/D (dfnrb1)	0.00	1.70 f
data arrival time		1.70
clock pllout (rise edge)	0.00	0.00
clock source latency	-1.98	-1.98
PLL/OUT (DUMMYPLL)	0.00	-1.98 r
clktree_root/I (bufbd1)	0.00	-1.98 r
clktree_root/Z (bufbd1)	2.20 *	0.22 r
clktree_1/I (bufbd1)	0.00	0.22 r
clktree_1/Z (bufbd1)	0.30 *	0.52 r
din_reg/CP (dfnrb1)	0.00	0.52 r
clock reconvergence pessimism	0.00	0.52
library hold time	0.01	0.53
data required time		0.53
-----		
data required time		0.53
data arrival time		-1.70
-----		
slack (MET)		1.17

So, PT has *excess pessimism* due to the fact that it doesn't know about the pll adjustment explicitly.

## 6 OCV/PLL excess pessimism workarounds

### 6.1 Forcing OCV off on the fb path

Recognize that the io path calculation will *always* be wrong as long as the ios are referenced to the external clock and the feedback path component of the source latency has multiple values. Why? Because the effect of a delay in the feedback path is the opposite of what it is in the main path. A larger delay will produce an *earlier* source latency, but a longer clock path. So, if it doesn't know about the clock reconvergence, PT will either use early source latency / slow clock or late source latency / fast clock, both of which are wrong.

Therefore, one approach to working around this problem that has been suggested (and I believe is used by at least one vendor) is to turn off OCV on all the elements in the fb path .

Before showing the code to do this, I want to show the effect..

Here's our example again, with the values forced to max:

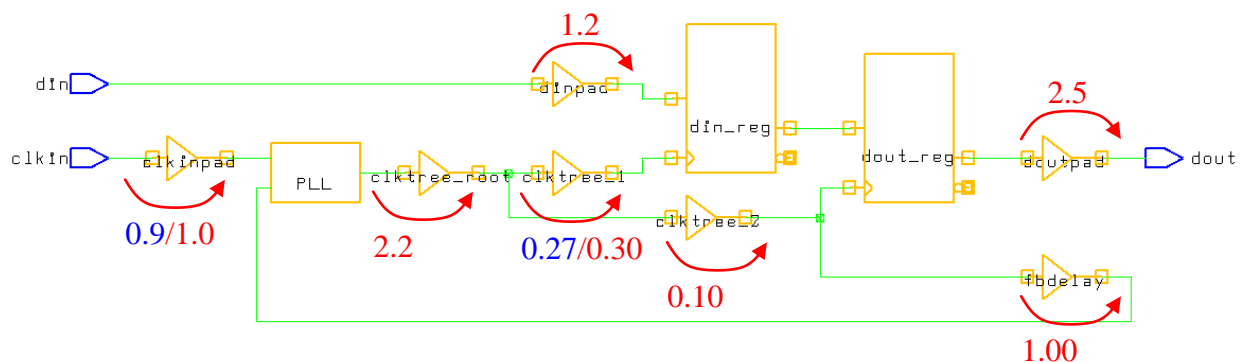


Figure 6-1

The OCV effect is retained on **clk\_inpad** (Tref), since it will never share gates with the clock path. Therefore, the code uses “set\_clock\_latency -early” and “set\_clock\_latency -late”, but only a single value for `$_fb_delay`.

```
set_clock_latency -early -source \  
  [expr $_ref_delay_min - $_fb_delay] \  
  [get_clocks pllout]  
  
set_clock_latency -late -source \  
  [expr $_ref_delay_max - $_fb_delay] \  
  [get_clocks pllout]
```

The early source latency value will be  $0.9 - (2.2 + 0.1 + 1.0) = -2.40$ .

The late source latency value will be  $1.0 - (2.2 + 0.1 + 1.0) = -2.30$ .

OCV is retained on clktree\_1 as well, since it is not part of the feedback path.

This approach works fine on the din setup path:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

```
Startpoint: din (input port clocked by clkin)
Endpoint:   din_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type:  max
```

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
din_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.40	7.60
PLL/OUT (DUMMYPLL)	0.00	7.60 r
clktree_root/I (bufbd1)	0.00	7.60 r
clktree_root/Z (bufbd1)	2.20 *	9.80 r
clktree_1/I (bufbd1)	0.00	9.80 r
clktree_1/Z (bufbd1)	0.27 *	10.07 r
din_reg/CP (dfnrb1)	0.00	10.07 r
clock reconvergence pessimism	0.00	10.07
library setup time	-0.08	9.99
data required time		9.99
-----		
data required time		9.99
data arrival time		-9.20
-----		
slack (MET)		0.79

But it doesn't work on the hold path:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din -
delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

```
Startpoint: din (input port clocked by clkin)
Endpoint: din_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type: min
```

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	0.50	0.50 f
din (in)	0.00	0.50 f
dinpad/I (bufbd1)	0.00	0.50 f
dinpad/Z (bufbd1)	1.20 *	1.70 f
din_reg/D (dfnrb1)	0.00	1.70 f
data arrival time		1.70
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
clock reconvergence pessimism	0.00	0.20
library hold time	0.01	0.21
data required time		0.21
-----		
data required time		0.21
data arrival time		-1.70
-----		
slack (MET)		1.49

The correct slack is 1.37. So this approach results in excess *optimism*.



If I instead force the feedback path values to all mins, the hold path works but the setup path is optimistic:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06-SP1
*****
```

```
Startpoint: din (input port clocked by clkin)
Endpoint:   din_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type:  max
```

Point	Incr	Path
clock clkin (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	8.00	8.00 r
din (in)	0.00	8.00 r
dinpad/I (bufbd1)	0.00	8.00 r
dinpad/Z (bufbd1)	1.20 *	9.20 r
din_reg/D (dfnrb1)	0.00	9.20 r
data arrival time		9.20
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.08	7.92
PLL/OUT (DUMMYPLL)	0.00	7.92 r
clktree_root/I (bufbd1)	0.00	7.92 r
clktree_root/Z (bufbd1)	2.00 *	9.92 r
clktree_1/I (bufbd1)	0.00	9.92 r
clktree_1/Z (bufbd1)	0.27 *	10.19 r
din_reg/CP (dfnrb1)	0.00	10.19 r
clock reconvergence pessimism	0.00	10.19
library setup time	-0.08	10.11
data required time		10.11
-----		
data required time		10.11
data arrival time		-9.20
-----		
slack (MET)		0.91

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din -
delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06-SP1
*****
```

Startpoint: din (input port clocked by clkln)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
-----		
clock clkln (rise edge)	0.00	0.00
clock network delay (propagated)	0.00	0.00
input external delay	0.50	0.50 f
din (in)	0.00	0.50 f
dinpad/I (bufbd1)	0.00	0.50 f
dinpad/Z (bufbd1)	1.20 *	1.70 f
din_reg/D (dfnrb1)	0.00	1.70 f
data arrival time		1.70
clock pllout (rise edge)	0.00	0.00
clock source latency	-1.98	-1.98
PLL/OUT (DUMMYPLL)	0.00	-1.98 r
clktree_root/I (bufbd1)	0.00	-1.98 r
clktree_root/Z (bufbd1)	2.00 *	0.02 r
clktree_1/I (bufbd1)	0.00	0.02 r
clktree_1/Z (bufbd1)	0.30 *	0.32 r
din_reg/CP (dfnrb1)	0.00	0.32 r
clock reconvergence pessimism	0.00	0.32
library hold time	0.01	0.33
data required time		0.33
-----		
data required time		0.33
data arrival time		-1.70
-----		
slack (MET)		1.37

The effect is similar for internal paths. Here's the all-max trace for setup:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_1/I (bufbd1)	0.00	-0.10 r
clktree_1/Z (bufbd1)	0.30 *	0.20 r
din_reg/CP (dfnrb1)	0.00	0.20 r
din_reg/Q (dfnrb1) <-	0.32 *	0.52 r
dout_reg/D (dfnrb1)	0.00	0.52 r
data arrival time		0.52
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.40	7.60
PLL/OUT (DUMMYPLL)	0.00	7.60 r
clktree_root/I (bufbd1)	0.00	7.60 r
clktree_root/Z (bufbd1)	2.20 *	9.80 r
clktree_2/I (bufbd1)	0.00	9.80 r
clktree_2/Z (bufbd1)	0.10 *	9.90 r
dout_reg/CP (dfnrb1)	0.00	9.90 r
clock reconvergence pessimism	0.10	10.00
library setup time	-0.08	9.92
data required time		9.92
-----		
data required time		9.92
data arrival time		-0.52
-----		
slack (MET)		9.40

The problem here is more subtle. Since this is an internal path, PT knows about the clock reconvergence pessimism due to the min/max source latency, and applies the correct adjustment. But notice that buffer clktree\_2 is in the feedback path *and* in the path leading to the capture flop. Because we have forced OCV off on this buffer, the very real slack reduction due to OCV has been removed. Again, the slack is optimistic.

But the all-max approach *is* correct on internal hold:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg -delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.40	-2.40
PLL/OUT (DUMMYPLL)	0.00	-2.40 r
clktree_root/I (bufbd1)	0.00	-2.40 r
clktree_root/Z (bufbd1)	2.20 *	-0.20 r
clktree_1/I (bufbd1)	0.00	-0.20 r
clktree_1/Z (bufbd1)	0.27 *	0.07 r
din_reg/CP (dfnrb1)	0.00	0.07 r
din_reg/Q (dfnrb1) <-	0.32 *	0.39 f
dout_reg/D (dfnrb1)	0.00	0.39 f
data arrival time		0.39
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.30	-2.30
PLL/OUT (DUMMYPLL)	0.00	-2.30 r
clktree_root/I (bufbd1)	0.00	-2.30 r
clktree_root/Z (bufbd1)	2.20 *	-0.10 r
clktree_2/I (bufbd1)	0.00	-0.10 r
clktree_2/Z (bufbd1)	0.10 *	-0.00 r
dout_reg/CP (dfnrb1)	0.00	0.00 r
clock reconvergence pessimism	-0.10	-0.10
library hold time	0.01	-0.09
data required time		-0.09
-----		
data required time		-0.09
data arrival time		-0.39
-----		
slack (MET)		0.48

As you might expect, the all-min approach gets the setup slack right, but the hold is optimistic:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg -delay max
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-1.98	-1.98
PLL/OUT (DUMMYPLL)	0.00	-1.98 r
clktree_root/I (bufbd1)	0.00	-1.98 r
clktree_root/Z (bufbd1)	2.00 *	0.02 r
clktree_1/I (bufbd1)	0.00	0.02 r
clktree_1/Z (bufbd1)	0.30 *	0.32 r
din_reg/CP (dfnrb1)	0.00	0.32 r
din_reg/Q (dfnrb1) <-	0.32 *	0.64 r
dout_reg/D (dfnrb1)	0.00	0.64 r
data arrival time		0.64
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.08	7.92
PLL/OUT (DUMMYPLL)	0.00	7.92 r
clktree_root/I (bufbd1)	0.00	7.92 r
clktree_root/Z (bufbd1)	2.00 *	9.92 r
clktree_2/I (bufbd1)	0.00	9.92 r
clktree_2/Z (bufbd1)	0.08 *	10.00 r
dout_reg/CP (dfnrb1)	0.00	10.00 r
clock reconvergence pessimism	0.10	10.10
library setup time	-0.08	10.02
data required time		10.02
-----		
data required time		10.02
data arrival time		-0.64
-----		
slack (MET)		9.38

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg -delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
clock pllout (rise edge)	0.00	0.00
clock source latency	-2.08	-2.08
PLL/OUT (DUMMYPLL)	0.00	-2.08 r
clktree_root/I (bufbd1)	0.00	-2.08 r
clktree_root/Z (bufbd1)	2.00 *	-0.08 r
clktree_1/I (bufbd1)	0.00	-0.08 r
clktree_1/Z (bufbd1)	0.27 *	0.19 r
din_reg/CP (dfnrb1)	0.00	0.19 r
din_reg/Q (dfnrb1) <-	0.32 *	0.51 f
dout_reg/D (dfnrb1)	0.00	0.51 f
data arrival time		0.51
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	-1.98	-1.98
PLL/OUT (DUMMYPLL)	0.00	-1.98 r
clktree_root/I (bufbd1)	0.00	-1.98 r
clktree_root/Z (bufbd1)	2.00 *	0.02 r
clktree_2/I (bufbd1)	0.00	0.02 r
clktree_2/Z (bufbd1)	0.08 *	0.10 r
dout_reg/CP (dfnrb1)	0.00	0.10 r
clock reconvergence pessimism	-0.10	0.00
library hold time	0.01	0.01
data required time		0.01
-----		
data required time		0.01
data arrival time		-0.51
-----		
slack (MET)		0.50

So, neither all-max nor all-min gives the correct timing. However, running each separately and combining the results *will* give the correct result. Recall that when each approach was wrong, it was optimistic. So, the minimum slack will always refer to the more pessimistic, i.e. correct, result ***if you run both.***

So, instead of running “worst-case” analysis and “best-case” analysis, we will have to run “worst-case, early pll”, “worst-case, late pll”, “best-case, early pll”, and “best-case, late pll” (where early pll means fb all-max and late pll means fb all-min).

So, how do we force the all-min and all-max values onto the gates in the fb loop? To illustrate this better, I'm going to change the example a little bit. Up to now, we have ignored net delays to simplify the examples. Now I want to include at least one net delay. So that the expected slack numbers don't change, I'll split the delay in the "fbdelay" buffer into 2 components – a cell delay and a net delay, like this:

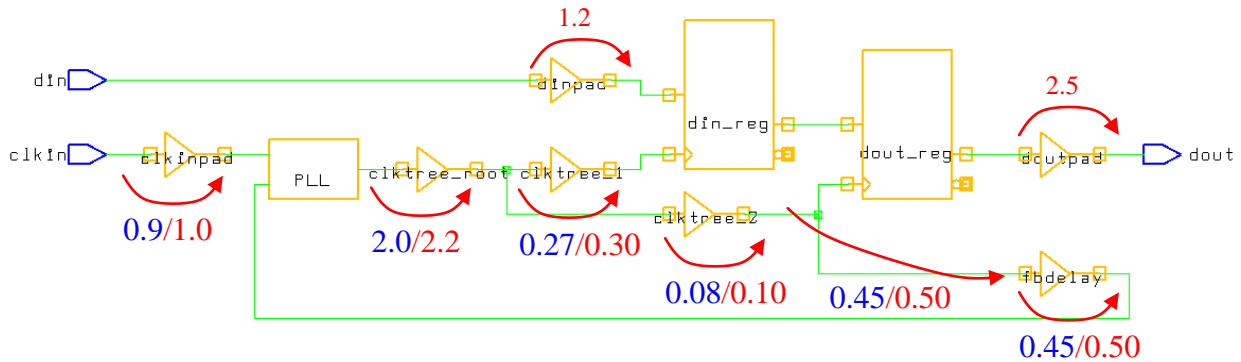


Figure 6-2

I've done this so that the net delay will be seen more explicitly.

There are a variety of ways to force the delay values along the feedback loop. You could use perl to modify the SDF and create fb-max and fb-min (early\_pll and late\_pll) versions. You could use the tricks shown in reference [4] to perl the output of report\_timing to create set\_annotated\_delay commands, or even a mini-SDF file. But I think the most straightforward way is to use get\_timing\_paths to fetch the desired timing information after the OCV SDF has been read in and then use set\_annotated\_delay to force the desired value for min and max.

The basic approach is to do a get\_timing\_paths on the feedback path, then extract a collection called "points", which is an attribute of the path. We can then foreach (foreach\_in\_collection) our way through these points, and use the arrival attribute of each point to calculate the appropriate value for set\_annotated\_delay.

So, here's my proc for doing this - `&build_cmd_to_fix_delays`. It takes one argument – the path returned from `get_timing_paths`. It returns a command for setting the annotated delays:

```

proc &build_cmd_to_fix_delays {_path} {
  # init internal variables
  set _cmd {}
  set _prevpin {}
  set _prevarrival 0
  # Go through points in path
  foreach_in_collection _point [get_attribute $_path points] {
    set _object [get_attribute $_point object]
    set _pinname [get_object_name $_object]
    set _class [get_attribute $_object object_class]
    set _arrival [get_attribute $_point arrival]
    # If this isn't the first pin
    if {$_prevpin != {}} {
      # And it *is* a pin (not a port)
      if {$_class == "pin"} {
        set _pin_direction [get_attribute $_object pin_direction]
        # If this is an input, then the path is a net
        if {$_pin_direction == "in"} {
          set_nextcmd "set_annotated_delay -net "
          # If this is an output, then the path is a cell
        } elseif {$_pin_direction == "out"} {
          set_nextcmd "set_annotated_delay -cell "
        } else {
          echo "Err" "or - $_pinname not an in or out port - direction is
$_pin_direction"
        }
        # Calculate incremental delay
        set_incr [expr $_arrival - $_prevarrival]
        # Finish command
        set_nextcmd "$_nextcmd -from $_prevpin -to $_pinname $_incr ;\n"
        # Append to list of commands
        set _cmd "$_cmd $_nextcmd"
      }
    }
    # Prepare for next
    set _prevpin $_pinname
    set _prevarrival $_arrival
  }
  # Return list of commands
  return $_cmd
}

```



A few notes on this code:

- 1) The `get_timing_paths` command itself (using the `-delay` option) will determine whether the min or max values are used in the `set_annotated_delay` commands.
- 2) The `set_annotated_delay` command needs to know whether the delay belongs to a net or a gate. To determine this, we can look at the pin direction of the point under consideration. If it is an input, then the previous point must have been an output, and the path is a net. If it is an output, then the previous point must have been an input, and the path is a “cell” (gate) path.
- 3) The check for “`$_class == pin`” is to handle the case where the timing path starts or ends on a port. The “net” between the port and the pad pin is a Verilog anomaly –see reference [1].
- 4) We don’t really want to do the `set_annotated_delay` commands as we go, because this will cause many unnecessary timing updates. Instead, we will build a set of `set_annotated_delay` commands in a variable for later processing.

Now we can use this proc in our script.

First, we turn on CRPR and create the clocks as usual:

```
set timing_remove_clock_reconvergence_pessimism true

create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout
```

Now we get the feedback path in the usual way. Notice, however, that the `delay_type` argument is now based on a variable “`_delay_type`”. This would be set to “`max_rise`” for `early_pll` and “`min_rise`” for `late_pll`.

```
set _path [get_timing_paths -delay $_delay_type \
  -from [get_pins PLL/OUT] \
  -to [get_pins PLL/FB] \
]
```

Since we specified the delay type with `get_timing_paths`, it is perfectly safe to use the arrival value as the Tfb. We will be overriding the delays along this path, but the override values will be the same as what was used on this path.

```
set _fb_delay [get_attribute $_path arrival]
```

Now we can create the `set_annotated_delay` commands.

```
set _sadcmts [&build_cmd_to_fix_delays $_path]
```

We put the command string in a variable for later execution. We'll execute it at the end before generating reports to avoid an unnecessary timing updates.

Now we fetch the min and max refclk delays:

```
set _path [get_timing_paths -delay max_rise \  
  -from [get_ports clkkin] \  
  -to [get_pins PLL/CKREF] \  
]  
set _ref_delay_max [get_attribute $_path arrival]  
  
set _path [get_timing_paths -delay min_rise \  
  -from [get_ports clkkin] \  
  -to [get_pins PLL/CKREF] \  
]  
set _ref_delay_min [get_attribute $_path arrival]
```

Now we can set the min and max source latencies. But now the min and max come only from the min and max ref delays; there's only one feedback delay value.

```
set_clock_latency -early -source \  
  [expr $_ref_delay_min - $_fb_delay] \  
  [get_clocks pllout]  
  
set_clock_latency -late -source \  
  [expr $_ref_delay_max - $_fb_delay] \  
  [get_clocks pllout]
```

We set the i/o constraints as usual:

```
set_input_delay -max 8.0 -clock clkkin [get_ports din]  
set_input_delay -min 0.5 -clock clkkin [get_ports din]  
set_output_delay -max 2.0 -clock clkkin [get_ports dout]  
set_output_delay -min [expr -1.0 * 0.5] -clock clkkin [get_ports dout]
```

And now we finally execute the set\_annotated\_delay command string:

```
eval $_sadcmts
```

If we run all this with \$\_delay\_type set to "min\_rise", and generate summary reports, we get:

```
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay max  
*****  
Report : timing  
        -path summary  
        -delay max  
        -max_paths 100  
Design : ocv_pll  
Version: V-2004.06  
*****
```

Startpoint	Endpoint	Slack
dout_reg/CP (dfnrb1)	dout (out)	5.08
din (in)	din_reg/D (dfnrb1)	0.91
din_reg/CP (dfnrb1)	dout_reg/D (dfnrb1)	9.38

```
1
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay min
*****
```

```
Report : timing
        -path summary
        -delay min
        -max_paths 100
```

```
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint	Endpoint	Slack
dout_reg/CP (dfnrb1)	dout (out)	2.32
din_reg/CP (dfnrb1)	dout_reg/D (dfnrb1)	0.50
din (in)	din_reg/D (dfnrb1)	1.37

If we run it with `$_delay_type` set to “max\_rise”, the summary reports look like this:

```
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay max
*****
```

```
Report : timing
        -path summary
        -delay max
        -max_paths 100
```

```
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint	Endpoint	Slack
dout_reg/CP (dfnrb1)	dout (out)	5.18
din (in)	din_reg/D (dfnrb1)	0.79
din_reg/CP (dfnrb1)	dout_reg/D (dfnrb1)	9.40

```
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay min
*****
```

```
Report : timing
        -path summary
        -delay min
        -max_paths 100
```

```
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint	Endpoint	Slack
dout_reg/CP (dfnrb1)	dout (out)	2.22
din_reg/CP (dfnrb1)	dout_reg/D (dfnrb1)	0.48
din (in)	din_reg/D (dfnrb1)	1.49

The slack values in red are the correct calculations. The corresponding path in the other run will be more optimistic and so it will not matter.

Let's take a look at that `_sadcmts` variable. With `$_delay_type` set to "min\_rise", it looks like this:

```
pt_shell> echo $_sadcmts
set_annotated_delay -net -from PLL/OUT -to clktree_root/I 0.0 ;
set_annotated_delay -cell -from clktree_root/I -to clktree_root/Z 2.0 ;
set_annotated_delay -net -from clktree_root/Z -to clktree_2/I
4.29999999998e-05 ;
set_annotated_delay -cell -from clktree_2/I -to clktree_2/Z 0.08 ;
set_annotated_delay -net -from clktree_2/Z -to fbdelay/I 0.45 ;
set_annotated_delay -cell -from fbdelay/I -to fbdelay/Z 0.45 ;
set_annotated_delay -net -from fbdelay/Z -to PLL/FB 0.0 ;
```

You can see it is using min delays (2.0 for the `clktree_root`, for example). You can also see the net delay I inserted (0.45) from `clktree_2/Z` to `fbdelay/I`.

With `$_delay_type` set to "max\_rise", it looks like this:

```
pt_shell> echo $_sadcmts
set_annotated_delay -net -from PLL/OUT -to clktree_root/I 0.0 ;
set_annotated_delay -cell -from clktree_root/I -to clktree_root/Z 2.2 ;
set_annotated_delay -net -from clktree_root/Z -to clktree_2/I
4.29999999998e-05 ;
set_annotated_delay -cell -from clktree_2/I -to clktree_2/Z 0.1 ;
set_annotated_delay -net -from clktree_2/Z -to fbdelay/I 0.5 ;
set_annotated_delay -cell -from fbdelay/I -to fbdelay/Z 0.5 ;
set_annotated_delay -net -from fbdelay/Z -to PLL/FB 0.0 ;
```

Now it's using max delays.

So, we can use this code and run "late\_pll" and "early\_pll" modes to correct for the pll excess pessimism problem.

The code shown above only works for SDF analysis. Since the `set_timing_derate` command will derate our `set_annotated_delay` values, we need to turn off timing derate for the cells and nets that we are working on. This means adding “`set_timing_derate 1.0 [get_cells <cellname>]`” and “`set_timing_derate [get_nets <netname>]`” to the `set_annotated_delay` commands. Without going into detailed proofs, here’s the code that handles parasitics:

```

proc &build_cmd_to_fix_delays {_path} {
  # init internal variables
  set _cmd {}
  set _prevpin {}
  set _prevarrival 0
  # Go through points in path
  foreach_in_collection _point [get_attribute $_path points] {
    set _object [get_attribute $_point object]
    set _pinname [get_object_name $_object]
    set _class [get_attribute $_object object_class]
    set _arrival [get_attribute $_point arrival]
    # If this isn't the first pin
    if {$ _prevpin != {} } {
      # And it *is* a pin (not a port)
      if {$ _class == "pin" } {
        set _pin_direction [get_attribute $_object pin_direction]
        # If this is an input, then the path is a net
        if {$ _pin_direction == "in" } {
          set _sadcmt "set_annotated_delay -net "
          set _netname [get_object_name [all_connected $_object]]
          set _deratecmd "set_timing_derate 1.0 \[get_nets $_netname\] ;\n"
          # If this is an output, then the path is a cell
        } elseif {$ _pin_direction == "out" } {
          set _sadcmt "set_annotated_delay -cell "
          set _cellname [get_object_name [cell_of $_pinname]]
          set _deratecmd "set_timing_derate 1.0 \[get_cells $_cellname\] ;\n"
        } else {
          echo "Err" "or - $_pinname not an in or out port - direction is
$_pin_direction"
        }
        # Calculate incremental delay
        set _incr [expr $_arrival - $_prevarrival]
        # Finish command
        set _sadcmt "$_sadcmt -from $_prevpin -to $_pinname $_incr ;\n"
        # Append to list of commands
        set _cmd "$_cmd $_sadcmt $_deratecmd"
      }
    }
    # Prepare for next
    set _prevpin $_pinname
    set _prevarrival $_arrival
  }
  # Return list of commands
  return $_cmd
}

```

With `$_delay_type` set to “min\_rise”, `$_sdcmds` looks like this:

```
set_annotated_delay -net -from PLL/OUT -to clktree_root/I 0.0 ;
set_timing_derate 1.0 [get_nets pllout] ;
set_annotated_delay -cell -from clktree_root/I -to clktree_root/Z 0.076052 ;
set_timing_derate 1.0 [get_cells clktree_root] ;
set_annotated_delay -net -from clktree_root/Z -to clktree_2/I 3.9e-05 ;
set_timing_derate 1.0 [get_nets clktree] ;
set_annotated_delay -cell -from clktree_2/I -to clktree_2/Z 0.093738 ;
set_timing_derate 1.0 [get_cells clktree_2] ;
set_annotated_delay -net -from clktree_2/Z -to fbdelay/I 4e-05 ;
set_timing_derate 1.0 [get_nets clktree2] ;
set_annotated_delay -cell -from fbdelay/I -to fbdelay/Z 0.065468 ;
set_timing_derate 1.0 [get_cells fbdelay] ;
set_annotated_delay -net -from fbdelay/Z -to PLL/FB 0.0 ;
set_timing_derate 1.0 [get_nets pllfb] ;
```

The actual delays are different (since I didn't load in my SDF), but the important thing to notice is the new “set\_timing\_derate” commands.

## 6.2 Referencing the i/os to the feedback clock

It sure would be nice to find a way to avoid the OCV/PLL excess pessimism problem without having to double our PT runs for tapeout. After a lot of messing around, I have come up with a least a partial workaround for this excess pessimism problem, but it is decidedly weird.

Recall that the problem comes in via source latency which contains path information that is hidden from PT. So, we're going to have to eliminate the source latency somehow. But if we do that, how do we get the timing to match between the refclk and the clock tree?

My basic idea is to reference the i/os to the FB pin on the pll. Remember that this has the same timing as the REF pin on the pll (that's what all that source latency calculation did), so its timing is "near" that of the refclk – it is offset by the refclk delay  $T_{ref}$ .

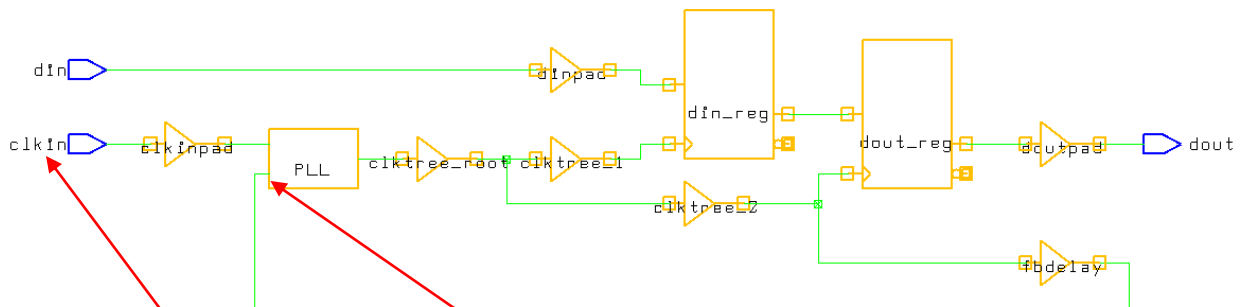


Figure 6-3

The clock *here* is the same as the clock *here*, offset by  $T_{ref}$

So, if we adjust the i/o constraint *values* by  $T_{ref}$  (min or max), we should be able to force the i/o timing to be calculated correctly. But it will be calculated *without reference to refclk*. As we shall see, PT will have the feedback loop timing in the calculation twice, just as with internal paths. So, it will see the clock reconvergence pessimism and apply the correct adjustment!

There won't be any effect on the internal paths, since the source latency, combined with the CRPR adjustment, always cancels.

In detail, the steps are these:

1. Don't apply any source latency to the pll output clock.
2. Create a divide-by 1 generated clock at the FB pin of the pll.
3. Reference input/output delays to this new generated clock, adjusting the value by  $T_{ref\_min}$  or  $T_{ref\_max}$  as required.

The code looks like this:

```
create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout

# Create the gen'd clock on the fb pin
create_generated_clock \
  -source [get_pins PLL/OUT] \
  -name pll_fb_clk \
  -divide_by 1 \
  [get_pins PLL/FB]
set_propagated_clock pll_fb_clk

set_path [get_timing_paths -delay max_rise \
  -from [get_ports clkin] \
  -to [get_pins PLL/CKREF] \
]
set_ref_delay_max [get_attribute $_path arrival]

set_path [get_timing_paths -delay min_rise \
  -from [get_ports clkin] \
  -to [get_pins PLL/CKREF] \
]
set_ref_delay_min [get_attribute $_path arrival]
```

Remember that we have to adjust the input/output delays using the ref clk delay, as follows:

```
# Now set input and output delays using pll_fb_clk
set_input_delay -max \
  -clock pll_fb_clk \
  [expr 8.0 - $_ref_delay_min] \
  [get_ports din]

set_input_delay -min \
  -clock pll_fb_clk \
  [expr 0.5 - $_ref_delay_max] \
  [get_ports din]

set_output_delay -max \
  -clock pll_fb_clk \
  [expr 2.0 + $_ref_delay_max] \
  [get_ports dout]

set_output_delay -min \
  -clock pll_fb_clk \
  [expr (-1.0 * 0.5) + $_ref_delay_min] \
  [get_ports dout]
```

First we'll run this without any OCV to verify it reproduces the slack results when the pll/ocv problem is not present.

Here's what the summary timing reports look like for this circuit using the traditional technique, without OCV:



```

pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay max
*****
Report : timing
        -path summary
        -delay max
        -max_paths 100
Design : ocv_pll
Version: V-2004.06
*****
Startpoint                Endpoint                Slack
-----
dout_reg/CP (dfnrb1)      dout (out)              5.18
din (in)                  din_reg/D (dfnrb1)     0.92
din_reg/CP (dfnrb1)      dout_reg/D (dfnrb1)    9.40

```

```

1
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay min
*****
Report : timing
        -path summary
        -delay max
        -max_paths 100
Design : ocv_pll
Version: V-2004.06
*****
Startpoint                Endpoint                Slack
-----
dout_reg/CP (dfnrb1)      dout (out)              2.32
din_reg/CP (dfnrb1)      dout_reg/D (dfnrb1)    0.51
din (in)                  din_reg/D (dfnrb1)     1.49

```

And the summary reports using these constraints are the same:

```

pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay max
*****
Report : timing
        -path summary
        -delay max
        -max_paths 100
Design : ocv_pll
Version: V-2004.06
*****
Startpoint                Endpoint                Slack
-----
dout_reg/CP (dfnrb1)      dout (out)              5.18
din (in)                  din_reg/D (dfnrb1)     0.92
din_reg/CP (dfnrb1)      dout_reg/D (dfnrb1)    9.40

```

```

1
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay min
*****
Report : timing
        -path summary
        -delay max
        -max_paths 100
Design : ocv_pll
Version: V-2004.06
*****

```

Startpoint	Endpoint	Slack
dout_reg/CP (dfnrb1)	dout (out)	2.32
din_reg/CP (dfnrb1)	dout_reg/D (dfnrb1)	0.51
din (in)	din_reg/D (dfnrb1)	1.49

So we didn't break anything.

Now let's look at the din path with OCV *turned on* using the new constraints:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din (input port clocked by pll\_fb\_clk)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
clock pll_fb_clk (rise edge)	0.00	0.00
clock pllout (source latency)	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_2/I (bufbd1)	0.00	2.20 r
clktree_2/Z (bufbd1)	0.10 *	2.30 r
fbdelay/I (bufbd1)	0.00	2.30 r
fbdelay/Z (bufbd1)	1.00 *	3.30 r
PLL/FB (DUMMYPLL) (gclock source)	0.00	3.30 r
input external delay	7.10	10.40 r
din (in)	0.00	10.40 r
dinpad/I (bufbd1)	0.00	10.40 r
dinpad/Z (bufbd1)	1.20 *	11.60 r
din_reg/D (dfnrb1)	0.00	11.60 r
data arrival time		11.60
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	0.00	10.00
PLL/OUT (DUMMYPLL)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.00 *	12.00 r
clktree_1/I (bufbd1)	0.00	12.00 r
clktree_1/Z (bufbd1)	0.27 *	12.27 r
din_reg/CP (dfnrb1)	0.00	12.27 r
clock reconvergence pessimism	0.20	12.47
library setup time	-0.08	12.39
data required time		12.39
-----		
data required time		12.39
data arrival time		-11.60
-----		
slack (MET)		0.79

Now we get the expected slack value of 0.79. Look carefully at the report. Both paths now go through the clock tree, so PT recognizes that it used the min delay for clktree\_root in one place and max delay in another and gives us back the lost 0.20 as clock reconvergence pessimism.

How about hold on din?

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from din -
delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din (input port clocked by pll\_fb\_clk)  
 Endpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
-----		
clock pll_fb_clk (rise edge)	0.00	0.00
clock pllout (source latency)	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.00 *	2.00 r
clktree_2/I (bufbd1)	0.00	2.00 r
clktree_2/Z (bufbd1)	0.08 *	2.08 r
fbdelay/I (bufbd1)	0.00	2.08 r
fbdelay/Z (bufbd1)	0.90 *	2.98 r
PLL/FB (DUMMYPLL) (gclock source)	0.00	2.98 r
input external delay	-0.50	2.48 f
din (in)	0.00	2.48 f
dinpad/I (bufbd1)	0.00	2.48 f
dinpad/Z (bufbd1)	1.20 *	3.68 f
din_reg/D (dfnrb1)	0.00	3.68 f
data arrival time		3.68
clock pllout (rise edge)	0.00	0.00
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_1/I (bufbd1)	0.00	2.20 r
clktree_1/Z (bufbd1)	0.30 *	2.50 r
din_reg/CP (dfnrb1)	0.00	2.50 r
clock reconvergence pessimism	-0.20	2.30
library hold time	0.01	2.31
data required time		2.31
-----		
data required time		2.31
data arrival time		-3.68
-----		
slack (MET)		1.37

Right again!

### Internal path setup:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_1/I (bufbd1)	0.00	2.20 r
clktree_1/Z (bufbd1)	0.30 *	2.50 r
din_reg/CP (dfnrb1)	0.00	2.50 r
din_reg/Q (dfnrb1) <-	0.32 *	2.82 r
dout_reg/D (dfnrb1)	0.00	2.82 r
data arrival time		2.82
clock pllout (rise edge)	10.00	10.00
clock source latency	0.00	10.00
PLL/OUT (DUMMYPLL)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.00 *	12.00 r
clktree_2/I (bufbd1)	0.00	12.00 r
clktree_2/Z (bufbd1)	0.08 *	12.08 r
dout_reg/CP (dfnrb1)	0.00	12.08 r
clock reconvergence pessimism	0.20	12.28
library setup time	-0.08	12.20
data required time		12.20
-----		
data required time		12.20
data arrival time		-2.82
-----		
slack (MET)		9.38

This is correct. Note that the slack has been reduced by 0.02 relative to the non-OCV case. This is the effect of clktree\_2 being faster that I pointed out earlier. PT has done it right.

**Internal path hold:**

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din_reg -to dout_reg -delay min
*****
Report : timing
        -path full_clock_expanded
        -delay min
        -input_pins
        -max_paths 1
Design : ocv_pll
Version: V-2004.06
*****
```

Startpoint: din\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Endpoint: dout\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: min

Point	Incr	Path
-----		
clock pllout (rise edge)	0.00	0.00
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.00 *	2.00 r
clktree_1/I (bufbd1)	0.00	2.00 r
clktree_1/Z (bufbd1)	0.27 *	2.27 r
din_reg/CP (dfnrb1)	0.00	2.27 r
din_reg/Q (dfnrb1) <-	0.32 *	2.59 f
dout_reg/D (dfnrb1)	0.00	2.59 f
data arrival time		2.59
clock pllout (rise edge)	0.00	0.00
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.00	0.00 r
clktree_root/Z (bufbd1)	2.20 *	2.20 r
clktree_2/I (bufbd1)	0.00	2.20 r
clktree_2/Z (bufbd1)	0.10 *	2.30 r
dout_reg/CP (dfnrb1)	0.00	2.30 r
clock reconvergence pessimism	-0.20	2.10
library hold time	0.01	2.11
data required time		2.11
-----		
data required time		2.11
data arrival time		-2.59
-----		
slack (MET)		0.48

This is also correct.

Here are the complete summary reports for setup and hold. They match the desired results (the circled values in the summary reports from the traditional workaround) exactly:

```
pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay max
*****
```

```
Report : timing
        -path summary
        -delay max
        -max_paths 100
```

```
Design : ocv_pll
Version: V-2004.06-SP1
*****
```

Startpoint	Endpoint	Slack
-----		
dout_reg/CP (dfnrb1)	dout (out)	5.08
din (in)	din_reg/D (dfnrb1)	0.79
din_reg/CP (dfnrb1)	dout_reg/D (dfnrb1)	9.38

```

pt_shell> report_timing -path_type summary -max_paths 100 -nworst 1 -delay min
*****
Report : timing
        -path summary
        -delay min
        -max_paths 100
Design : ocv_pll
Version: V-2004.06-SP1
*****

Startpoint                                Endpoint                                  Slack
-----
dout_reg/CP (dfnrb1)                      dout (out)                               2.22
din_reg/CP (dfnrb1)                      dout_reg/D (dfnrb1)                     0.48
din (in)                                  din_reg/D (dfnrb1)                      1.37

```

I have tried this trick with all sorts of circuits and paths, and it seems to work – as long as no flops are directly connected to clkIn (see below). But I have yet to apply it in a tape-out situation. I'm not sure if I will in the future. Although it works, the reports don't model reality, so it makes me nervous. So, proceed with caution.

### 6.3 The shell game

But there's still a fly in the ointment. What if the design has registers directly clocked by clkIn?

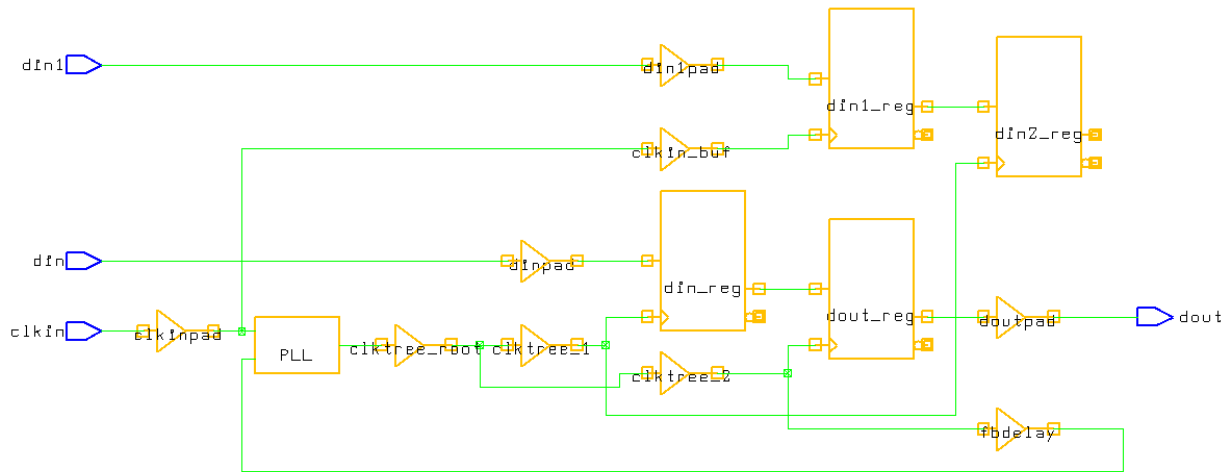


Figure 6-4

Now what? Even without OCV, this isn't going to work. With OCV turned off, this is what the path between din1\_reg (on clkIn) and din2\_reg (on pllout) should look like:

```

pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din1_reg -to din2_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll_wclkin
Version: V-2004.06
*****

```

```

Startpoint: din1_reg (rising edge-triggered flip-flop clocked by clkkin)
Endpoint:   din2_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type:  max

```

Point	Incr	Path
-----		
clock clkkin (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clkkin (in)	0.00	0.00 r
clkkinpad/I (bufbd1)	0.00	0.00 r
clkkinpad/Z (bufbd1)	1.00 *	1.00 r
clkkin_buf/I (bufbd1)	0.00	1.00 r
clkkin_buf/Z (bufbd1)	0.40 *	1.40 r
din1_reg/CP (dfnrb1)	0.00	1.40 r
din1_reg/Q (dfnrb1) <-	0.32	1.72 r
din2_reg/D (dfnrb1)	0.00	1.72 r
data arrival time		1.72
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	-2.30	7.70
PLL/OUT (DUMMYPLL)	0.00	7.70 r
clktree_root/I (bufbd1)	0.00	7.70 r
clktree_root/Z (bufbd1)	2.20 *	9.90 r
clktree_1/I (bufbd1)	0.00	9.90 r
clktree_1/Z (bufbd1)	0.30 *	10.20 r
din2_reg/CP (dfnrb1)	0.00	10.20 r
clock reconvergence pessimism	0.00	10.20
library setup time	-0.08	10.12
data required time		10.12
-----		
data required time		10.12
data arrival time		-1.72
-----		
slack (MET)		8.40



But this is what it looks like (again, with OCV turned off) using this new technique:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from
din1_reg -to din2_reg
*****
Report : timing
        -path full_clock_expanded
        -delay max
        -input_pins
        -max_paths 1
Design : ocv_pll_wclkin
Version: V-2004.06
*****
```

```
Startpoint: din1_reg (rising edge-triggered flip-flop clocked by clkin)
Endpoint:   din2_reg (rising edge-triggered flip-flop clocked by pllout)
Path Group: pllout
Path Type:  max
```

Point	Incr	Path
-----		
clock clkin (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clkin (in)	0.00	0.00 r
clkinpad/I (bufbd1)	0.00	0.00 r
clkinpad/Z (bufbd1)	1.00 *	1.00 r
clkin_buf/I (bufbd1)	0.00	1.00 r
clkin_buf/Z (bufbd1)	0.40 *	1.40 r
din1_reg/CP (dfnrb1)	0.00	1.40 r
din1_reg/Q (dfnrb1) <-	0.32	1.72 r
din2_reg/D (dfnrb1)	0.00	1.72 r
data arrival time		1.72
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	0.00	10.00
PLL/OUT (DUMMYPLL)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.20 *	12.20 r
clktree_1/I (bufbd1)	0.00	12.20 r
clktree_1/Z (bufbd1)	0.30 *	12.50 r
din2_reg/CP (dfnrb1)	0.00	12.50 r
clock reconvergence pessimism	0.00	12.50
library setup time	-0.08	12.42
data required time		12.42
-----		
data required time		12.42
data arrival time		-1.72
-----		
slack (MET)		10.70

The problem is that we are no longer taking the PLL's time shift into account. The trick of referencing the i/os to the feedback clock basically moved the operation out in time. If we leave clkin at time zero, the path from din1\_reg to din2\_reg (from clkin to the pll clock) will be all messed up - it will fail to take into account the behavior of the pll.

We can fix this by adding the appropriate amount of source latency to clkkin:

```
set_clock_latency -source \  
  [expr $_fb_delay - $_ref_delay] \  
  [get_clocks clkkin]
```

Now the path works (without OCV):

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from  
din1_reg -to din2_reg
```

```
*****
```

```
Report : timing  
  -path full_clock_expanded  
  -delay max  
  -input_pins  
  -max_paths 1
```

```
Design : ocv_pll_wclkkin
```

```
Version: V-2004.06
```

```
*****
```

```
Startpoint: din1_reg (rising edge-triggered flip-flop clocked by clkkin)  
Endpoint: din2_reg (rising edge-triggered flip-flop clocked by pllout)  
Path Group: pllout  
Path Type: max
```

Point	Incr	Path
-----		
clock clkkin (rise edge)	0.00	0.00
clock source latency	2.30	2.30
clkkin (in)	0.00	2.30 r
clkkinpad/I (bufbd1)	0.00	2.30 r
clkkinpad/Z (bufbd1)	1.00 *	3.30 r
clkkin_buf/I (bufbd1)	0.00	3.30 r
clkkin_buf/Z (bufbd1)	0.40 *	3.70 r
din1_reg/CP (dfnrb1)	0.00	3.70 r
din1_reg/Q (dfnrb1) <-	0.32	4.02 r
din2_reg/D (dfnrb1)	0.00	4.02 r
data arrival time		4.02
clock pllout (rise edge)	10.00	10.00
clock source latency	0.00	10.00
PLL/OUT (DUMMYPLL)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.20 *	12.20 r
clktree_1/I (bufbd1)	0.00	12.20 r
clktree_1/Z (bufbd1)	0.30 *	12.50 r
din2_reg/CP (dfnrb1)	0.00	12.50 r
clock reconvergence pessimism	0.00	12.50
library setup time	-0.08	12.42
data required time		12.42
-----		
data required time		12.42
data arrival time		-4.02
-----		
slack (MET)		8.40

Now let's turn on OCV. We'll have to do the source latency on clkin using both min and max values:

```
set_clock_latency -early -source \  
  [expr $_fb_delay_min - $_ref_delay_max] \  
  [get_clocks clkin]
```

```
set_clock_latency -late -source \  
  [expr $_fb_delay_max - $_ref_delay_min] \  
  [get_clocks clkin]
```

Now let's look at the timing trace:

```
pt_shell> report_timing -input_pins -path_type full_clock_expanded -from  
din1_reg -to din2_reg  
*****  
Report : timing  
  -path full_clock_expanded  
  -delay max  
  -input_pins  
  -max_paths 1  
Design : ocv_pll_wclkin  
Version: V-2004.06  
*****
```

Startpoint: din1\_reg (rising edge-triggered flip-flop clocked by clkin)  
 Endpoint: din2\_reg (rising edge-triggered flip-flop clocked by pllout)  
 Path Group: pllout  
 Path Type: max

Point	Incr	Path
clock clkin (rise edge)	0.00	0.00
clock source latency	2.40	2.40
clkin (in)	0.00	2.40 r
clkinpad/I (bufbd1)	0.00	2.40 r
clkinpad/Z (bufbd1)	1.00 *	3.40 r
clkin_buf/I (bufbd1)	0.00	3.40 r
clkin_buf/Z (bufbd1)	0.40 *	3.80 r
din1_reg/CP (dfnrb1)	0.00	3.80 r
din1_reg/Q (dfnrb1) <-	0.32	4.12 r
din2_reg/D (dfnrb1)	0.00	4.12 r
data arrival time		4.12
-----		
clock pllout (rise edge)	10.00	10.00
clock source latency	0.00	10.00
PLL/OUT (DUMMYPLL)	0.00	10.00 r
clktree_root/I (bufbd1)	0.00	10.00 r
clktree_root/Z (bufbd1)	2.00 *	12.00 r
clktree_1/I (bufbd1)	0.00	12.00 r
clktree_1/Z (bufbd1)	0.27 *	12.27 r
din2_reg/CP (dfnrb1)	0.00	12.27 r
clock reconvergence pessimism	0.00	12.27
library setup time	-0.08	12.19
data required time		12.19
-----		
data required time		12.19
data arrival time		-4.12
-----		
slack (MET)		8.07

The value of “2.40” for clock source latency comes from the “max”, which is  $\$_{fb\_delay\_max} - \$_{ref\_delay\_min}$ .  $\$_{fb\_delay\_max}$ , in turn, comes from having all the gates in the feedback path at max. But the bottom half of the trace uses clktree\_root at 2.0 – its min value. And the clock reconvergence pessimism value is 0. So, even without going through the detail case1/case2 analysis, we can see that the OCV/PLL excess pessimism is back.

I also tried “superimposing” the two approaches – creating and propagating both the time-zero and source-latency pll output clocks and controlling which ones interact – only allowing clkin to interact with the adjusted pllout clock. That also makes the timing work without OCV, but still exhibits the excess pessimism problem.

I think the problem here is something fundamental. Any time that there is a path that crosses between the clkin domain and the pllout domain, there will be excess pessimism, because one or the other of them will always be a calculated value that PT won’t know contains min/max paths.

The reason the “reference the ios to the feedback clock” trick worked was that I eliminated any reference to clkin (I could do a remove\_clock on clkin and it wouldn’t change the timing

reports). Once there are flops directly tied to clkin, there is no way to avoid this domain crossing.

You can, however, move the problem around. Originally, the excess pessimism appears on the i/o paths. By referencing the i/os to the feedback clock, and either adding source latency to clkin or propagating both sets of clocks, you can move the problem to the din1\_reg/din2\_reg interface – an internal interface, where there *might* be more timing slack available.

## **7 Conclusion**

Complete and accurate modeling of PLLs is a surprisingly complex task, but a very important one. There are a lot of issues and effects to consider. Hopefully this paper will provide the reader with the tools necessary to tackle the project successfully.

## **8 Acknowledgements**

The author would like to acknowledge the following people for their assistance and review:

Karthik Rajan for doing a complete, thorough review of the entire paper.

Chris Papademetrious for doing a complete, thorough review of the entire paper.

Matt Weber for lending his expertise in OCV and his insights on jitter, as well as slogging his way through the entire paper.

Steve Golson for his insights on the issues with PLL models in a parasitics environment, jitter and for his help in reviewing the entire paper.

## 9 References

- (1) Complex Clocking Situations Using PrimeTime  
Paul Zimmer  
Synopsys Users Group 2000 San Jose  
(available at [www.zimmerdesignservices.com](http://www.zimmerdesignservices.com))
- (2) Working with DDRs in PrimeTime  
Paul Zimmer, Andrew Cheng  
Synopsys Users Group 2001 San Jose  
(available at [www.zimmerdesignservices.com](http://www.zimmerdesignservices.com))
- (3) My Favorite DC/PT Shell Tricks  
Paul Zimmer  
Synopsys Users Group 2002 San Jose  
(available at [www.zimmerdesignservices.com](http://www.zimmerdesignservices.com))
- (4) Accounting for PLL Correction in Static Timing Analysis  
Solvnet article 000096.html
- (5) My Head Hurts, My Timing Stinks, and I Don't Love On-Chip Variation  
Matt Weber – Silicon Logic Engineering  
Available from Solvnet



## 10 Appendix

### 10.1 The PLL model itself

When using primetime with parasitics rather than SDF, it is important that the PLL model correctly reflect the drive and loading properties of the real PLL device. Although in many cases the delays may cancel, the transition time effects can propagate. This is particularly true if the clock tree connected to the PLL output doesn't begin with a single buffer.

Getting this right isn't as easy as you might think.

#### 10.1.1 Modeling the PLL using a verilog wrapper

One way of modeling the PLL would be to find out from the vendor what sort of standard buffers best represent the PLL's inputs and outputs, then build a verilog model "wrapper" around the empty PLL shell. In practice, this may cause problems because the parasitic extraction files will refer to pins on the PLL that are now just internal hierarchy pins. So, you probably can't use this technique directly without hacking the extraction files. Still, it is useful to illustrate some of the issues to be dealt with and to serve as a baseline.

For example, if the PLL instantiation looks like this:

```
PLLXYZ PLL (.OUT(p11out), .FB(p11fb), .CKREF(iclkin) );
```

I'll build a 2-stage verilog model that looks like this:

```
module PLLXYZ (
    OUT,
    FB,
    CKREF
);

output OUT;
input CKREF;
input FB;

    DUMMYPLL DUMMYPLL (.DUMMYOUT(pllout), .DUMMYFB(pllfb),
.DUMMYCKREF(pllrefclk) );

bufbd1 fbinbuf(.I(FB), .Z(pllfb));
bufbd1 ckrefinbuf(.I(CKREF), .Z(pllrefclk));
bufbd1 outbuf(.I(pllout), .Z(OUT));

endmodule

module DUMMYPLL (
    DUMMYOUT,
    DUMMYFB,
    DUMMYCKREF
);

output DUMMYOUT;
input DUMMYCKREF;
input DUMMYFB;

endmodule
```

In this case, I have modeled the PLL's i/os as standard library buffers (bufbd1's).

Now for the tricky bit. If I load up the data and report the timing before I have created any clocks, I get something like this:

```
pt_shell> set timing_report_unconstrained_paths true
true
pt_shell> report_timing -input_pins -to [get_pins PLL/DUMMYPLL/DUMMYFB]
*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****
```

```

Startpoint: PLL/outbuf/I
              (internal pin)
Endpoint: PLL/DUMMYPLL/DUMMYFB
              (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
PLL/outbuf/I (bufbd1)	0.00	0.00 r
PLL/outbuf/Z (bufbd1)	0.30	0.30 r
PLL/OUT (PLLXYZ)	0.00 +	0.30 r
clktree_root/I (bufbd1)	0.04 +	0.34 r
clktree_root/Z (bufbd1)	0.41 +	0.75 r
clktree_2/I (bufbd1)	0.04 +	0.79 r
clktree_2/Z (bufbd1)	0.40 +	1.19 r
fbdelay/I (bufbd1)	0.04 +	1.23 r
fbdelay/Z (bufbd1)	0.40 +	1.63 r
PLL/FB (PLLXYZ)	0.00	1.63 r
PLL/fbinbuf/I (bufbd1)	0.04	1.67 r
PLL/fbinbuf/Z (bufbd1)	0.14	1.81 r
PLL/DUMMYPLL/DUMMYFB (DUMMYPLL)	0.00	1.81 r
data arrival time		1.81

(Path is unconstrained)

Now create the clocks. The output of the pll is now “PLL/outbuf/Z”, so I’ll create the pll output clock there (“PLL/OUT” is no longer a real physical pin).

```

set _plloutpin_name {PLL/outbuf/Z}

create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins ${_plloutpin_name}]
set_propagated_clock pllout

```

If I re-run the report, I get this:

```

pt_shell> report_timing -input_pins -to [get_pins PLL/DUMMYPLL/DUMMYFB]
*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

```

```

Startpoint: PLL/outbuf/Z
             (clock source 'pllout')
Endpoint: PLL/DUMMYPLL/DUMMYFB
           (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
clock source latency	0.00	0.00
PLL/outbuf/Z (bufbd1)	0.00	0.00 r
PLL/OUT (PLLXYZ)	0.00 +	0.00 r
clktree_root/I (bufbd1)	0.04 +	0.04 r
clktree_root/Z (bufbd1)	0.31 +	0.35 r
clktree_2/I (bufbd1)	0.04 +	0.39 r
clktree_2/Z (bufbd1)	0.40 +	0.79 r
fbdelay/I (bufbd1)	0.04 +	0.83 r
fbdelay/Z (bufbd1)	0.40 +	1.23 r
PLL/FB (PLLXYZ)	0.00	1.23 r
PLL/fbinbuf/I (bufbd1)	0.04	1.27 r
PLL/fbinbuf/Z (bufbd1)	0.14	1.41 r
PLL/DUMMYPLL/DUMMYFB (DUMMYPLL)	0.00	1.41 r
data arrival time		1.41

(Path is unconstrained)

The startpoint has changed due to the clock, so the delay of PLL/outbuf doesn't show up in the report, which is to be expected. But something else has changed. Notice that the delay through clktree\_root has changed (from 0.41 to 0.31). What happened?

If we do “report\_delay\_calculation -from clktree\_root/I -to clktree\_root/Z” before and after the clock creation commands and compare the files, we find that the root cause of the change is that:

(X) input\_pin\_transition = 0.610997

changed to:

(X) input\_pin\_transition = 0

The create\_clock has a side-effect of forcing a 0 transition time on the clock driving pin. Create\_generated\_clock does the same thing.

This transition time should be a function of the cell drive characteristics and the load. It shouldn't be 0.

One way around this is to create the clock on the “empty shell” pin PLL/DUMMYPLL/DUMMYOUT. The problem is that we'll then get the net delay from there to the PLL/outbuf/I pin, plus the cell delay of PLL/outbuf, in the clock path. In theory, these delays should cancel, but since these aren't real devices, I'd prefer to get rid of them.

So, I'll create the clock on the PLL/outbuf/I pin (which will get rid of the net delay), and use set\_annotated\_delay -cell to get rid of the delay through PLL/outbuf:

```
set_plloutpin_name {PLL/outbuf/I}
set_annotated_delay 0.0 -cell -from PLL/outbuf/I -to PLL/outbuf/Z

create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins ${_plloutpin_name}]
set_propagated_clock pllout
```

Now if we run the reports after the clocks are created, we get the same values as before the clocks are created:

```
pt_shell> report_timing -input_pins -to [get_pins PLL/DUMMYPLL/DUMMYFB]
*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

Startpoint: PLL/outbuf/I
            (clock source 'pllout')
Endpoint:  PLL/DUMMYPLL/DUMMYFB
            (internal pin)
Path Group: (none)
Path Type: max

Point                                     Incr          Path
-----
clock source latency                       0.00          0.00
PLL/outbuf/I (bufbd1)                      0.00          0.00 r
PLL/outbuf/Z (bufbd1)                      0.00 *        0.00 r
PLL/OUT (PLLXYZ)                           0.00 +        0.00 r
clktree_root/I (bufbd1)                    0.04 +        0.04 r
clktree_root/Z (bufbd1)                    0.41 +        0.45 r
clktree_2/I (bufbd1)                       0.04 +        0.49 r
clktree_2/Z (bufbd1)                       0.40 +        0.89 r
fbdelay/I (bufbd1)                         0.04 +        0.93 r
fbdelay/Z (bufbd1)                         0.40 +        1.33 r
PLL/FB (PLLXYZ)                             0.00          1.33 r
PLL/fbinbuf/I (bufbd1)                     0.04          1.36 r
PLL/fbinbuf/Z (bufbd1)                     0.14          1.51 r
PLL/DUMMYPLL/DUMMYFB (DUMMYPLL)            0.00          1.51 r
data arrival time                          1.51
-----
(Path is unconstrained)
```

By the way, if you look at the delay calculation on the PLL/outbuf itself, you find that PT used an input transition time of 0:

```
pt_shell> report_delay_calculation -from PLL/outbuf/I -to PLL/outbuf/Z
*****
Report : delay_calculation
Design : idc_pll_example
Version: V-2004.06
*****

From pin:                                PLL/outbuf/I
To pin:                                   PLL/outbuf/Z
Main Library Units: 1ns 1pF 1kOhm

Library: 'cb13fs120_tsmc_max'
Library Units: 1ns 1pF 1kOhm
Library Cell: 'bufbd1'
arc sense:                                positive_unate
arc type:                                  cell
Units: 1ns 1pF 1kOhm

Rise Delay

cell delay = 0.301583
Table is indexed by
(X) input_pin_transition = 0
. . .
```

What we are telling PT is that the PLL output behaves like a bufbd1 *with a zero input transition time*. This may or may not be accurate – only your PLL vendor can tell you for sure. But since this transition time is NOT a function of the load, you can use set\_annotated\_transition to force this to the desired value:

```

pt_shell> set_annotated_transition 0.50 [get_pins PLL/outbuf/I]
1
pt_shell> report_delay_calculation -from PLL/outbuf/I -to PLL/outbuf/Z
*****
Report : delay_calculation
Design : idc_pll_example
Version: V-2004.06
*****

From pin:                PLL/outbuf/I
To pin:                  PLL/outbuf/Z
Main Library Units:     1ns 1pF 1kOhm

Library: 'cb13fs120_tsmc_max'
Library Units:          1ns 1pF 1kOhm
Library Cell: 'bufbd1'
arc sense:              positive_unate
arc type:                cell
Units:                  1ns 1pF 1kOhm

Rise Delay

    cell delay = 0.382833
    Table is indexed by
    (X) input_pin_transition = 0.5
. . .

```

Note that the fake buffers PLL/fbinbuf and PLL/ckrefinbuf also have non-zero delay. So, the source latency calculation code will need to use the buffer input pins as endpoints:

```

set_pllfbpin_name {PLL/fbinbuf/I}
set_pllrefpin_name {PLL/ckrefinbuf/I}

set_path [get_timing_paths -delay max_rise \
  -from [get_ports clkin] \
  -to [get_pins ${_pllrefpin_name}] \
]

set_ref_delay [get_attribute $_path arrival]

set_path [get_timing_paths -delay max_rise \
  -from [get_pins ${_plloutpin_name}] \
  -to [get_pins ${_pllfbpin_name}] \
]

set_fb_delay [get_attribute $_path arrival]

```

### 10.1.2 Modeling the PLL using a qtm model

The verilog wrapper is useful for illustrating the clock transition time problem, but it isn't the cleanest solution because it won't match the parasitic files. Another solution that will match the parasitic files is to use a QTM (Quick Timing Model) model.

Using the code in reference [4] as an example, here is the code I used to create a QTM model for my PLL:

```
# Make sure library has been read in
if {[sizeof_collection [get_libs -quiet *]] == 0} {
    read_db $cell_lib_max
}

set pllname DUMMYPLL
create_qtm_model $pllname
set_qtm_technology -library cb13fs120_tsmc_max

create_qtm_drive_type -lib_cell bufbd1 plldrive
create_qtm_load_type -lib_cell bufbd1 pllckrefload
create_qtm_load_type -lib_cell bufbd1 pllfbload

create_qtm_port {CKREF FB} -type input
set_qtm_port_load -type pllckrefload -factor 1 CKREF
set_qtm_port_load -type pllfbload -factor 1 FB
create_qtm_port {OUT} -type output
set_qtm_port_drive -type plldrive OUT

#create_qtm_delay_arc -from CKREF -to OUT -value 0.0

report_qtm_model
save_qtm_model -library_cell -output ${pllname} -format {db lib}
```

Notice that I commented out the “create\_qtm\_delay\_arc” line from the example in reference [4]. More on this in the next section. Also, I added code to create drive types and load types and tied the i/os to these types. I modeled the PLL i/os to be bufbd1’s as in the verilog wrapper example.

When I run this in primetime, it creates two files:

```
DUMMYPLL_lib.db
DUMMYPLL.lib
```

The first one is the QTM model. The second one is a .lib description that can be used to create a technology library model.

To link this in, we do something like this:

```
set link_path "$link_path DUMMYPLL_lib.db"
```



If we load everything up and time the feedback path, we get:

```
pt_shell> set timing_report_unconstrained_paths true
true
pt_shell> report_timing -input_pins -to [get_pins PLL/FB]
Information: Using automatic max wire load selection group 'predcaps'. (ENV-003)
Information: Using automatic min wire load selection group 'predcaps'. (ENV-003)
*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

Startpoint: PLL/OUT_drive
            (internal pin)
Endpoint:  PLL/FB (internal pin)
Path Group: (none)
Path Type: max

Point                                     Incr          Path
-----
PLL/OUT_drive (DUMMYPLL)                 0.00          0.00 r
PLL/OUT (DUMMYPLL)                       0.30 +        0.30 r
clktree_root/I (bufbd1)                  0.04 +        0.34 r
clktree_root/Z (bufbd1)                  0.41 +        0.75 r
clktree_2/I (bufbd1)                     0.04 +        0.79 r
clktree_2/Z (bufbd1)                     0.40 +        1.19 r
fbdelay/I (bufbd1)                       0.04 +        1.23 r
fbdelay/Z (bufbd1)                       0.40 +        1.63 r
PLL/FB (DUMMYPLL)                        0.04 +        1.67 r
data arrival time                        1.67

-----
(Path is unconstrained)
```

This trace is very similar to the trace with the verilog wrapper. The delay of the input buffer (fbinbuf) is gone, but the 0.30 delay of the “dummy” driving buffer (outbuf) is there, even though the driver isn’t! So, it isn’t surprising that when we create the clock on PLL/OUT, we have the same transition time problem as we had with the verilog wrapper:

```
create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins PLL/OUT]
set_propagated_clock pllout
```

```

pt_shell> report_timing -input_pins -to [get_pins PLL/FB]
*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

```

```

Startpoint: PLL/OUT (clock source 'pllout')
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
clock source latency	0.00	0.00
PLL/OUT (DUMMYPLL)	0.00	0.00 r
clktree_root/I (bufbd1)	0.04 +	0.04 r
clktree_root/Z (bufbd1)	0.31 +	0.35 r
clktree_2/I (bufbd1)	0.04 +	0.39 r
clktree_2/Z (bufbd1)	0.40 +	0.79 r
fbdelay/I (bufbd1)	0.04 +	0.83 r
fbdelay/Z (bufbd1)	0.40 +	1.23 r
PLL/FB (DUMMYPLL)	0.04 +	1.27 r
data arrival time		1.27

(Path is unconstrained)

But now there's no "outbuf/I" on which to create the clock! Fortunately, there is a funny little port on the QTM model called "OUT\_drive" that is the equivalent of the outbuf/I. So, we'll create the clock there:

```

create_clock -period 10.0 -name clkin [get_ports clkin]
set_propagated_clock clkin

create_clock -period 10.0 -name pllout [get_pins PLL/OUT_drive]
set_propagated_clock pllout

```

```

pt_shell> report_timing -input_pins -to [get_pins PLL/FB]
*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

```

```

Startpoint: PLL/OUT_drive
             (clock_source 'pllout')
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
clock source latency	0.00	0.00
PLL/OUT_drive (DUMMYPLL)	0.00	0.00 r
PLL/OUT_ (DUMMYPLL)	0.30 +	0.30 r
clktree_root/I (bufbd1)	0.04 +	0.34 r
clktree_root/Z (bufbd1)	0.41 +	0.75 r
clktree_2/I (bufbd1)	0.04 +	0.79 r
clktree_2/Z (bufbd1)	0.40 +	1.19 r
fbdelay/I (bufbd1)	0.04 +	1.23 r
fbdelay/Z (bufbd1)	0.40 +	1.63 r
PLL/FB (DUMMYPLL)	0.04 +	1.67 r
data arrival time		1.67

(Path is unconstrained)

This fixes the transition time problem, but we still have the dummy buffer delay. As before, we can use `set_annotated_delay` to zero this value:

```

pt_shell> set_annotated_delay 0.0 -cell -from PLL/OUT_drive -to PLL/OUT
1

```

```

pt_shell> report_timing -input_pins -to [get_pins PLL/FB]

```

```

*****
Report : timing
        -path full
        -delay max
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****

```

```

Startpoint: PLL/OUT_drive
             (clock_source 'pllout')
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
clock source latency	0.00	0.00
PLL/OUT_drive (DUMMYPLL)	0.00	0.00 r
PLL/OUT_ (DUMMYPLL)	0.00 *	0.00 r
clktree_root/I (bufbd1)	0.04 +	0.04 r
clktree_root/Z (bufbd1)	0.41 +	0.45 r
clktree_2/I (bufbd1)	0.04 +	0.49 r
clktree_2/Z (bufbd1)	0.40 +	0.89 r
fbdelay/I (bufbd1)	0.04 +	0.93 r
fbdelay/Z (bufbd1)	0.40 +	1.33 r
PLL/FB (DUMMYPLL)	0.04 +	1.36 r
data arrival time		1.36

(Path is unconstrained)

Since there are no dummy buffers on the inputs, the source latency calculation code can still refer to the PLL/CKREF and PLL/FB pins.

### 10.1.3 Beware of the arc!

Now back to that line I commented out in the QTM model script. This line creates a timing arc from CKREF to OUT. This time I'll generate the model with this arc in place:

```
# Make sure library has been read in
if {[sizeof_collection [get_libs -quiet *]] == 0} {
    read_db $cell_lib_max
}

set pllname DUMMYPLL
create_qtm_model $pllname
set_qtm_technology -library cb13fs120_tsmc_max

create_qtm_drive_type -lib_cell bufbd1 plldrive
create_qtm_load_type -lib_cell bufbd1 pllckrefload
create_qtm_load_type -lib_cell bufbd1 pllfbload

create_qtm_port {CKREF FB} -type input
set_qtm_port_load -type pllckrefload -factor 1 CKREF
set_qtm_port_load -type pllfbload -factor 1 FB
create_qtm_port {OUT} -type output
set_qtm_port_drive -type plldrive OUT

create_qtm_delay_arc -from CKREF -to OUT -value 0.0

report_qtm_model
save_qtm_model -library_cell -output ${pllname} -format {db lib}
```

The reason this is in reference [4] is that they use a slightly different technique for modeling the pll circuitry. This is discussed in more detail in the next part of the appendix.

If I load up this model and report the timing to the FB pin, I get:

```
pt_shell> set timing_report_unconstrained_paths true
true
pt_shell> report_timing -input_pins -delay max_rise -to PLL/FB
*****
Report : timing
        -path full
        -delay max_rise
        -input_pins
        -max_paths 1
Design : idc_pll_example
Version: V-2004.06
*****
```

```

Startpoint: clk_in (input port)
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max

```

Point	Incr	Path
input external delay	0.00	0.00 r
clk_in (in)	0.26 +	0.26 r
clk_inpad/I (bufbd1)	0.04 +	0.30 r
clk_inpad/Z (bufbd1)	0.70 +	1.00 r
PLL/CKREF (DUMMYPLL)	0.10 +	1.10 r
PLL/OUT_drive (DUMMYPLL)	0.00	1.10 r
PLL/OUT (DUMMYPLL)	0.47 +	1.56 r
clktree_root/I (bufbd1)	0.04 +	1.60 r
clktree_root/Z (bufbd1)	0.41 +	2.01 r
clktree_2/I (bufbd1)	0.04 +	2.05 r
clktree_2/Z (bufbd1)	0.40 +	2.45 r
fbdelay/I (bufbd1)	0.04 +	2.49 r
fbdelay/Z (bufbd1)	0.40 +	2.89 r
PLL/FB (DUMMYPLL)	0.04 +	2.93 r
data arrival time		2.93

(Path is unconstrained)

Notice that the path now starts at the clk\_in input port and runs through the PLL. That's fine, but notice also that the delay on the "OUT\_drive" to "OUT" path is now 0.47. Let's look at the delay calculation for this path:

```

pt_shell> report_delay_calculation -from PLL/OUT_drive -to PLL/OUT
*****
Report : delay_calculation
Design : idc_pll_example
Version: V-2004.06
*****

```

```

From pin:                PLL/OUT_drive
To pin:                  PLL/OUT
Main Library Units: 1ns 1pF 1kOhm

Library: 'DUMMYPLL_lib'
Library Units: 1ns 1pF 1kOhm
Library Cell: 'DUMMYPLL'
arc sense:                positive_unate
arc type:                 cell
Units: 1ns 1pF 1kOhm

```

Rise Delay

```

cell delay = 0.466338
Table is indexed by
(X) input_pin_transition = 1.4439
(Y) output_net_total_cap = 0.07875
. . .

```

This transition time comes from the CKREF pin.

```
pt_shell> report_timing -input_pins -delay max_rise -to PLL/FB -
transition_time
```

```
*****
```

```
Report : timing
        -path full
        -delay max_rise
        -input_pins
        -max_paths 1
        -transition_time
```

```
Design : idc_pll_example
```

```
Version: V-2004.06
```

```
*****
```

```
Startpoint: clkkin (input port)
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max
```

Point	Trans	Incr	Path
-----			
input external delay		0.00	0.00 r
clkkin (in)	0.26	0.26 +	0.26 r
clkkinpad/I (bufbd1)	0.26	0.04 +	0.30 r
clkkinpad/Z (bufbd1)	1.44	0.70 +	1.00 r
PLL/CKREF (DUMMYPLL)	1.44	0.10 +	1.10 r
PLL/OUT_drive (DUMMYPLL)	1.44	0.00	1.10 r
PLL/OUT (DUMMYPLL)	0.62	0.47 +	1.56 r
clktree_root/I (bufbd1)	0.62	0.04 +	1.60 r
clktree_root/Z (bufbd1)	0.64	0.41 +	2.01 r
clktree_2/I (bufbd1)	0.64	0.04 +	2.05 r
clktree_2/Z (bufbd1)	0.61	0.40 +	2.45 r
fbdelay/I (bufbd1)	0.61	0.04 +	2.49 r
fbdelay/Z (bufbd1)	0.61	0.40 +	2.89 r
PLL/FB (DUMMYPLL)	0.61	0.04 +	2.93 r
data arrival time			2.93

```
-----
(Path is unconstrained)
```

But this is *not* correct! The PLL output drive should be independent of the CKREF *input* transition time.

If we break this arc, we get back to the expected timing:

```
pt_shell> report_timing -input_pins -delay max_rise -to PLL/FB -
transition_time
*****
Report : timing
        -path full
        -delay max_rise
        -input_pins
        -max_paths 1
        -transition_time
Design : idc_pll_example
Version: V-2004.06
*****

Startpoint: PLL/OUT_drive
            (internal pin)
Endpoint: PLL/FB (internal pin)
Path Group: (none)
Path Type: max

Point                                     Trans      Incr      Path
-----
-
PLL/OUT_drive (DUMMYPLL)                  0.00      0.00      0.00 r
PLL/OUT_ (DUMMYPLL)                       0.61      0.30 +    0.30 r
clktree_root/I (bufbd1)                   0.61      0.04 +    0.34 r
clktree_root/Z (bufbd1)                   0.64      0.41 +    0.75 r
clktree_2/I (bufbd1)                      0.64      0.04 +    0.79 r
clktree_2/Z (bufbd1)                      0.61      0.40 +    1.19 r
fbdelay/I (bufbd1)                       0.61      0.04 +    1.23 r
fbdelay/Z (bufbd1)                       0.61      0.40 +    1.63 r
PLL/FB (DUMMYPLL)                         0.61      0.04 +    1.67 r
data arrival time                         1.67
-----
-
(Path is unconstrained)
```

It might be possible to fix this by forcing transition times, but since I haven't found any advantage in timing through this arc, I'd rather just not have the arc (disable it using `set_disable_timing` if the vendor-supplied model has it).

Note that there's one case where the arc is desirable. If the PLL has a bypass mode for test, then you need this arc to accurately time test mode. Ideally, the model will make this arc conditional on the bypass mode pin(s). But if the vendor hasn't done that, you will need to disable it using `set_disable_timing` when you run functional mode STA.

### 10.1.4 Vendor-supplied PLL models

It should be evident from the examples above that the details of the PLL model can be tricky to get right. If your vendor provides a model of the PLL, make sure you know what's inside. I got one once that was like the verilog wrapper (instantiated buffers), but the buffers inputs were tied low! This caused bizarre behavior in Primetime.

If you're using parasitics, be sure to use "report\_delay\_calculation" and examine the output carefully to make sure the model is doing the right thing. There should not be any edge-rate dependency of the OUT pin on the REF or FB pins (unless the vendor is *sure* this is real).

Or get rid of the vendor model and use the QTM method shown above.

## 10.2 Why I do it my way

Many people use the approach described in the Synopsys App Note (Reference [4]). This paper wouldn't be complete if I didn't explain why I do it differently.

Leaving aside the minor detail of how the source latency is extracted, the Synopsys App Note uses a fundamentally different approach than the one I have described in this paper. They create a dummy model *with a timing arc from the REF pin to the pll OUT pin*. They then apply the source latency as an annotated delay on this arc. The result is that there is only *one clock* – the reference clock, which propagates through the pll with negative delay.

I see several problems with this approach:

1. Duty cycle will be inherited from the ref clock instead of being independent.
2. I don't see how you can model jitter correctly with only one clock.
3. The annotated delay will get scaled by OCV.
4. It creates a link between the input transition time at the ref pin of the PLL and the output driver of the PLL, which isn't correct.

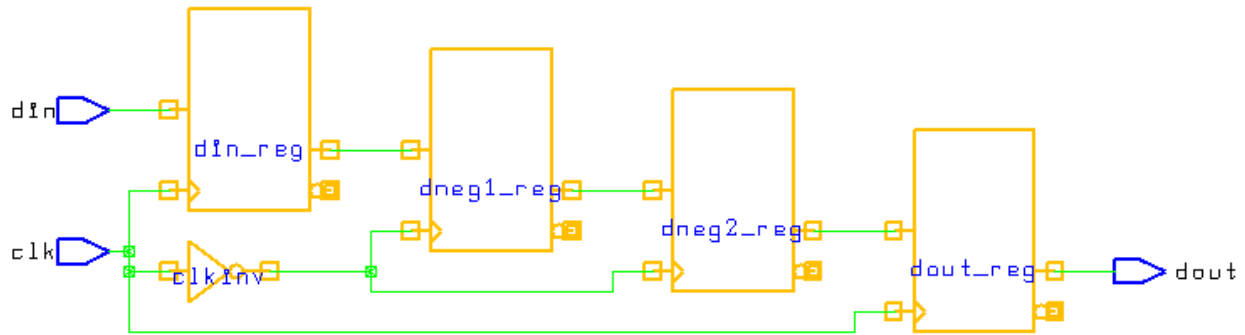
Although there may be solutions to these problems (I've never used this approach, so I've never tried to find any), I think they all stem from a fundamental problem – *this approach doesn't model reality*. A real pll does *not* simply forward the ref clock with a negative delay, it creates a *whole new clock* on its output pin. The characteristics of this clock depend on phase and frequency information from the REF and FB inputs, but the fact remains that it is a *new* clock. That's why my approach creates a new clock on the pll OUT pin.



### 10.3 Modeling duty cycle using set\_clock\_latency early/late

The “set\_clock\_latency” command can be used to model clock duty cycle, but only by making use of features intended for other uses.

Here is our example circuit again:



Notice that set\_clock\_latency has both –early/-late and –min/-max options:

```
pt_shell> help -v set_clock_latency
set_clock_latency      # Capture actual or predicted clock latency
  [-rise]              (Specify clock rise latency)
  [-fall]              (Specify clock fall latency)
  [-min]               (Specify clock rise and fall min condition latency)
  [-max]               (Specify clock rise and fall max condition latency)
  [-source]            (Specify clock rise and fall source latency)
  [-late]              (Specify clock rise and fall late source latency)
  [-early]             (Specify clock rise and fall early source latency)
  delay                (Clock latency)
  object_list          (List of clocks, ports or pins)
```

The –early/-late set corresponds to early and late arrival. The –min/-max set corresponds to early and late arrival in “min-max” analysis mode. Either can be tricked into doing the duty cycle calculation.

We’ll start with –early/-late. If we just try this out-of-the-box, it doesn’t work.

```
set _period 10.0
set _duty_cycle_min 0.40
set _duty_cycle_max [expr 1.0 - $_duty_cycle_min]
create_clock -period $_period -name clk \
  [get_ports clk]
set_propagated_clock clk

set_clock_latency -source -fall -early \
  [expr ($_duty_cycle_min - 0.5) * $_period] \
  [get_clocks clk]
set_clock_latency -source -fall -late \
  [expr ($_duty_cycle_max - 0.5) * $_period] \
  [get_clocks clk]
```

The input and output traces are ok:

```
pt_shell> report_timing -input_pins -path_type full_clock -from din_reg -to
dneg1_reg
```

```
*****
```

```
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
```

```
Design : duty_cycle_piclk
```

```
Version: V-2004.06
```

```
*****
```

```
Startpoint: din_reg (rising edge-triggered flip-flop clocked by clk)
Endpoint:   dneg1_reg (rising edge-triggered flip-flop clocked by clk')
Path Group: clk
Path Type:  max
```

Point	Incr	Path
-----		
clock clk (rise edge)	0.00	0.00
clock source latency	0.00	0.00
clk (in)	0.00	0.00 r
din_reg/CP (dfnrb1)	0.00	0.00 r
din_reg/Q (dfnrb1) <-	0.31	0.31 r
dneg1_reg/D (dfnrb1)	0.00	0.31 r
data arrival time		0.31
clock clk' (rise edge)	5.00	5.00
clock source latency	-1.00	4.00
clk (in)	0.00	4.00 f
clkinv/I (inv0d2)	0.00	4.00 f
clkinv/ZN (inv0d2)	1.00 *	5.00 r
dneg1_reg/CP (dfnrb1)	0.00	5.00 r
library setup time	-0.09	4.91
data required time		4.91
-----		
data required time		4.91
data arrival time		-0.31
-----		
slack (MET)		4.60

```

pt_shell> report_timing -input_pins -path_type full_clock -from dneg2_reg -to
dout_reg
*****
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_piclk
Version: V-2004.06
*****

```

```

Startpoint: dneg2_reg (rising edge-triggered flip-flop clocked by clk')
Endpoint:   dout_reg (rising edge-triggered flip-flop clocked by clk)
Path Group: clk
Path Type:  max

```

Point	Incr	Path
-----		
clock clk' (rise edge)	5.00	5.00
clock source latency	1.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dneg2_reg/CP (dfnrb1)	0.00	7.00 r
dneg2_reg/Q (dfnrb1) <-	0.32	7.32 r
dout_reg/D (dfnrb1)	0.00	7.32 r
data arrival time		7.32
clock clk (rise edge)	10.00	10.00
clock source latency	0.00	10.00
clk (in)	0.00	10.00 r
dout_reg/CP (dfnrb1)	0.00	10.00 r
library setup time	-0.10	9.90
data required time		9.90
-----		
data required time		9.90
data arrival time		-7.32
-----		
slack (MET)		2.59

But the falling-edge to falling-edge trace is not:

```

pt_shell> report_timing -input_pins -path_type full_clock -from dneg1_reg -to
dneg2_reg
*****
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_piclk
Version: V-2004.06
*****

```

Startpoint: dneg1\_reg (rising edge-triggered flip-flop clocked by clk')  
 Endpoint: dneg2\_reg (rising edge-triggered flip-flop clocked by clk')  
 Path Group: clk  
 Path Type: max

Point	Incr	Path
-----		
clock clk' (rise edge)	5.00	5.00
clock source latency	1.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dneg1_reg/CP (dfnrb1)	0.00	7.00 r
dneg1_reg/Q (dfnrb1) <-	0.32	7.32 r
dneg2_reg/D (dfnrb1)	0.00	7.32 r
data arrival time		7.32
clock clk' (rise edge)	15.00	15.00
clock source latency	-1.00	14.00
clk (in)	0.00	14.00 f
clkinv/I (inv0d2)	0.00	14.00 f
clkinv/ZN (inv0d2)	1.00 *	15.00 r
dneg2_reg/CP (dfnrb1)	0.00	15.00 r
library setup time	-0.09	14.91
data required time		14.91
-----		
data required time		14.91
data arrival time		-7.32
-----		
slack (MET)		7.59

The problem here is that the launch clock has used 1.0 and the capture clock has used -1.0. This is common path pessimism. So, if we turn on CRPR:

```
set timing_remove_clock_reconvergence_pessimism true
```

The path is now correct:

```
pt_shell> report_timing -input_pins -path_type full_clock -from dneg1_reg -to
dneg2_reg
*****
Report : timing
        -path full_clock
        -delay max
        -input_pins
        -max_paths 1
Design : duty_cycle_piclk
Version: V-2004.06
*****
```

Startpoint: dneg1\_reg (rising edge-triggered flip-flop clocked by clk')  
 Endpoint: dneg2\_reg (rising edge-triggered flip-flop clocked by clk')  
 Path Group: clk  
 Path Type: max

Point	Incr	Path
-----		
clock clk' (rise edge)	5.00	5.00
clock source latency	1.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dneg1_reg/CP (dfnrb1)	0.00	7.00 r
dneg1_reg/Q (dfnrb1) <-	0.32	7.32 r
dneg2_reg/D (dfnrb1)	0.00	7.32 r
data arrival time		7.32
clock clk' (rise edge)	15.00	15.00
clock source latency	-1.00	14.00
clk (in)	0.00	14.00 f
clkinv/I (inv0d2)	0.00	14.00 f
clkinv/ZN (inv0d2)	1.00 *	15.00 r
dneg2_reg/CP (dfnrb1)	0.00	15.00 r
clock reconvergence pessimism	2.00	17.00
library setup time	-0.09	16.91
data required time		16.91
-----		
data required time		16.91
data arrival time		-7.32
-----		
slack (MET)		9.59

I'm uncomfortable with this approach, however, because it uses a feature intended for OCV analysis (CRPR) for something unrelated to OCV. And to use it in conjunction with OCV will distort the numbers. I prefer the multiclock method.

## 10.4 Modeling duty cycle using set\_clock\_latency min/max

You can abuse the min-max version of set\_clock\_latency to do duty cycle analysis as well. The commands look like this:

```
set_clock_latency -source -fall -min -early \  
  [expr ($_duty_cycle_min - 0.5) * $_period] \  
  [get_clocks clk]  
set_clock_latency -source -fall -max -early \  
  [expr ($_duty_cycle_min - 0.5) * $_period] \  
  [get_clocks clk]  
set_clock_latency -source -fall -min -late \  
  [expr ($_duty_cycle_max - 0.5) * $_period] \  
  [get_clocks clk]  
set_clock_latency -source -fall -max -late \  
  [expr ($_duty_cycle_max - 0.5) * $_period] \  
  [get_clocks clk]
```

Note: If you're using a version earlier than 2004.06, PT won't let you do the above commands until you put it in "min-max" mode. Try something like this:

```
set_operating_conditions \  
-min cb13fs120_tsmc_max \  
-max cb13fs120_tsmc_max
```

The falling-edge to falling-edge setup path looks like this:

```
pt_shell> report_timing -input_pins -path_type full_clock -from dneg1_reg -to  
dneg2_reg
```

```
*****
```

```
Report : timing  
-path full_clock  
-delay max  
-input_pins  
-max_paths 1
```

```
Design : duty_cycle_piclk  
Version: V-2004.06
```

```
*****
```

```
Startpoint: dneg1_reg (rising edge-triggered flip-flop clocked by clk')  
Endpoint: dneg2_reg (rising edge-triggered flip-flop clocked by clk')  
Path Group: clk  
Path Type: max
```

Point	Incr	Path
-----		
clock clk' (rise edge)	5.00	5.00
clock source latency	1.00	6.00
clk (in)	0.00	6.00 f
clkinv/I (inv0d2)	0.00	6.00 f
clkinv/ZN (inv0d2)	1.00 *	7.00 r
dneg1_reg/CP (dfnrb1)	0.00	7.00 r
dneg1_reg/Q (dfnrb1) <-	0.32	7.32 r
dneg2_reg/D (dfnrb1)	0.00	7.32 r
data arrival time		7.32
clock clk' (rise edge)	15.00	15.00
clock source latency	-1.00	14.00
clk (in)	0.00	14.00 f
clkinv/I (inv0d2)	0.00	14.00 f
clkinv/ZN (inv0d2)	1.00 *	15.00 r
dneg2_reg/CP (dfnrb1)	0.00	15.00 r
clock reconvergence pessimism	2.00	17.00
library setup time	-0.09	16.91
data required time		16.91
-----		
data required time		16.91
data arrival time		-7.32
-----		
slack (MET)		9.59

Again, we're dependent on CRPR to make this work.

And I still prefer the multiclock method. The reason is that I don't like having to "overload" the early/late or min/max number with duty cycle information. I need these values for other effects,

and I hate to just add duty cycle into these numbers. By having separate duty cycle clocks, I can be confident that I'm not somehow applying duty cycle variation in cases where it doesn't belong just because I rolled it into early/late or min/max latency.